

# The Configuration Space (C-space)

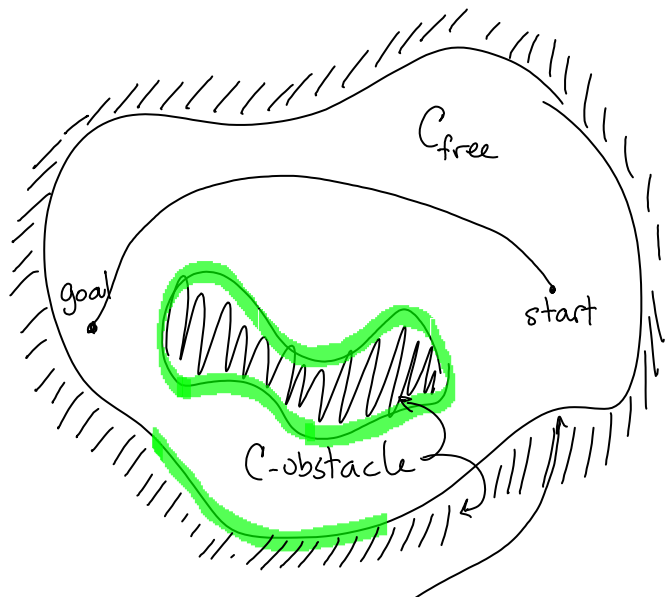
C-space is a (non-unique) space in which the robot is represented as a point

C-space is the set of points corresponding to every possible configuration of the robot, even those requiring overlap of bodies.

Robot motions correspond to continuous paths in

$C_{obstacle}$  = set of configs. corresponding to overlap

$C_{free}$  = set of configs. corresp. to no overlap or contact



$C_{contact}$  = set of configs where one or more objects touch but don't overlap.

Let C-space be denoted by  $C$ . Then

$$C = C_{obs} \cup C_{cont} \cup C_{free}$$

where  $C_{obs}$ ,  $C_{cont}$ ,  $C_{free}$  are mutually

where  $C_{obs}$ ,  $C_{free}$ ,  $C_{cont}$  are mutually exclusive.

All motion planning problems can be transformed into finding a continuous path in  $(C_{free} \cup C_{contact})$  connecting the starting and goal configs.

If we can approximate  $C$  as a discrete set of points with paths between them, then we can use discrete search methods!

For noncontact problems, it's usually best if the path is smooth (easier on actuators and structure of system).

Manipulation planning problems require a portion of the path to be in  $C_{cont}$ .

When dynamics is important planning may need to be done in state space  $X$ .

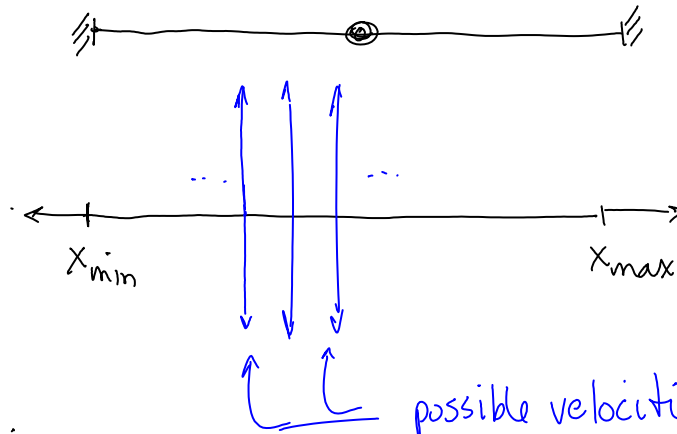
$$X = C \times T$$

↑  
↑  
space of velocities  
Cartesian product

At every point in  $C$ ,  $T$  is a "fiber" that is

At every point in  $C$ ,  $T$  is a "fiber" that is a space of velocities.

Bead on wire

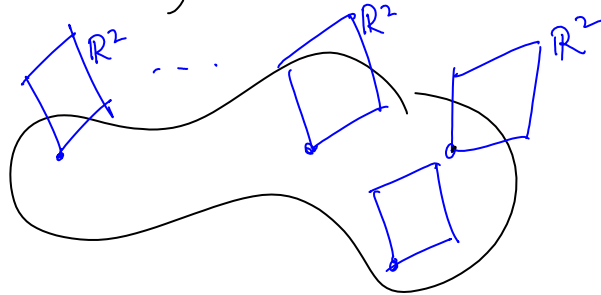


$$X = I' \times \mathbb{R}^1$$

if velocities are bounded,

then  $X = I' \times I' = D^2 = \text{the disk in } \mathbb{R}^2$

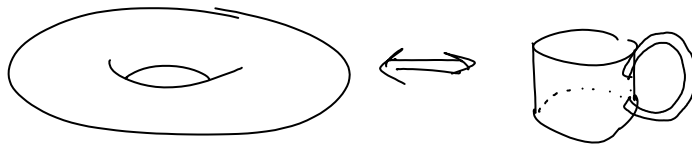
If  $C \subset \mathbb{R}^2$ ,  $T \subset \mathbb{R}^2$ , then  $X \subset \mathbb{R}^4$



A bit of topology

Homeomorphic - two shapes are homeomorphic if one can be deformed smoothly into the

other.



torus

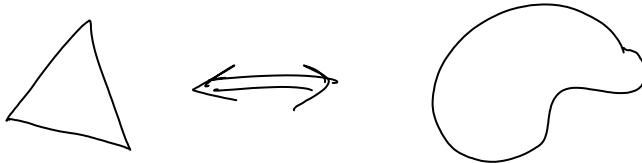
mug

both have one hole!

solid  
sphere



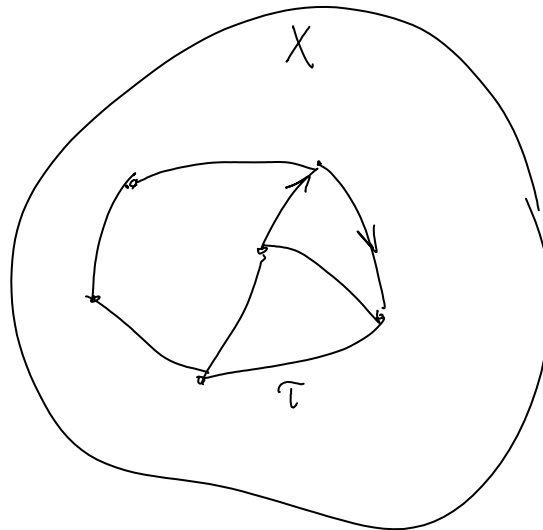
solid  
cube



### Topological Graphs

every vertex  
is in  $X$

every edge  
is in  $X$ .



