

# **Frontiers of Network Science**

## **Fall 2018**

### **Class 7: Small World and BA Networks** **(Chapters 4 and 5 in Textbook)**

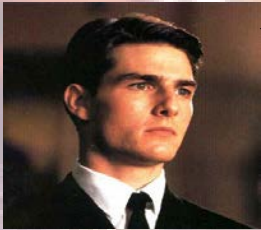
**Boleslaw Szymanski**

based on slides by  
Albert-László Barabási  
and Roberta Sinatra

## Examples

# ACTOR NETWORK

**Nodes:** actors  
**Links:** cast jointly



*Days of Thunder* (1990)  
*Far and Away* (1992)  
*Eyes Wide Shut* (1999)

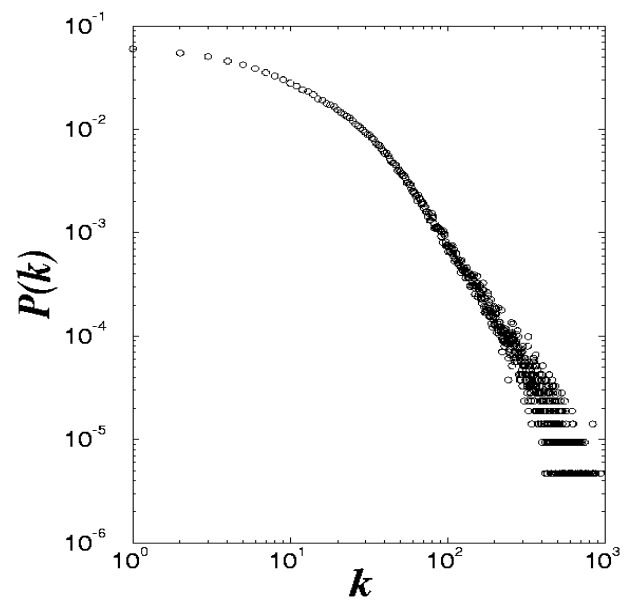


$N = 212,250$  actors

$\langle k \rangle = 28.78$

$P(k) \sim k^{-\gamma}$

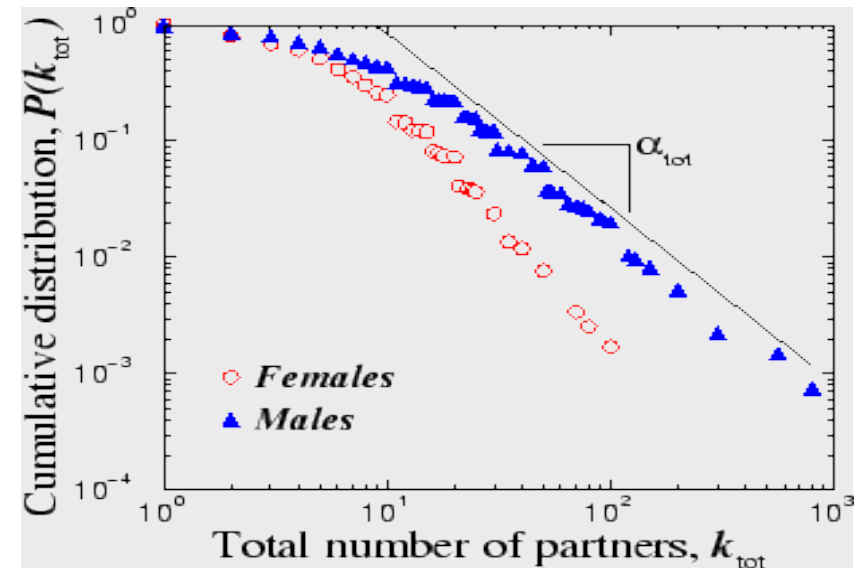
$\gamma = 2.3$





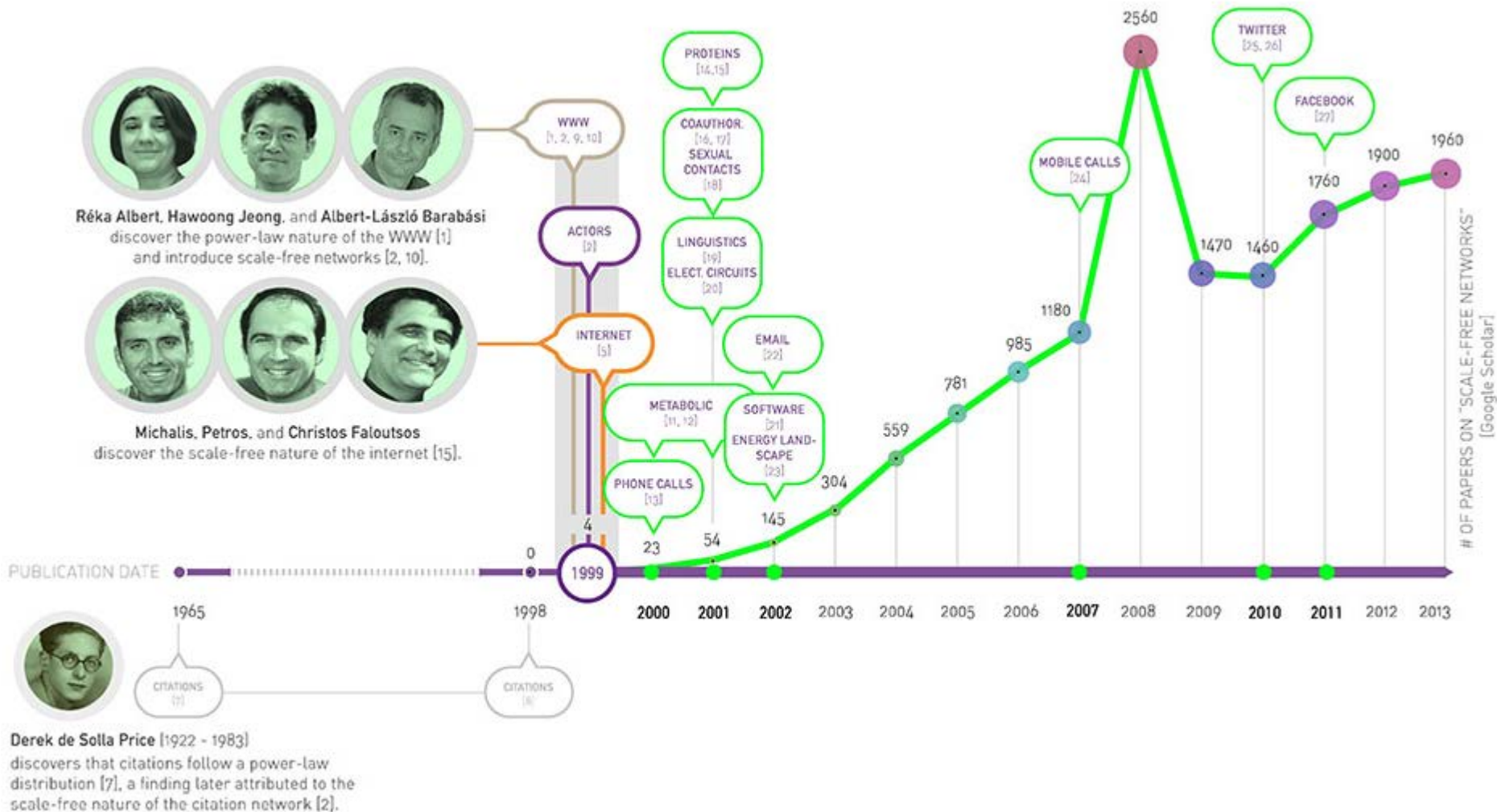
**Nodes:** people (Females; Males)