More precisely

every edge map  $I^1$  onto a curve in  $X$

$$\tau : [0, 1] \rightarrow X$$

(i.e., think of  $\tau$  as  $\tau(t)$ , such that

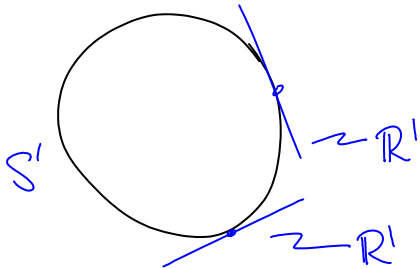
$$\tau(t) \in X \quad \forall t$$



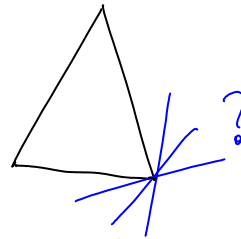
## Manifolds - Most C-spaces are manifolds

A set  $M$  is a manifold if  $\forall x \in M$ ,  $M$  is locally Euclidean.

This means there is a well-defined tangent plane at each point of  $M$



Circle is a 1D manifold



Triangle is not a 1D manifold

You need coord. frame on  $M$  for motion planning.

For  $S'$ , you could use a single parameter  $\theta \in [0, 2\pi)$

Sometimes it is more convenient to embed  $M$  in a

higher-dimensional space. e.g.

$S'$  can be embedded in  $\mathbb{R}^2$

$$S' = \{(x, y) \in \mathbb{R}^2 \mid x^2 + y^2 = 1\} \leftarrow \begin{array}{l} \text{2 dof reduced} \\ \text{to 1 by 1 eq.} \end{array}$$

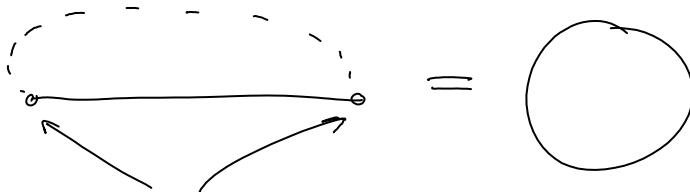
One could also embed in  $\mathbb{R}^3$ , but then you need 2 eqs. Here's a trivial example:

$$S' = \{(x, y, z) \in \mathbb{R}^3 \mid x^2 + y^2 = 1, z = 14\}$$

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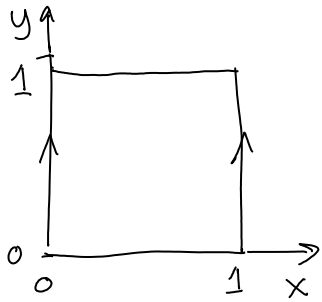
Identifications - must handle these carefully in motion planning.

$S' = I'$  with endpoints "identified"

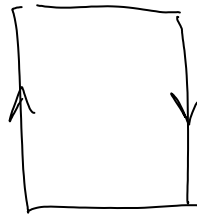


Treat those points as if they were identical.

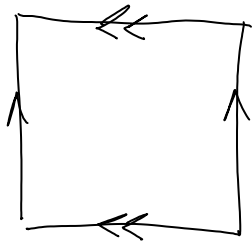
Some well-known identifications of  $I' \times I'$



$(0,y) \sim (1,y)$   
Cylinder

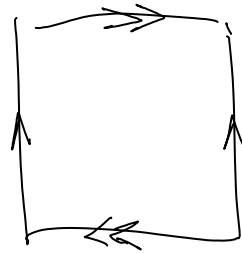


$(0,y) \sim (1,1-y)$   
Möbius Strip



$(0,y) \sim (1,y)$   
 $(x,0) \sim (x,1)$

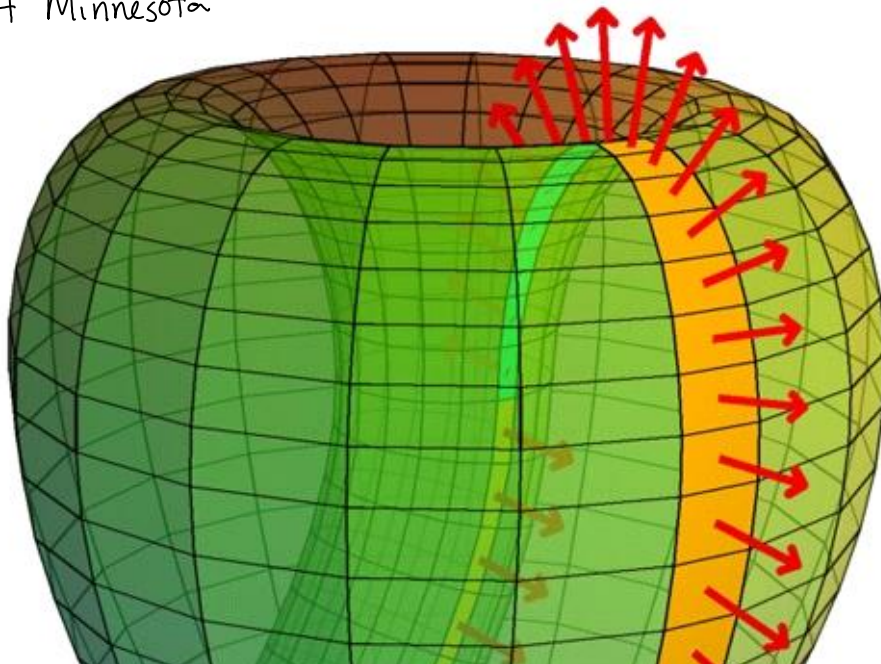
Torus



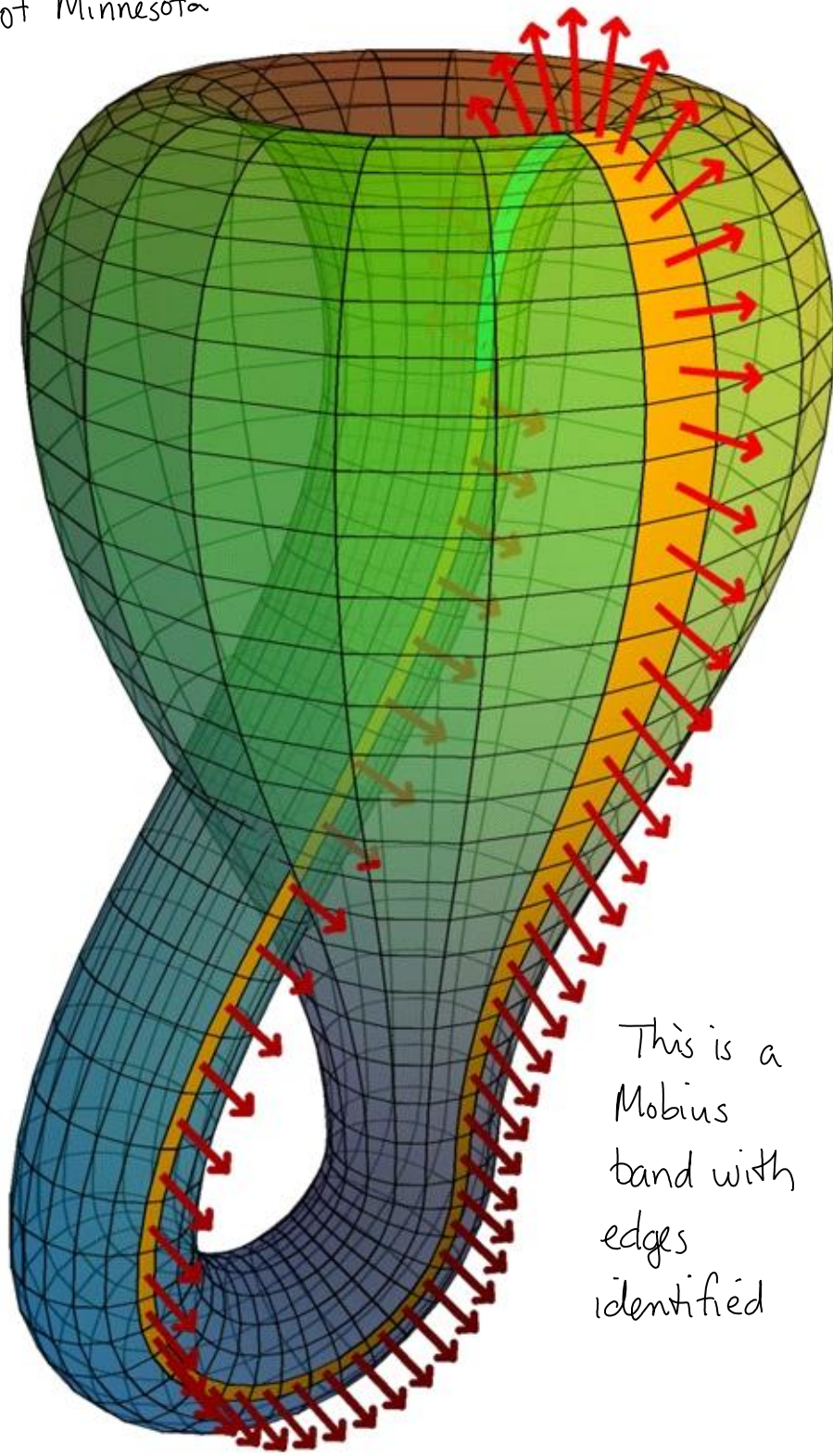
Klein Bottle



From Institute for Mathematics and its Application  
Univ of Minnesota



Univ of Minnesota



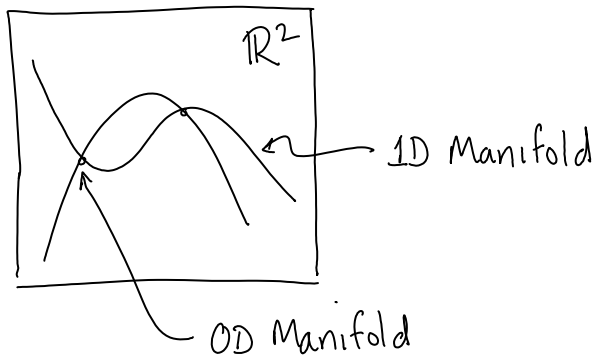
This is a  
Möbius  
band with  
edges  
identified

Higher-dimensional Manifolds



e.g.  $S^n = \{x \in \mathbb{R}^n \mid \|x\| = 1\}$

Typically  $m$  simultaneous equations in  $n > m$  variables yields a manifold



Intersection is a manifold of 2 distinct points  
(a point is considered a manifold of dimension zero)

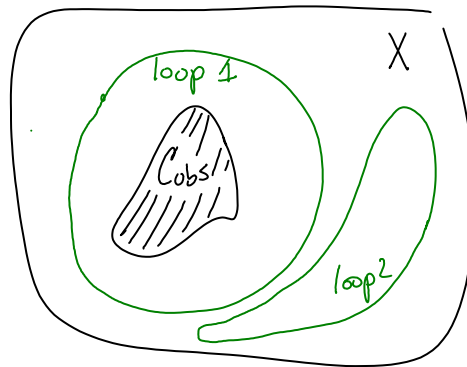
### Simply connected space

Every loop can be contracted to a point

Loop 1 cannot be contracted

Loop 2 can be

$X$  is not simply connected!



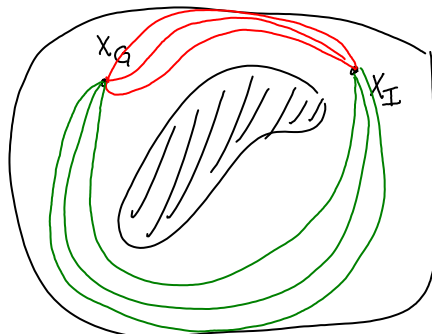
In a simply connected space all paths can be

In a simply connected space all paths can be morphed into any other path.

i.e. All paths are homotopic

" " " in the same homotopy class

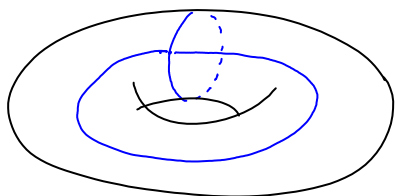
Green paths cannot be smoothly morphed into the red paths



Are there more classes of paths?

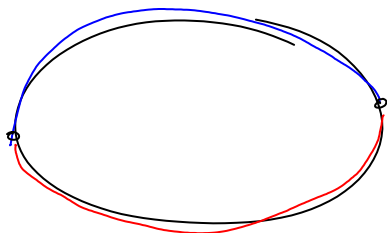
What about paths that encircle the obstacle?

Is the torus simply connected?



No. Not simply connected.

Is  $S^1$  simply connected?



No!

An example of homotopic paths:

$$r_1(s) = \begin{bmatrix} s - 1/2 \\ (s - 1/2)^2 \end{bmatrix}$$

$$r_2(s) = \begin{bmatrix} s - 1/2 \\ 1/2 - (s - 1/2)^2 \end{bmatrix}$$

$$v_1(s) = \left[ (s - \frac{1}{2})^2 \right]$$

$$v_2(s) = \left[ \frac{1}{2} - (s - \frac{1}{2})^2 \right]$$

$$\tau_1(0) = \begin{bmatrix} -\frac{1}{2} \\ \frac{1}{4} \end{bmatrix}$$

$$0 \leq s \leq 1$$

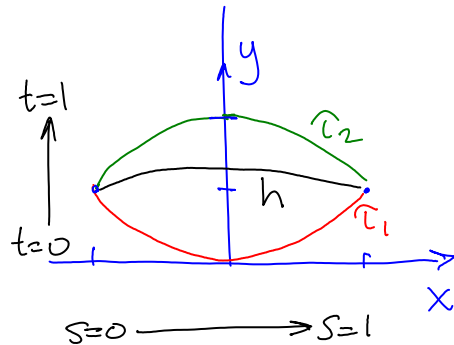
$$\tau_2(0) = \begin{bmatrix} -\frac{1}{2} \\ \frac{1}{4} \end{bmatrix}$$

$$\tau_1(1) = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{4} \end{bmatrix}$$

$$\tau_2(1) = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{4} \end{bmatrix}$$

$$\tau_1(\frac{1}{2}) = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$\tau_2(\frac{1}{2}) = \begin{bmatrix} 0 \\ \frac{1}{2} \end{bmatrix}$$



Define the morphing function,  $h(s,t)$

$$h(s,t) = (1-t) \tau_1(s) + t \tau_2(s)$$

This function must be continuous in  $t$  and  $s$ ,

i.e.:

- $h(\alpha,\beta)$  must exist at every  $\alpha,\beta$

For each  $\alpha,\beta$ , we must have:

- $\lim_{\alpha,\beta \rightarrow st} h(s,t)$  must exist at every  $\alpha,\beta$

- $\lim_{\alpha,\beta \rightarrow st} h(s,t) = h(\alpha,\beta)$



## LaValle 4.2 Defining C-space

$$\text{2D rigid bodies} - \mathcal{C} = \mathbb{R}^2 \times S^1 = SE(2)$$

$(x,y)$

$\theta$

2D rotations  
special Euclidean

space space

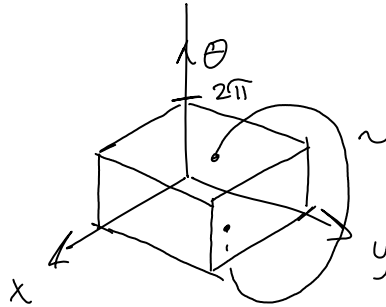
$$u = \begin{bmatrix} x \\ y \\ \theta \end{bmatrix}$$

$$S^1 = [0, 2\pi) / \sim$$



three parameters represent three degrees of freedom

A chunk of  $\mathcal{C}$  with bounds on  $(x, y)$ . Note identification of top & bottom surfaces.



Other representations:

base on  $T = \left[ \begin{array}{cc|c} a & c & x \\ b & d & y \\ \hline 0 & 0 & 1 \end{array} \right]$

$$u = [x \ y \ a \ b \ c \ d]^T \Rightarrow u \in \mathbb{R}^6$$

$$\text{but constraints: } a^2 + b^2 = 1$$

$$c^2 + d^2 = 1$$

$$ac + bd = 0$$

$\therefore \mathcal{C}$  is a 3D "variety" embedded in  $\mathbb{R}^6$

You could also take advantage of the fact that  $a=d$  and  $b=-c$ ,

$$\text{i.e. } R \in SO(2)$$

$$R = \begin{bmatrix} a & -b \\ b & a \end{bmatrix}$$

and  $a^2 + b^2 = 1$

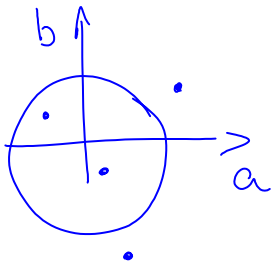
Determinant must = +1

special orthogonal group of  $2 \times 2$  matrices, Columns (and rows) are orthogonal unit vectors

$$u = [x \ y \ a \ b]^T$$

$\mathcal{C}$  is a 3-D surface embedded in  $\mathbb{R}^4$

Which representation of  $\mathcal{C}$ -space should you use?



Choosing points in  $(a,b)$ -space requires subsequent adjustment to satisfy  $a^2 + b^2 = 1$

Sampling  $\theta$  directly is simpler and can easily provide uniformly distributed values of  $\theta$ .

3D Rigid Bodies

$$\mathcal{C} = SE(3) = \mathbb{R}^3 \times SO(3)$$

$\uparrow$                        $\uparrow$   
 $(x, y, z)$                $(?)$

Elements of  $SO(3)$

$$\begin{bmatrix} a & d & a \end{bmatrix} \quad \begin{cases} a^2 + b^2 + c^2 = 1 \\ ad + be + cf = 0 \end{cases}$$

$$\begin{bmatrix} b & e & h \\ c & f & i \end{bmatrix} \iff \begin{cases} d^2 + e^2 + f^2 = 1 \\ g^2 + h^2 + i^2 = 1 \end{cases} \quad \begin{cases} dg + eh + fi = 0 \\ ga + hb + ic = 0 \end{cases}$$

Determinant must equal +1

$SO(3)$  is a 3-dimensional space embedded in  $\mathbb{R}^9$

Alternative representation:

① Euler angles or roll-pitch-yaw: 3 angles

② Quaternions: scalar & 3-vector 

Quaternions are preferred despite extra parameter, because mathematical properties are superior - i.e. no singularities, which improves uniformity of sampling

Let  $h$  denote a quaternion

$$h = \underbrace{a}_{\text{scalar}} + \underbrace{bi + cj + dk}_{\text{vector}} ; \quad a, b, c, d \in \mathbb{R}$$

$$i^2 = j^2 = k^2 = ijk = -1$$

$$ij = k, \quad jk = i, \quad ki = j$$

quaternion multiplication  $\equiv$  product of rotation matrices

$\therefore$  not commutative

Let  $h_1, h_2$  be quaternions

$$h_1 = a_1 + n_1 \quad h_2 = a_2 + n_2$$

↖ (eq. 4.19)

$$h_1 \cdot h_2 = (a_1 a_2 - n_1 \cdot n_2) + (a_1 n_2 + a_2 n_1 + n_1 \times n_2)$$

↑  
dot product

Compare computational cost

$$R_1 R_2 \Rightarrow 27 \text{ mults} + 18 \text{ adds}$$

$$h_1 h_2 \Rightarrow 16 \text{ mults} + 12 \text{ adds}$$

Note that the extra parameter gives some duplication:

$$\theta \curvearrowright n = -\theta \curvearrowright -n$$

(Think of the quaternion as an operation on points, NOT as an axis  $n$  fixed in a body which is rotated about  $n$ )

Given  $n, \theta$ , determine  $h$  (eq. 4.21):

$$h = \underbrace{\cos\left(\frac{\theta}{2}\right)}_a + \underbrace{(n_x \sin\left(\frac{\theta}{2}\right))}_b i + \underbrace{(n_y \sin\left(\frac{\theta}{2}\right))}_c j + \underbrace{(n_z \sin\left(\frac{\theta}{2}\right))}_d k$$

Determine the rotation matrix (eq. 4.20):

$$R(h) = \begin{bmatrix} 2(a^2 + \underline{b^2}) - 1 & 2(bc - \underline{ad}) & 2(bd + \underline{ac}) \\ 2(bc + \underline{ad}) & 2(a^2 + \underline{c^2}) - 1 & 2(cd - \underline{ab}) \\ 2(bd - \underline{ac}) & 2(cd + \underline{ab}) & 2(a^2 + \underline{d^2}) - 1 \end{bmatrix}$$



Note that  $R(h)$  is quadratic in the elements of  $h$ .

Convert  $R$  to  $h$  (eqs. 4.24 - 4.30)

$$R = \begin{bmatrix} r_{11} & r_{12} & r_{13} \\ r_{21} & r_{22} & r_{23} \\ r_{31} & r_{32} & r_{33} \end{bmatrix}$$

eqs. (4.24 - 4.27)

$$a = \frac{1}{2} \sqrt{r_{11} + r_{22} + r_{33} - 1}$$

$$b = \frac{r_{32} - r_{23}}{4a}$$

$$c = \frac{r_{13} - r_{31}}{4a}$$

$$d = \frac{r_{21} - r_{12}}{4a}$$

Not valid if  $a = 0$

If  $a = 0$ , then use equations (4.28 - 4.30)

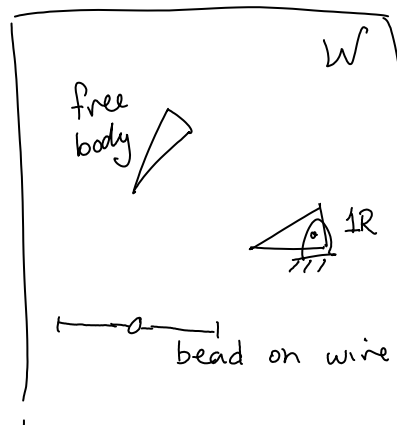
These also have possible zero denominators, but the problem can be corrected (eqs. 3.39 - 3.41).