**Links:** sexual relationships



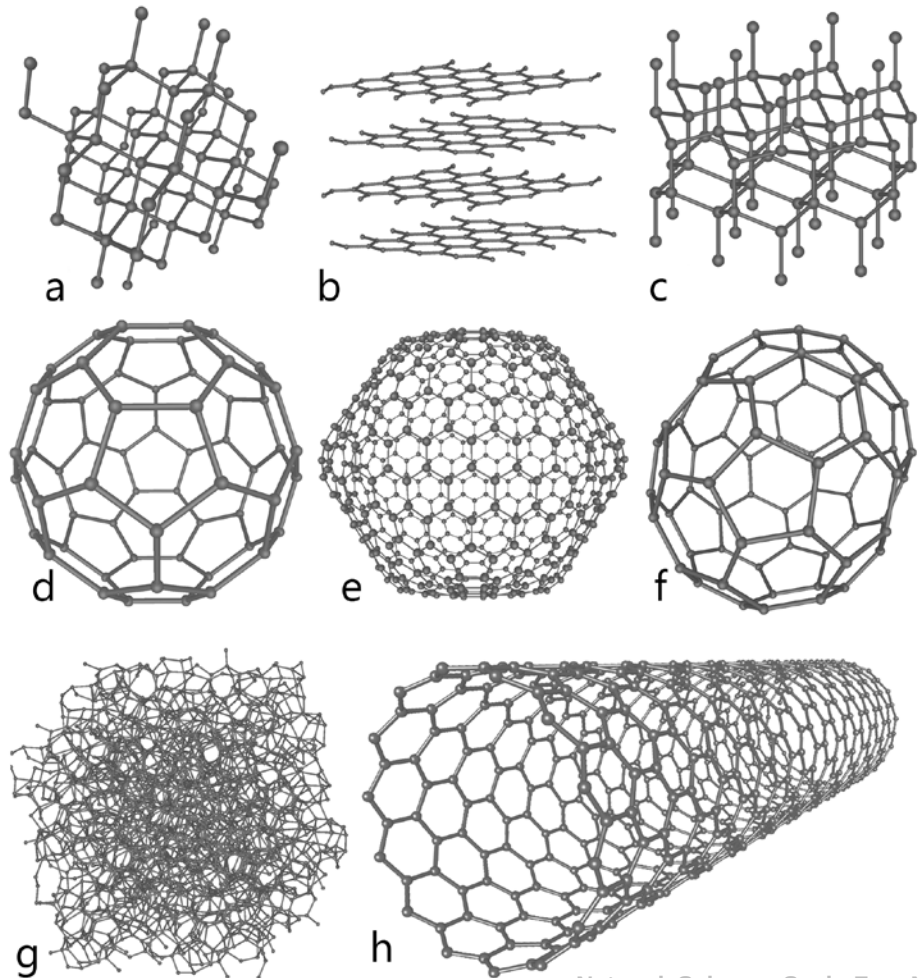
4781 Swedes; 18-74;  
59% response rate.

Liljeros et al. Nature 2001



# Not All Networks Are Scale-free

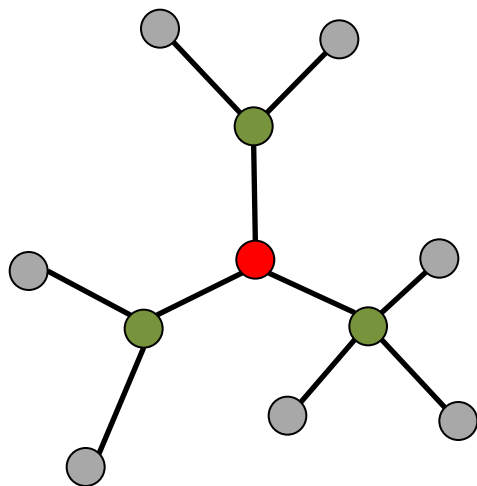
- *Networks appearing in material science, like the network describing the bonds between the atoms in crystalline or amorphous materials, where each node has exactly the same degree.*
- *The neural network of the *C.elegans* worm.*
- *The power grid, consisting of generators and switches connected by transmission lines*



# Ultra-small property

# DISTANCES IN RANDOM GRAPHS

Random graphs tend to have a tree-like topology with almost constant node degrees.



- number of first neighbors:

$$N_1 \cong \langle k \rangle$$

- number of second neighbors:

$$N_2 \cong \langle k \rangle^2$$

- number of neighbors at distance d:

$$N_d \cong \langle k \rangle^d$$

- estimation of maximum distance:

$$1 + \sum_{l=1}^{l_{\max}} \langle k \rangle^l = N \Rightarrow l_{\max} = \frac{\log N}{\log \langle k \rangle}$$

# SMALL WORLD BEHAVIOR IN SCALE-FREE NETWORKS

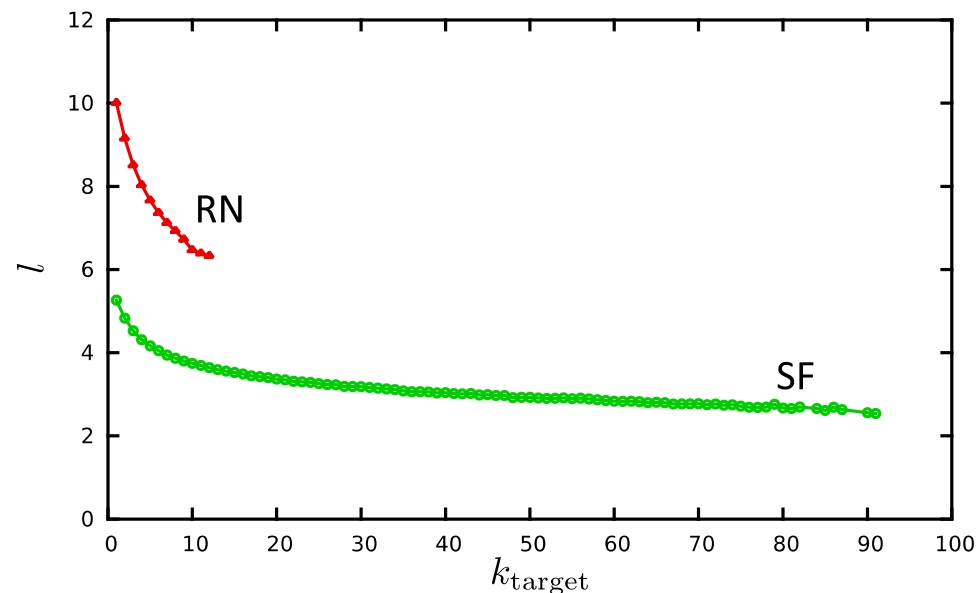
$$k_{\max} = k_{\min} N^{\frac{1}{\gamma-1}}$$

|                         |         |                                     |                  |                                                                                                                                                                                                                                                                                                    |
|-------------------------|---------|-------------------------------------|------------------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| Ultra<br>Small<br>World | < l > ~ | const.                              | $\gamma = 2$     | Size of the biggest hub is of order $O(N)$ . Most nodes can be connected within two layers of it, thus the average path length will be independent of the system size.                                                                                                                             |
|                         |         | $\frac{\ln \ln N}{\ln(\gamma - 1)}$ | $2 < \gamma < 3$ | The average path length increases slower than logarithmically. In a random network all nodes have comparable degree, thus most paths will have comparable length. In a scale-free network the vast majority of the path go through the few high degree hubs, reducing the distances between nodes. |
|                         |         | $\frac{\ln N}{\ln \ln N}$           | $\gamma = 3$     | Some key models produce $\gamma=3$ , so the result is of particular importance for them. This was first derived by Bollobas and collaborators for the network diameter in the context of a dynamical model, but it holds for the average path length as well.                                      |
| Small<br>World          |         | $\ln N$                             | $\gamma > 3$     | <u>T</u> he second moment of the distribution is finite, thus in many ways the network behaves as a random network. Hence the average path length follows the result that we derived for the random network model earlier.                                                                         |

Cohen, Havlin Phys. Rev. Lett. 90, 58701(2003); Cohen, Havlin and ben-Avraham, in Handbook of Graphs and Networks, Eds. Bornholdt and Shuster (Willy-VCH, NY, 2002) Chap. 4; Confirmed also by: Dorogovtsev et al (2002), Chung and Lu (2002); (Bollobas, Riordan, 2002; Bollobas, 1985; Newman, 2001

# We Are Always Close to the Hubs

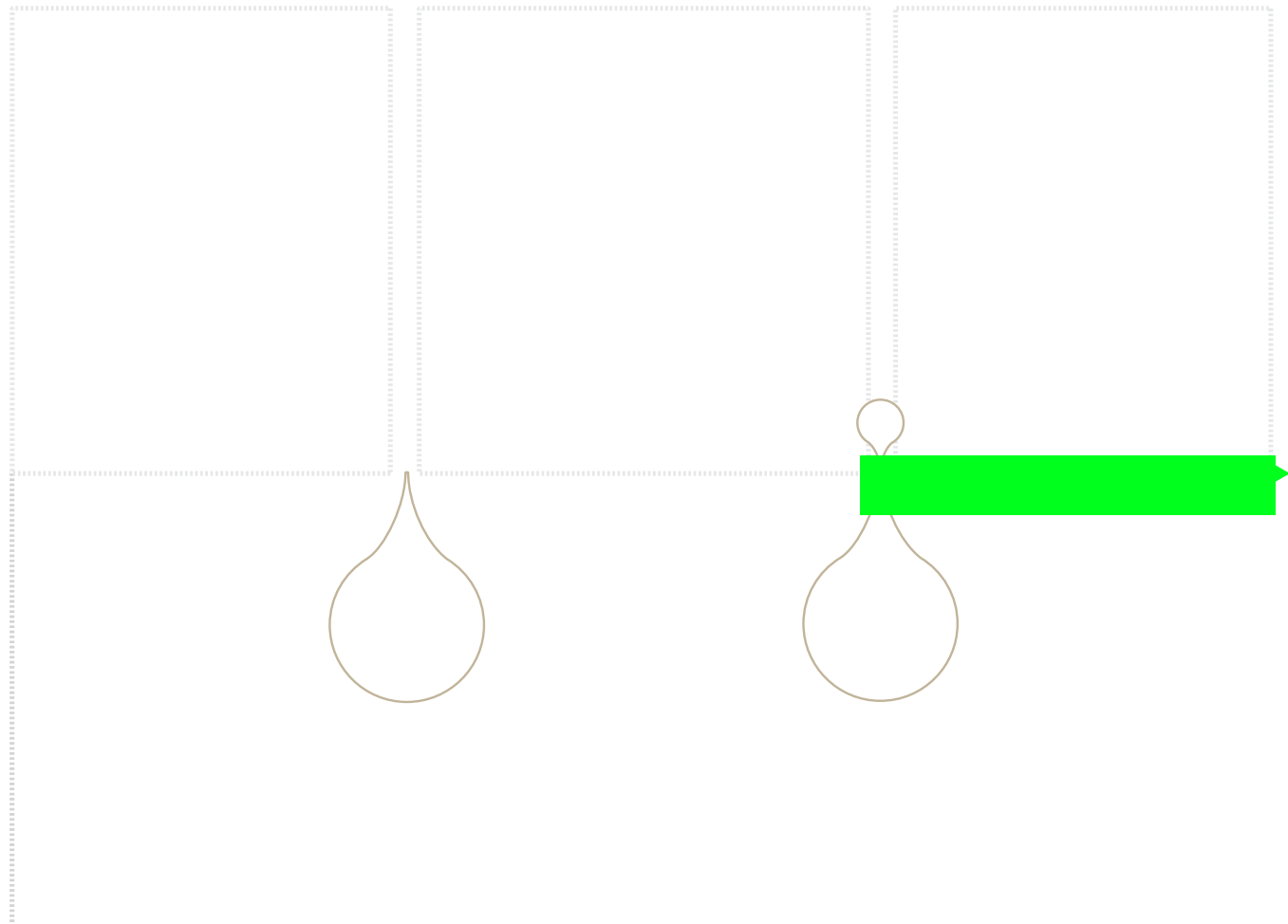
" it's always easier to find someone who knows a famous or popular figure than some run-the-mill, insignificant person."  
(Frigyes Karinthy, 1929)



Average person is less popular than this person's random friend!