C-space for multiple bodies

Let  $C_i$  be the C-space of body  $i$

$$C = C_1 \times C_2 \times \dots \times C_n$$

where  $n$  is # bodies

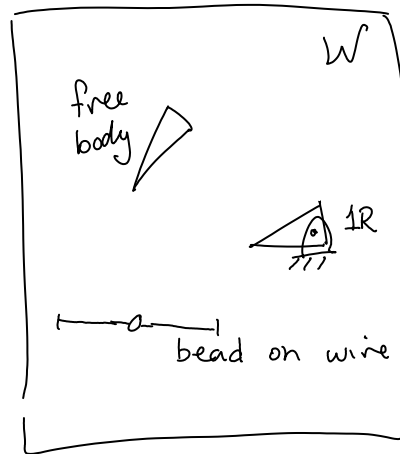




Let  $C_i$  be the C-space  
of body  $i$

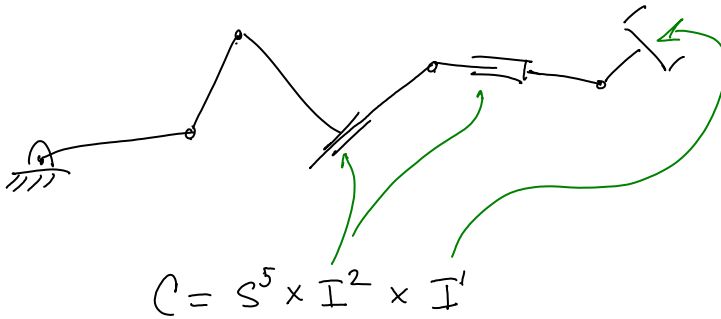
$$C = C_1 \times C_2 \times \dots \times C_n$$

where  $n$  is # bodies



$$C = I^1 \times S^1 \times SE(3)$$

Note: for every robot, it is practical to treat nearly every revolute joint as having C-space =  $I^1$ . The hard part is dealing with arbitrary orientations of free bodies and bodies connected by spherical joints.



$$C = S^5 \times I^2 \times I^1$$

Assuming revolute joints have limits,

$$C = I^8$$

$\therefore$  for planning purposes, just sample 8 intervals to get a point in C-space.

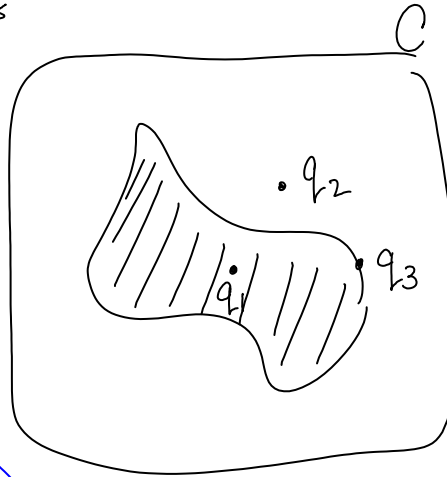
If base of robot is floating, then

$$C = SE(3) \times I^8$$

### 4.3 The C-space Obstacle, $C_{obs}$

Let  $q \in C$  denote  
a config. of the robot

Let  $A(q)$  be the robot  
in config  $q$ .

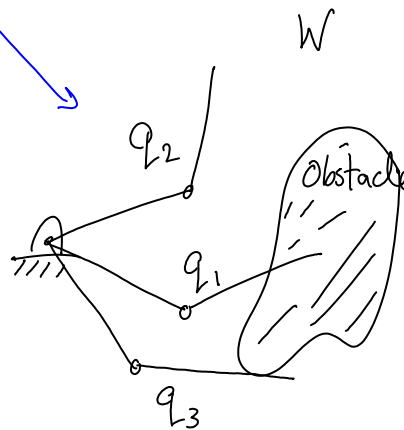


$$C_{obs} = \{q \in C \mid A(q) \cap \text{Int}(O) \neq \emptyset\}$$

$$C_{cont} = \{q \in C \mid A(q) \cap O \neq \emptyset, \\ A(q) \cap \text{Int}(O) = \emptyset\}$$

equivalently,

$$C_{cont} = \partial C_{obs}$$



$$C_{free} = C \setminus (C_{obs} \cup C_{cont}) = C \setminus \text{cl}(C_{obs})$$

↑ closure of a set

#### 4.3.2 Explicit representations of $C_{cont} = \partial C_{obs}$

(LaValle treats  $C_{obs}$  and  $C_{obs} \cup C_{cont}$ )

Definition: Minkowski difference of two sets in  $\mathbb{R}^n$

$$X \ominus Y = \{x - y \in \mathbb{R}^n \mid x \in X, y \in Y\}$$

Definition: Minkowski sum

$$X \oplus Y = \{x+y \in \mathbb{R}^n \mid x \in X, y \in Y\}$$

Note:  $X \ominus Y = X \oplus (-Y)$  where  $-Y$  is the set of negated elements of  $Y$ .

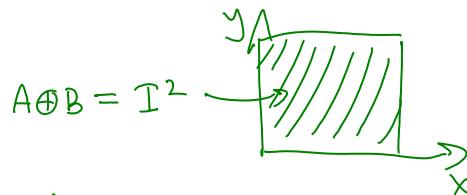
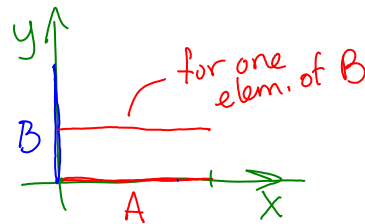
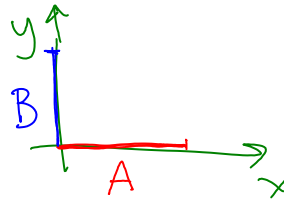
Example in  $\mathbb{R}^2$

Let  $A = I = ([0, 1], 0) =$  interval on x-axis

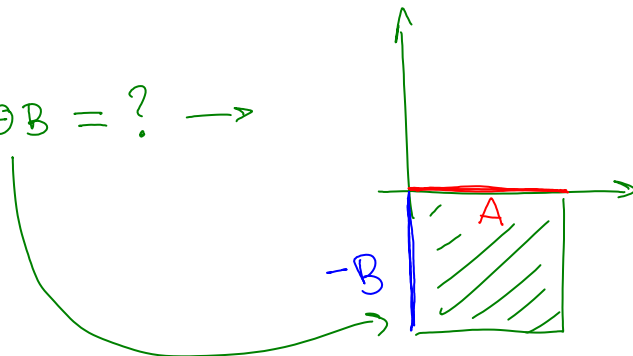
$B = I = (0, [0, 1]) =$  interval on y-axis

What is  $A \oplus B$ ?