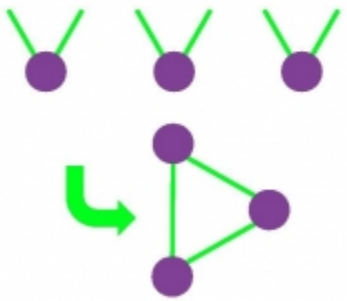
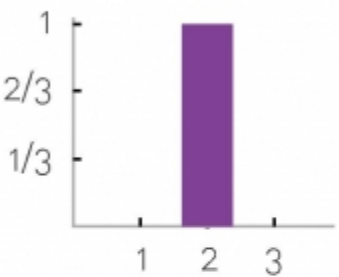
# The role of the degree exponent

# SUMMARY OF THE BEHAVIOR OF SCALE-FREE NETWORKS

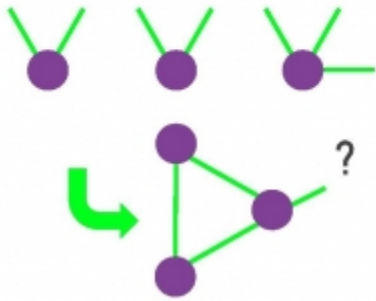
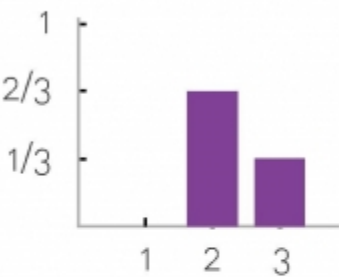


# Graphicality: No Large Networks for $\gamma < 2$

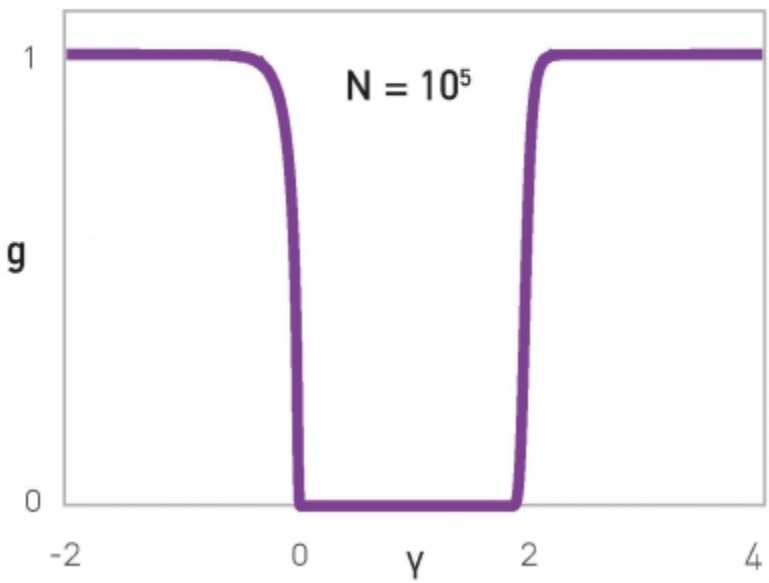
a. Graphical



b. Not Graphical



c.



In scale-free networks:  $k_{\max} = k_{\min} N^{\frac{1}{\gamma-1}}$

For  $\gamma < 2$ :  $1/(\gamma-2) > 1$

# Why Don't We See Networks with Exponents in the Range of $\gamma=4,5,6$ , etc?

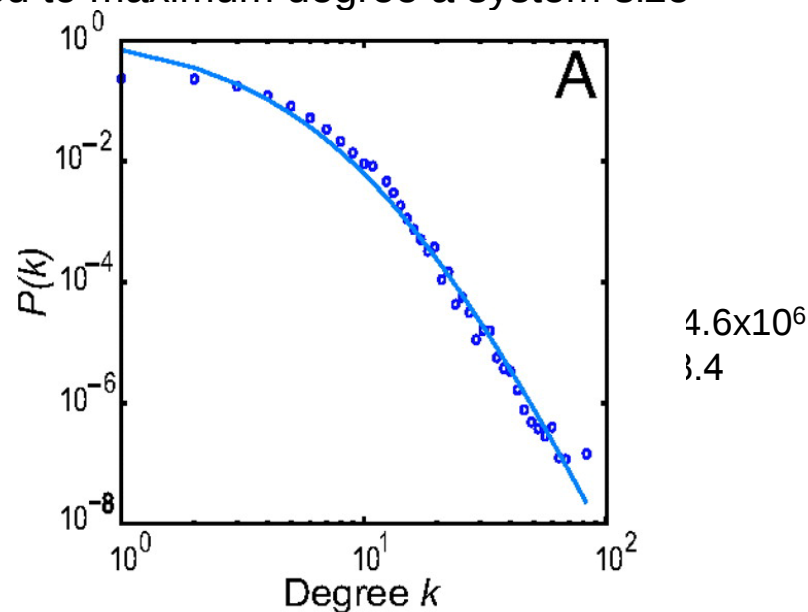
In order to document a scale-free networks, we need 2-3 orders of magnitude scaling.  
That is,  $K_{\max} \sim 10^3$

However, that constrains on the system size we require to document it.  
For example, to measure an exponent  $\gamma=5$ , we need to maximum degree a system size of the order of

$$K_{\max} = K_{\min} N^{\frac{1}{\gamma-1}}$$

$$N = \left( \frac{K_{\max}}{K_{\min}} \right)^{\gamma-1} \approx 10^8$$

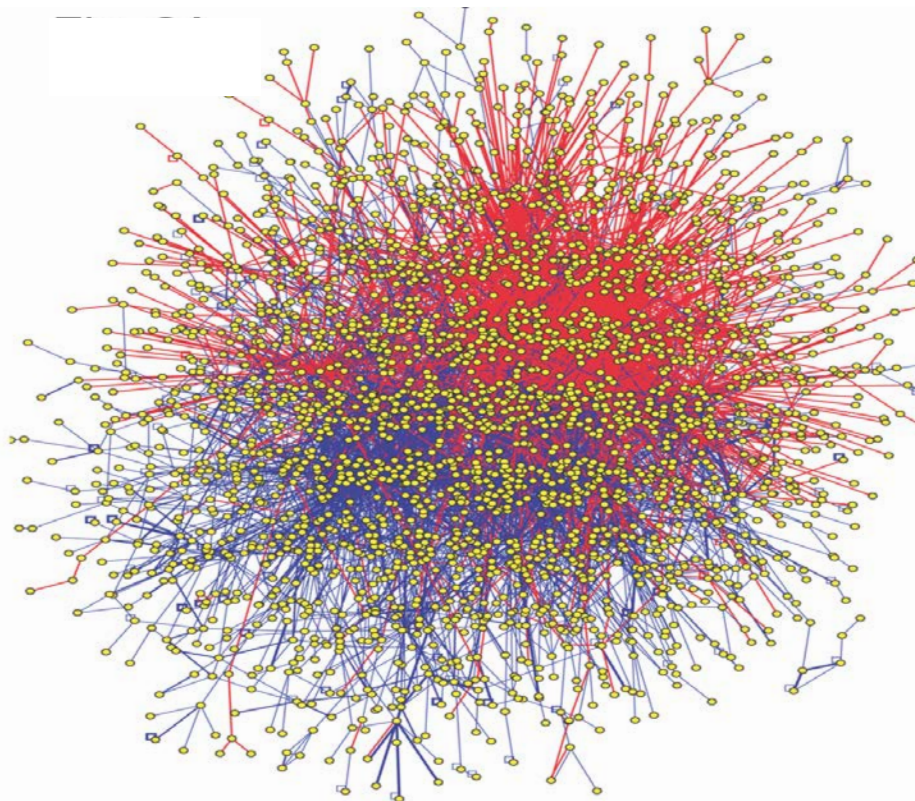
Mobile Call  
Network



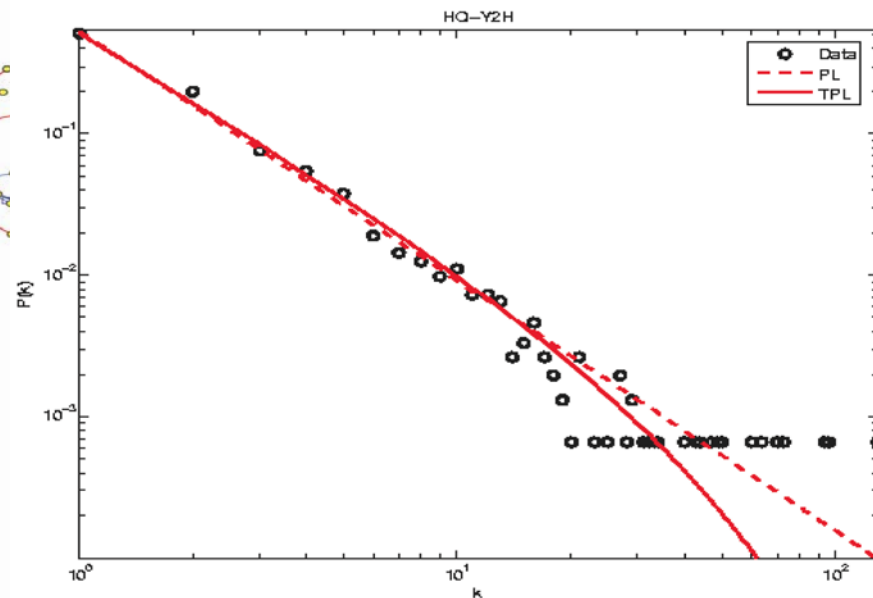
Characterizing the large-scale structure and the tie strengths of the mobile call graph. Vertex degrees are shown

# PLOTTING POWER LAWS

# HUMAN INTERACTION NETWORK

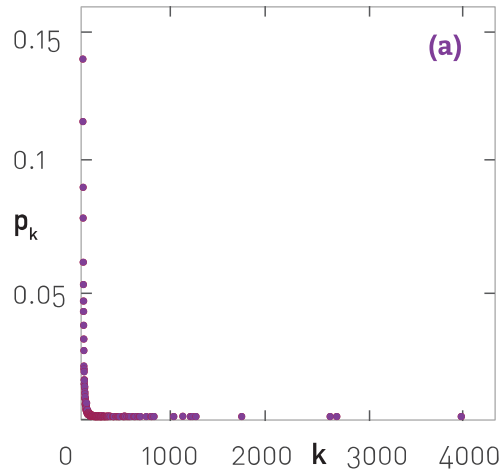


2,800 Y2H interactions  
4,100 binary LC interactions  
(HPRD, MINT, BIND, DIP, MIPS)