Add every element of  $A$   
to every element of  $B$



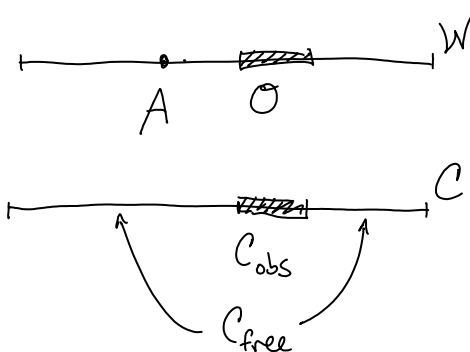
$A \ominus B = ? \rightarrow$



Using  $\Theta$  to determine  $C_{free}$  &  $C_{obs}$

Simplest case: A is a particle

1D world

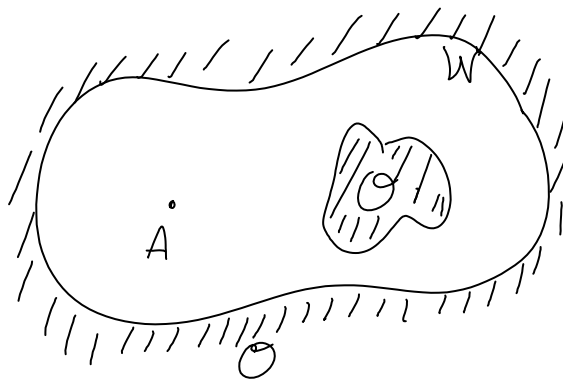


2D world

Again  $C_{obs} = \text{Int}(\Theta)$

$$C_{cont} = \partial\Theta$$

$$C_{free} = W \setminus \text{cl}(\Theta)$$



Same for 3D worlds

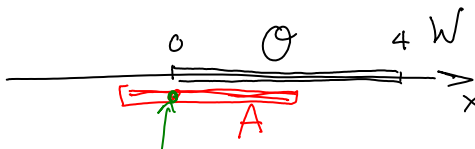
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Robots of finite extent:

1D:

$$\text{Let } A = [-1, 2]$$

$$\Theta = [0, 4]$$

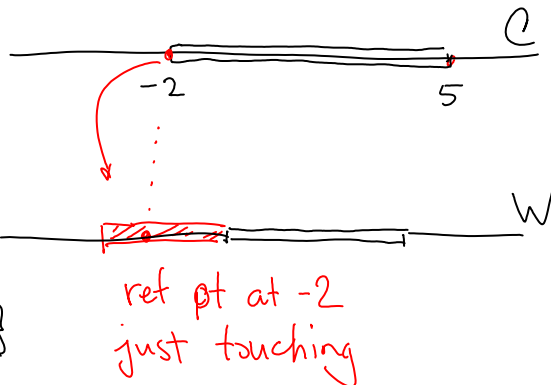


To create  $C_{obs}$ ,  $C_{cont}$ ,  $C_{free}$ ,  
 we need a ref. pt. on A.  
 (Choose the zero point)

$\mathcal{C}$  represents possible positions of A by possible positions of the ref. point.

$$cl(C_{obs}) = \mathcal{O} \oplus -A = [0, 4] \oplus [-2, 1]$$

$$cl(C_{obs}) = [-2, 5]$$

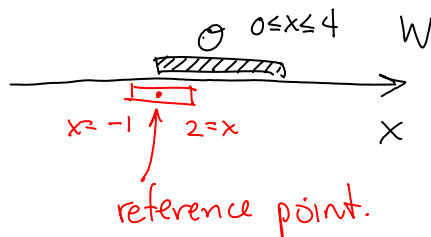


$$C_{free} = \{x \mid x < -2 \vee x > 5\}$$

$$C_{cont} = \{x \mid x = -2 \vee x = 5\}$$

Why do we define  $C_{obs}$  by  $\mathcal{O} \oplus (-A)$  instead of  $\mathcal{O} \oplus A$ ?

Consider 1D case:



Noncollision constraint  
 on left side of  $\mathcal{O}$

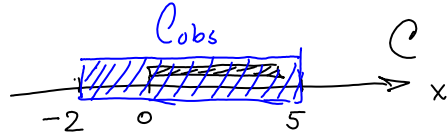
$$x + a < \sigma, \forall a \in A(x), \forall \sigma \in \mathcal{O}$$

$$\Rightarrow x < \sigma - a, \forall a \in A(x), \forall \sigma \in \mathcal{O}$$

This is Definition of  
 Minkowski difference

Worst case:  $x < -2$

Worst case:  $x < -2$

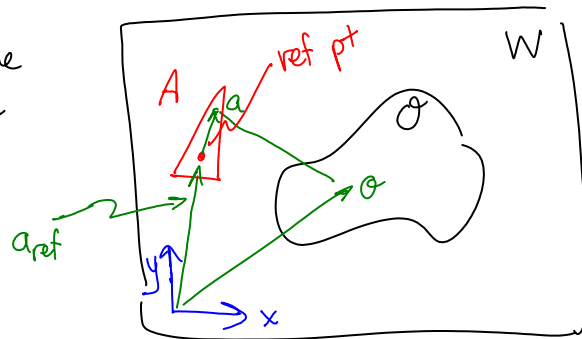


You can apply on right also

$$x > \sigma - a, \forall a \in A, \forall \sigma \in \mathcal{O}$$

Worst case:  $x > 5$

2D example works the same way, but constraints are more complicated



pt on robot,  $v$

$$v = a_{ref} + a, \quad a \in A(q)$$

$q \in C_{obs}$  if  $a_{ref} + a = \sigma$  for some  $\sigma \in \mathcal{O}$   
 for some  $a \in A$

$q \notin C_{obs}$  if  $a_{ref} \neq \sigma - a, \forall a \in A(q), \forall \sigma \in \mathcal{O}$

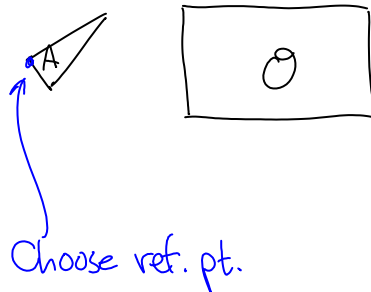
Minkowski difference.

Planar Translation Only Robot.