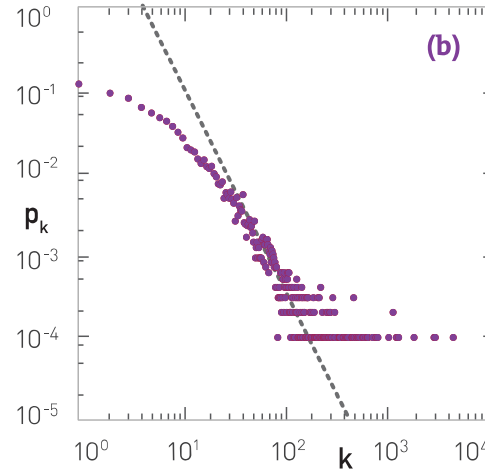


Rual et al. Nature 2005; Stelze et al. Cell 2005

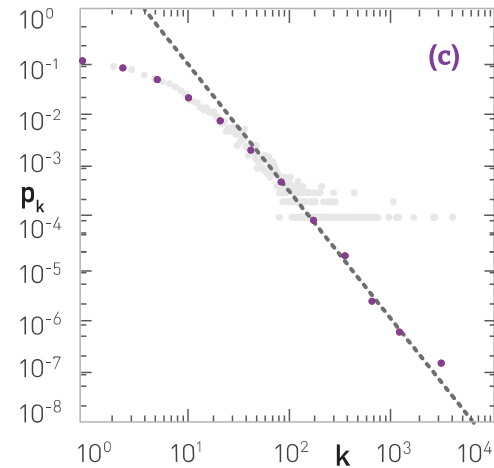
LINEAR SCALE



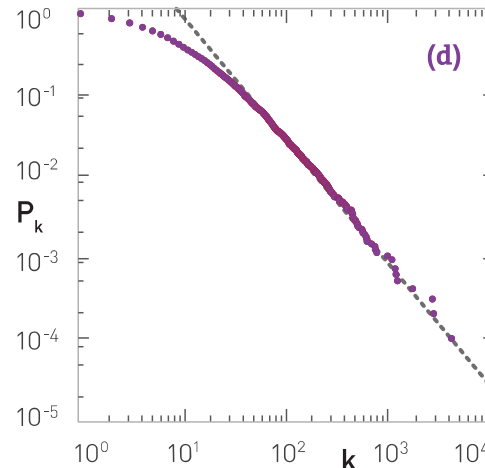
LINEAR BINNING



LOG-BINNING



CUMULATIVE



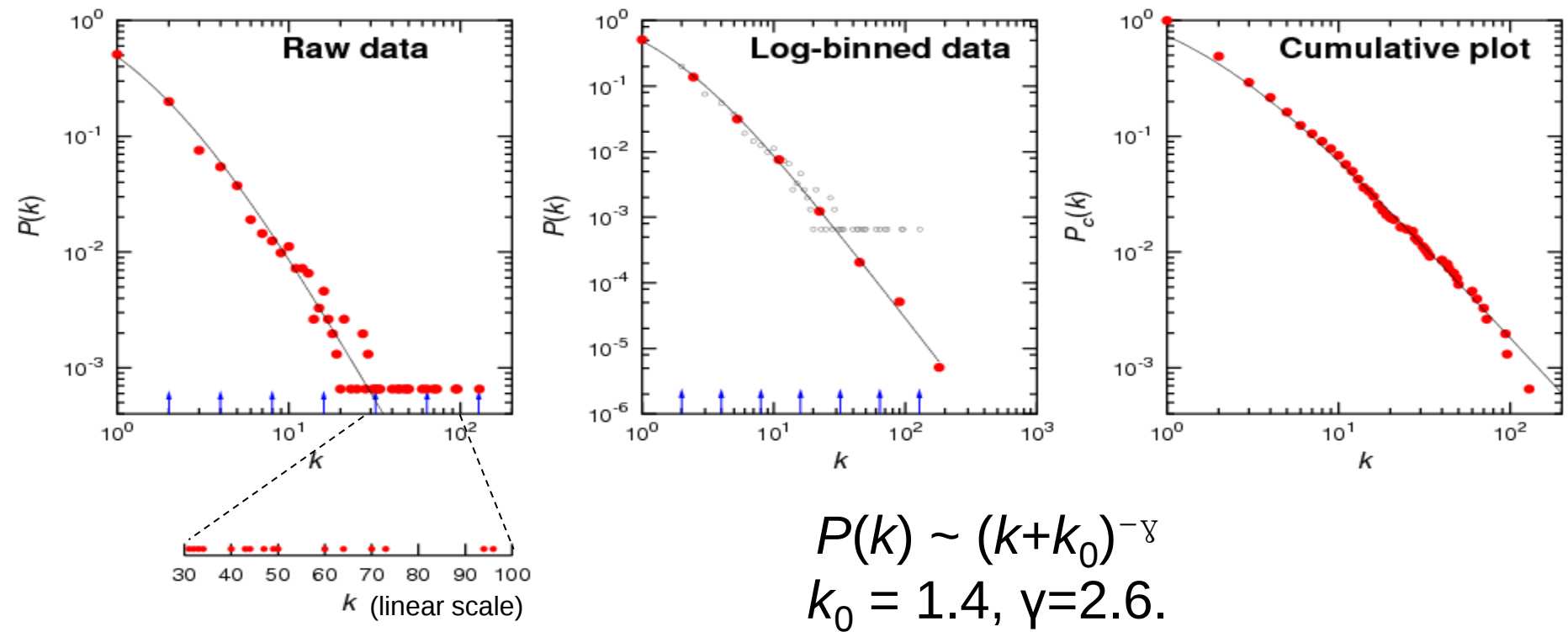
Use a Log-Log Plot

Avoid Linear Binning

Use Logarithmic Binning

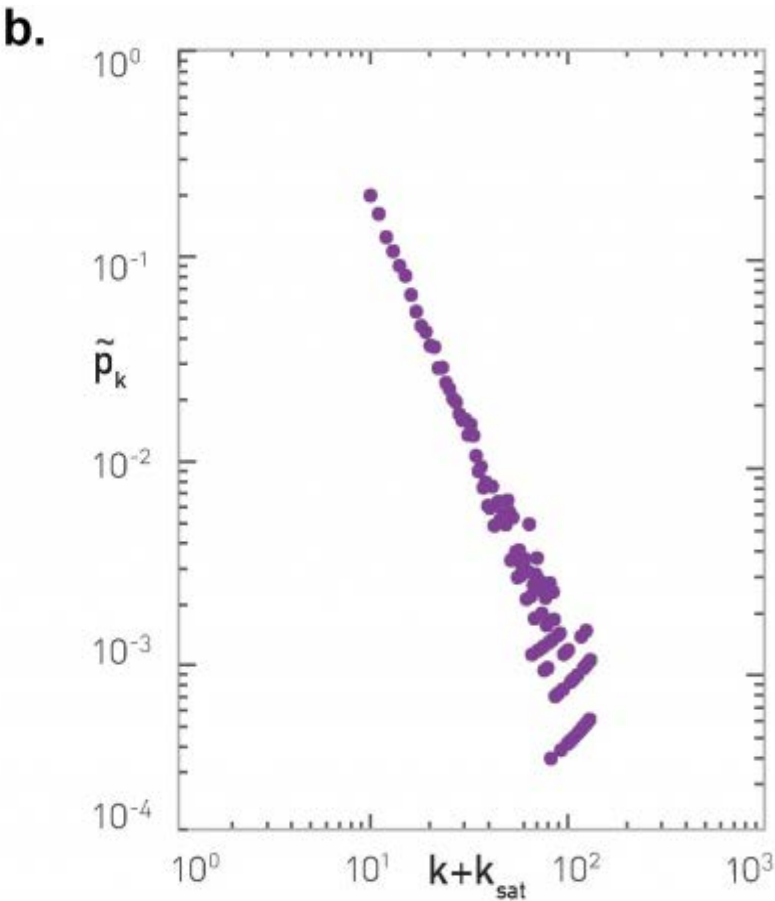
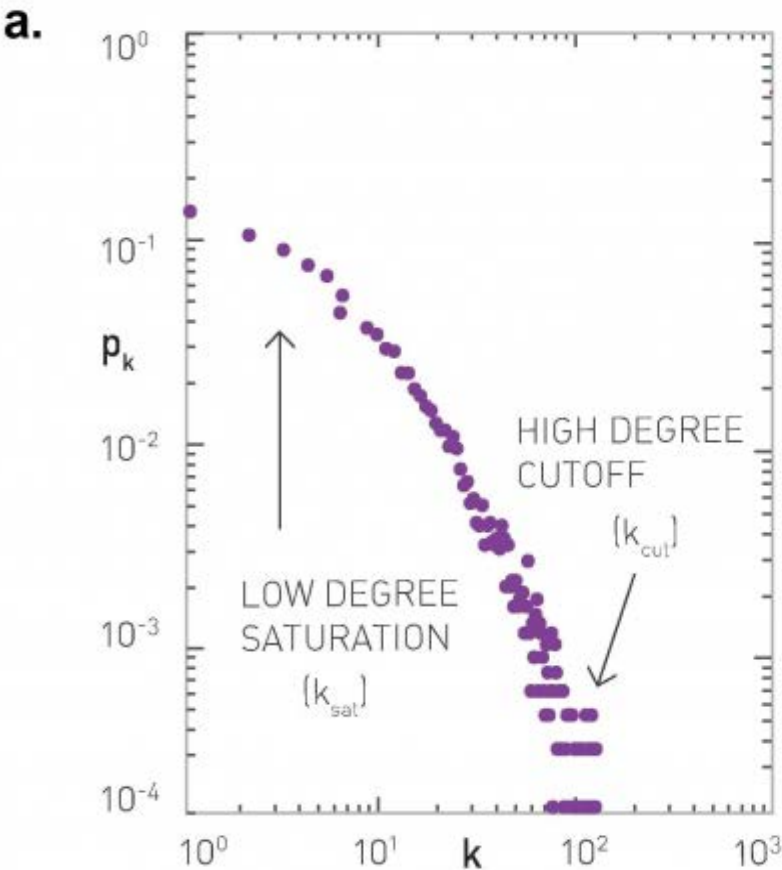
Use Cumulative Distribution

# HUMAN INTERACTION DATA BY RUAL ET AL.



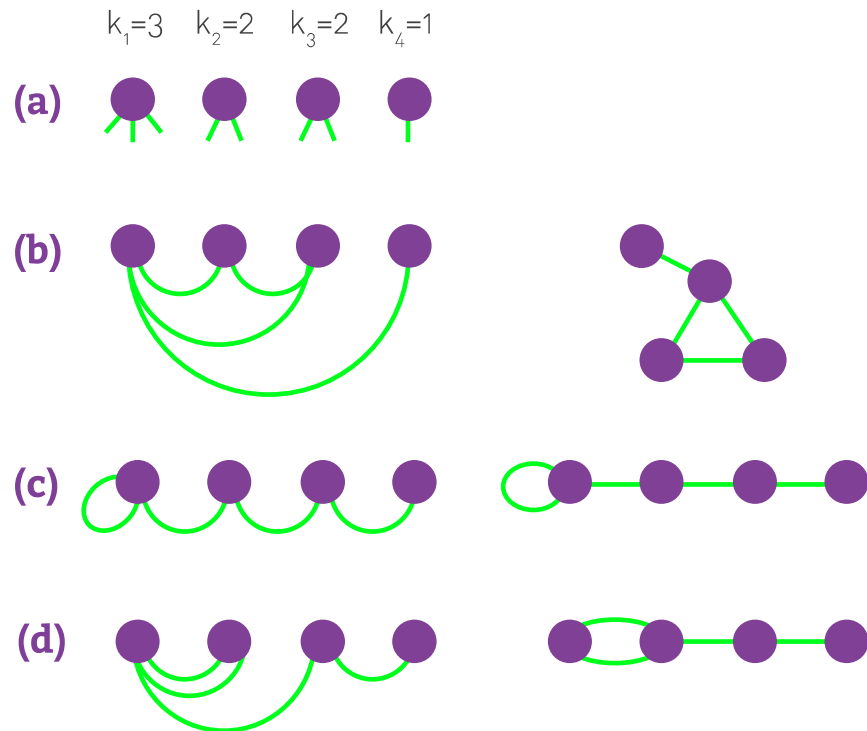
# COMMON MISCONCEPTIONS

$$p_k = a(k + k_{sat})^{-\gamma} \exp\left(-\frac{k}{k_{CW}}\right) \quad \bar{p}_k = p_k \exp\left(\frac{k}{k_{CW}}\right) \quad \tilde{k} = k + k_{sat} \quad p \sim \tilde{k}^{-\gamma}$$