$C_{free}$  is set of

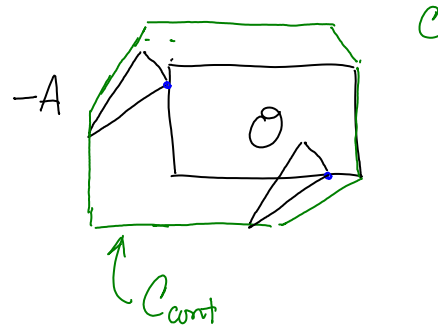
positions of ref.

point  $\ni A \cap \mathcal{O} = \emptyset$



$$cl(C_{obs}) = O \oplus -A$$

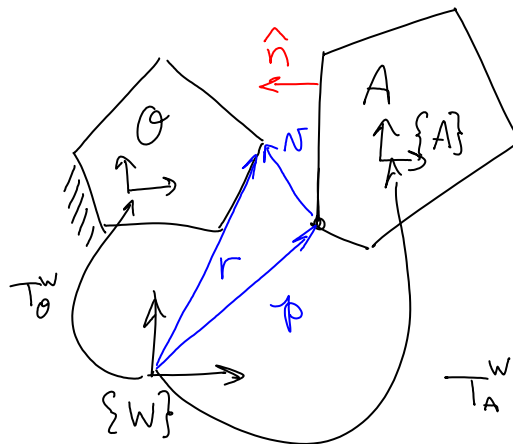
Add  $-A$  to every point in  $O$



Formulas for the boundary of  $C_{obs}$

for planar world with

$O$  &  $A$  defined as polygons



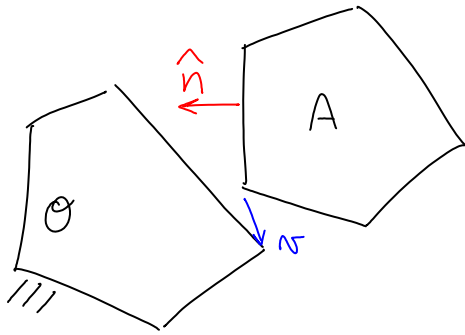
EV Contact -

edge of robot  
in contact with  
vertex of  $O$ .

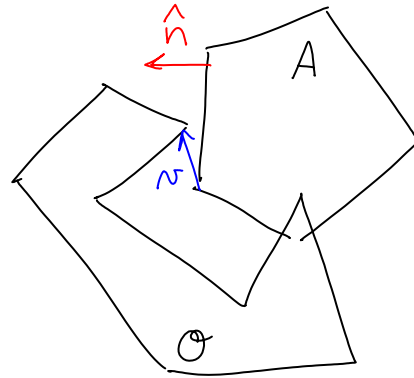
$\hat{n} \cdot N = 0$	<u>potential contact</u>
$\hat{n} \cdot N < 0$	<u>potential penetration</u>
$\hat{n} \cdot N > 0$	<u>local separation</u>

Why the qualifiers?





$\hat{n} \cdot n \leq 0$ ,  
but no contact  
or penetration

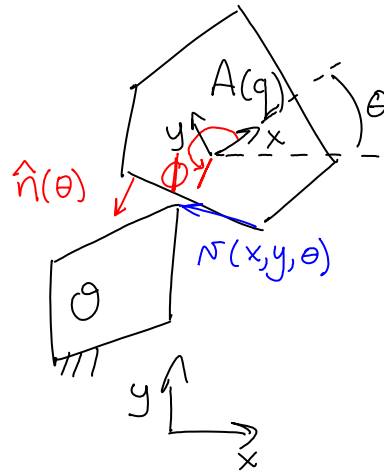


$\hat{n} \cdot n \geq 0$ , but  
penetration occurs  
(non-locally).

To obtain manifold of a given facet, write  
contact condition as a function of  $q = \begin{bmatrix} x \\ y \\ \theta \end{bmatrix}$

Necessary condition  
for contact :

$$\hat{n}(\theta) \cdot N(x, y, \theta) = 0$$



Let  $\phi$  denote the  
angle from the body-fixed  
x-axis to the normal  
direction

$$\hat{n}(\theta) = \begin{bmatrix} \cos(\phi + \theta) \\ \sin(\phi + \theta) \end{bmatrix}$$

$$N^W = r^W - p^W = r^W - T_A^W p^A$$

← position of p in  
frame A



$$N(x, y, \theta) = \begin{bmatrix} r_x - c_\theta p_x^A + s_\theta p_y^A - x \\ r_y - s_\theta p_x^A - c_\theta p_y^A - y \end{bmatrix}$$

position of obstacle vertex in  $W$ 
position of origin of robot frame in  $W$

Substitute into  $\hat{n}(\theta) \cdot N(x, y, \theta) = 0$

$$0 = \cos(\theta + \phi) [r_x - c_\theta p_x^A + s_\theta p_y^A - x] + \sin(\theta + \phi) [r_y - s_\theta p_x^A - c_\theta p_y^A - y]$$

(This eq. defines a 2D "variety" in 3D C-space  $(x, y, \theta)$ ).

Nonlinear trigonometric polynomial in  $\theta$   
(e.g.,  $\cos(\theta + \phi) \cdot \cos(\theta)$ )

Linear in  $(x, y)$

$$\boxed{ax + by + c = 0}$$

where  $a = -\cos(\theta + \phi)$ ,  $b = -\sin(\theta + \phi)$ ,

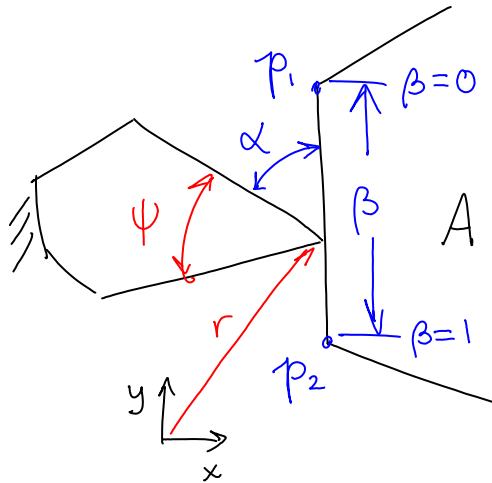
$$\text{and } c = \cos(\theta + \phi) [r_x - c_\theta p_x^A + s_\theta p_y^A] \\ + \sin(\theta + \phi) [r_y - s_\theta p_x^A - c_\theta p_y^A]$$

Setting this contact eq. = 0 yields a 2D manifold in a 3D space ( $\mathbb{R}^2 \times S^1$ ).

Pieces of this manifold appear in Brost's Cobs

How do we get the pieces?

Another useful formulation ...



$$0 \leq \alpha \leq \pi - \psi$$

$$0 \leq \beta \leq 1$$

$r$  must be on the line segment  $\overline{p_1 p_2}$

$$\begin{bmatrix} p_{ix}^w \\ p_{iy}^w \\ 1 \end{bmatrix} = \begin{bmatrix} c_\theta & -s_\theta & x \\ s_\theta & c_\theta & y \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} p_{ix}^A \\ p_{iy}^A \\ 1 \end{bmatrix}$$

Substitute  $p_1^w, p_2^w$  into:

$$i = 1, 2$$

given in geom. model

$$(1-\beta)p_1^w(x, y, \theta) + \beta p_2^w(x, y, \theta) = r$$

$$0 \leq \beta \leq 1$$

$$0 \leq \alpha \leq \pi - \psi$$

If  $A \neq \emptyset$  are convex, then this is a patch of Cobs

If  $A$  or  $\emptyset$  are nonconvex, then some of this patch may be cut away by nonlocal interpenetrations!

Other contact types.

$\hat{n}$  ////

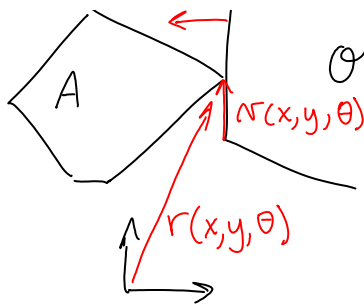
VE contact :

Same procedure, but

$\hat{n}$  is constant

$N$  is  $N(x,y,\theta)$

$r$  is  $r(x,y,\theta)$



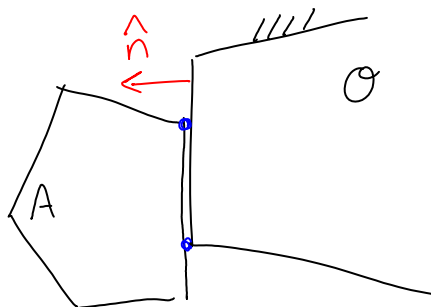
VE  $\Rightarrow$  2D facet of Cobs

EE contact :

Treat as two

contacts of type

EV or VE



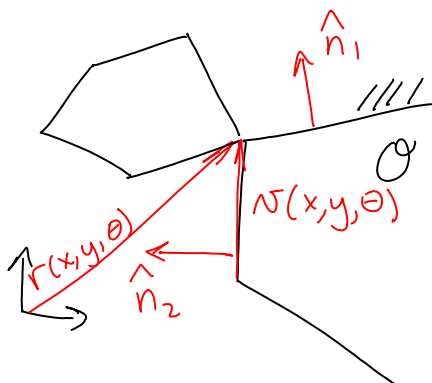
EE  $\Rightarrow$  1D facet of Cobs

VV contact :

Treat as two

contacts of type

EV or VE



VV  $\Rightarrow$  1D facet of Cobs

You can choose (EV, EV) or (VE, VE) or ~~(EV, VE)~~



~~doesn't always work~~

allowed



## Collision Checking for Polygons

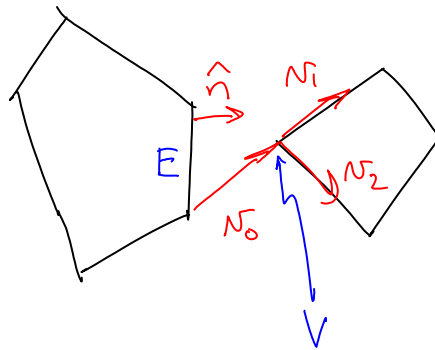
Assume polygons are convex  
(can partition non-convex ones)

A pair of convex polygons are separated iff  $\exists$  a vertex/edge pair  $(E, V) \ni$

$$\hat{n} \cdot N_0 > 0$$

$$\hat{n} \cdot N_1 > 0$$

$$\hat{n} \cdot N_2 > 0$$



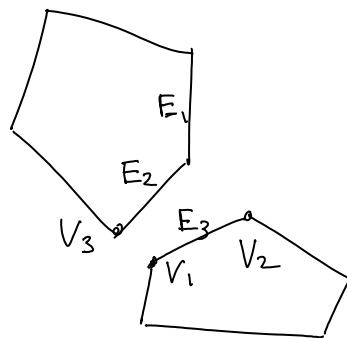
Example:  $\longrightarrow$

$(E_1, V_1)$  fails test

$(E_2, V_2)$  fails

$(E_3, V_3)$  succeeds

$(E_2, V_1)$  succeeds too



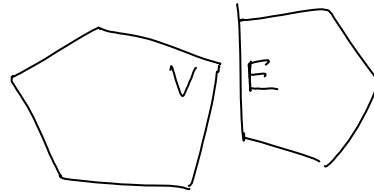
## Fast collision checking - intuition

If relative config does not change much, and

polygons were not in collision at previous check,  
it should be "easy" to find the definitive vertex/edge  
pair.

Exact collision checking w/o "hot start"

Consider 1 EV pair  $\rightarrow$



$$\text{let } H_A = H_1 \cup H_2 \cup H_3$$

$$\begin{aligned} \text{where } H_1 &= \{q \in C \mid \hat{n} \cdot n_0 \leq 0\} \\ H_2 &= \{q \in C \mid \hat{n} \cdot n_1 \leq 0\} \\ H_3 &= \{q \in C \mid \hat{n} \cdot n_2 \leq 0\} \end{aligned} \left. \vphantom{\begin{aligned} H_1 \\ H_2 \\ H_3 \end{aligned}} \right\} \begin{array}{l} \text{negations of} \\ \text{previously} \\ \text{used} \\ \text{inequalities} \end{array}$$

$H_A$  is a superset of  $C_{obs}$ , i.e.

if  $q \in H_1$  or  $H_2$  or  $H_3$ , then  $q$  might be  
an element of  $C_{obs}$

if  $q \in \bar{H}_1$  and  $\bar{H}_2$  and  $\bar{H}_3$ , then  $q \in C_{free}$

However, if  $q \in H_A$  for every EV and VE  
pair, then  $q \in C_{obs}$

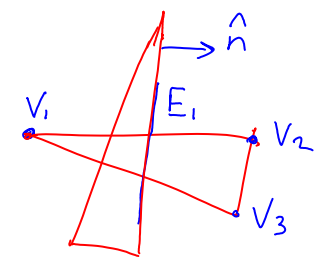
Testing: start checking all EV and VE pairs

until one is not satisfied. Then  $q \in C_{free}$ .  
 If all are satisfied, then  $q \in C_{obs}$ .

Does this work for this case?

$V_j E_i, i=1,2,3$   
 $j=1,2,3$  satisfied

$V_j E_i, i=4,5,6$   
 $j=4,5,6$  satisfied



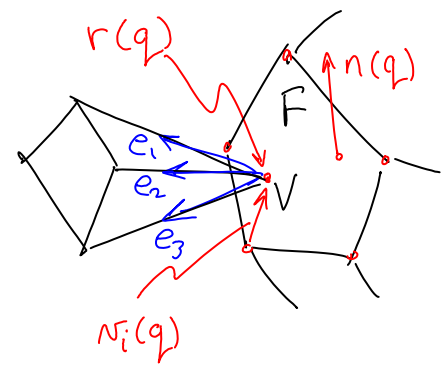
Yes! It works!

Simple exhaustive collision checking is  $O(m^2)$   
 where  $m$  is the maximum # of vertices  
 on each polygon.

More precisely # of predicates is  $2n_A n_O$  where  
 $n_A = \#$  vertices on robot  
 $n_O = \#$  vertices on obstacle

What about 3D?

Extension is analogous  
 but a bit more  
 complicated.



eg. Collision free  $q$  if :

$$\hat{n}_i(q) \cdot n_i(q) \geq 0$$

$$\hat{n}(q) \cdot e_i(q) \geq 0$$

Contact facet of  $C_{obs}$  defined by

$r(q)$  on plane of facet & within  
bounds defined by facet edges.

Collision checking can be formed by checking  
inequalities of all  $VF, FV, \& EE$  pairs.