# Generating networks with a pre-defined $p_k$

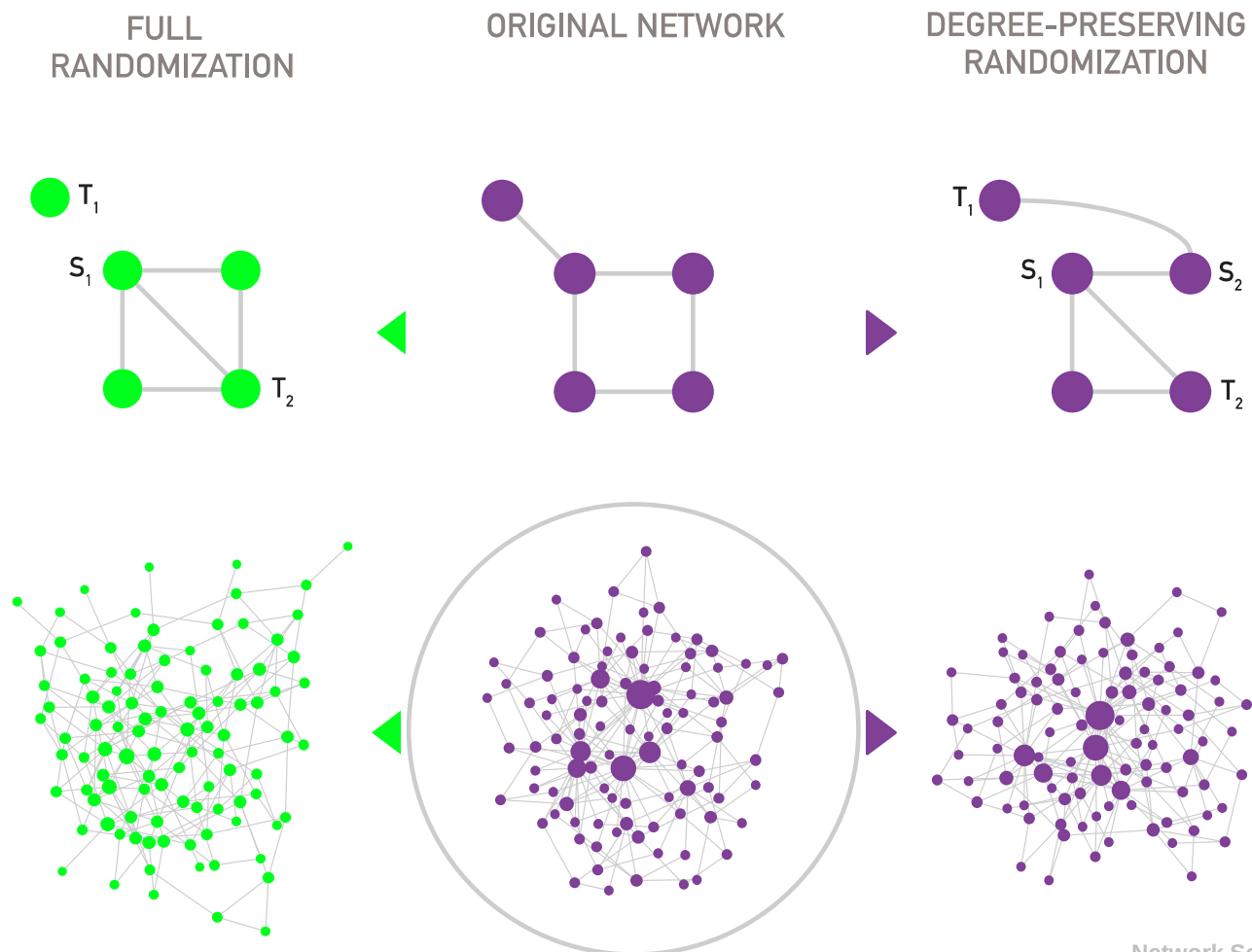
# Configuration model



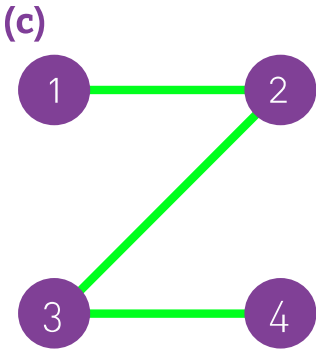
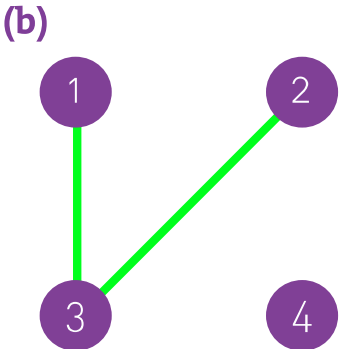
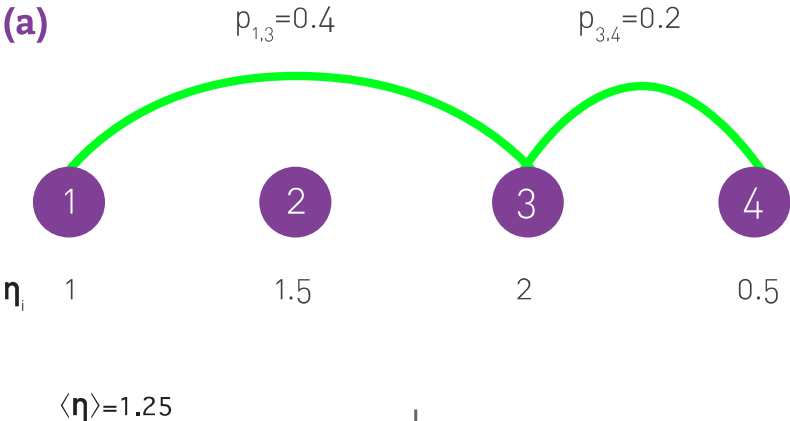
(1) **Degree sequence:** Assign a degree to each node, represented as stubs or half-links. The degree sequence is either generated analytically from a preselected distribution (Box 4.5), or it is extracted from the adjacency matrix of a real network. We must start from an even number of stubs, otherwise we will be left with unpaired stubs. (2) **Network assembly:** Randomly select a stub pair and connect them. Then randomly choose another pair from the remaining stubs and connect them. This procedure is repeated until all stubs are paired up. Depending on the order in which the stubs were chosen, we obtain different networks. Some networks include cycles (2a), others self-edges (2b) or multi-edges (2c). Yet, the expected number of self- and multi-edges goes to zero in the limit.

$$p_{ij} = \frac{k_i k_j}{2L - 1}$$

# Degree Preserving randomization



# Hidden parameter model



$$p(\eta_i, \eta_j) = \frac{\eta_i \eta_j}{\langle \eta \rangle N}$$

$$p_k = \int \frac{e^{-\eta} \eta^k}{k!} p(\eta) d\eta.$$

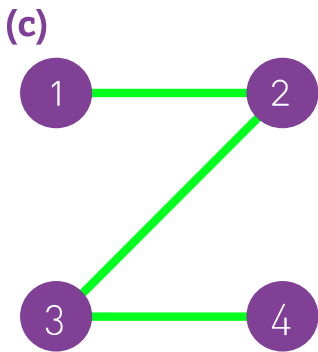
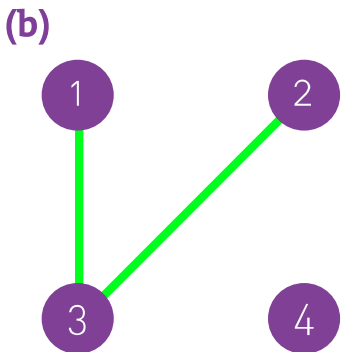
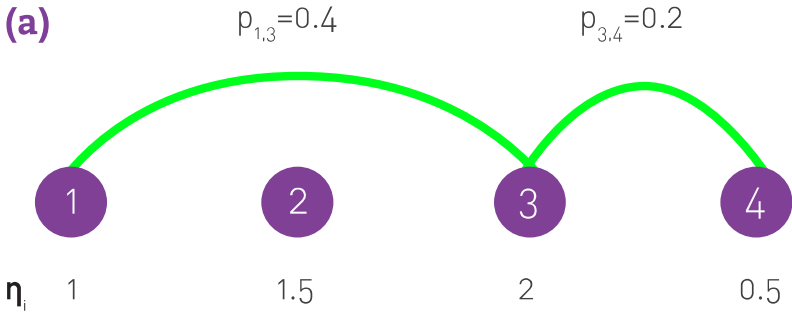
$$\{\eta_1, \eta_2, \dots, \eta_N\}.$$

$$p_k = \frac{1}{N} \sum_j \frac{e^{-\eta_j} \eta_j^k}{k!}.$$

$$\eta_j = \frac{c}{j^\alpha}, i = 1, \dots, N$$

$$p_k \sim k^{-(1+\frac{1}{\alpha})}$$

# Hidden parameter model



Start with  $N$  isolated nodes and assign to each node a “hidden parameter”  $\eta$ , which can be randomly selected from a  $p(\eta)$  distribution. We next connect each node pair with probability

$$p(\eta_i, \eta_j) = \frac{\eta_i \eta_j}{\langle \eta \rangle N}$$

For example, the figure shows the probability to connect nodes (1,3) and (3,4). After connecting the nodes, we end up with

the networks shown in (b) or (c), representing two independent realizations generated by the same hidden parameter sequence (a). The expected number of links in the obtained network is

$$L = \frac{1}{2} \sum_N \frac{\eta_i \eta_j}{\langle \eta \rangle N} = \frac{1}{2} \langle \eta \rangle N$$

$$p_k = \int \frac{e^{-\eta} \eta^k}{k!} p(\eta) d\eta.$$

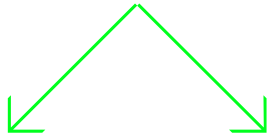
$$\{\eta_1, \eta_2, \dots, \eta_N\}, \quad p_k = \frac{1}{N} \sum_j \frac{e^{-\eta_j} \eta_j^k}{k!}.$$

$$\eta_j = \frac{c}{j^\alpha}, i = 1, \dots, N$$

$$p_k \sim k^{-(1+\frac{1}{\alpha})}$$

NETWORK OR DEGREE SEQUENCE

$$k_1, k_2, \dots, k_N$$



**FORBID**  
MULTI-LINKS

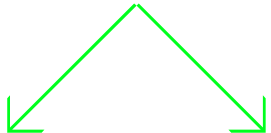
**ALLOW**  
MULTI-LINKS



**DEGREE-PRESERVING  
RANDOMIZATION**

DEGREE DISTRIBUTION

$$p_k$$

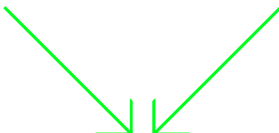


**ALLOW**  
MULTI-LINKS

**FORBID**  
MULTI-LINKS



**HIDDEN PARAMETER  
MODEL**



**CONFIGURATION  
MODEL**

# summary

## DEGREE DISTRIBUTION

Discrete form:

$$p_k = \frac{k^{-\gamma}}{\zeta(\gamma)}.$$

Continuous form:

$$p(k) = (\gamma - 1) k_{min}^{\gamma-1} k^{-\gamma}.$$

## SIZE OF THE LARGEST HUB

$$k_{max} \sim k_{min} N^{\frac{1}{\gamma-1}}.$$

MOMENTS OF  $p_k$  for  $N \rightarrow \infty$

$2 < \gamma < 3$ :  $\langle k \rangle$  finite,  $\langle k^2 \rangle$  diverges.

$\gamma > 3$ :  $\langle k \rangle$  and  $\langle k^2 \rangle$  finite.

## DISTANCES

$$\langle d \rangle \sim \begin{cases} \text{const.} & \gamma=2, \\ \frac{\ln \ln N}{\ln(\gamma-1)} & 2 < \gamma < 3, \\ \frac{\ln N}{\ln \ln N} & \gamma=3, \\ \ln N & \gamma > 3. \end{cases}$$

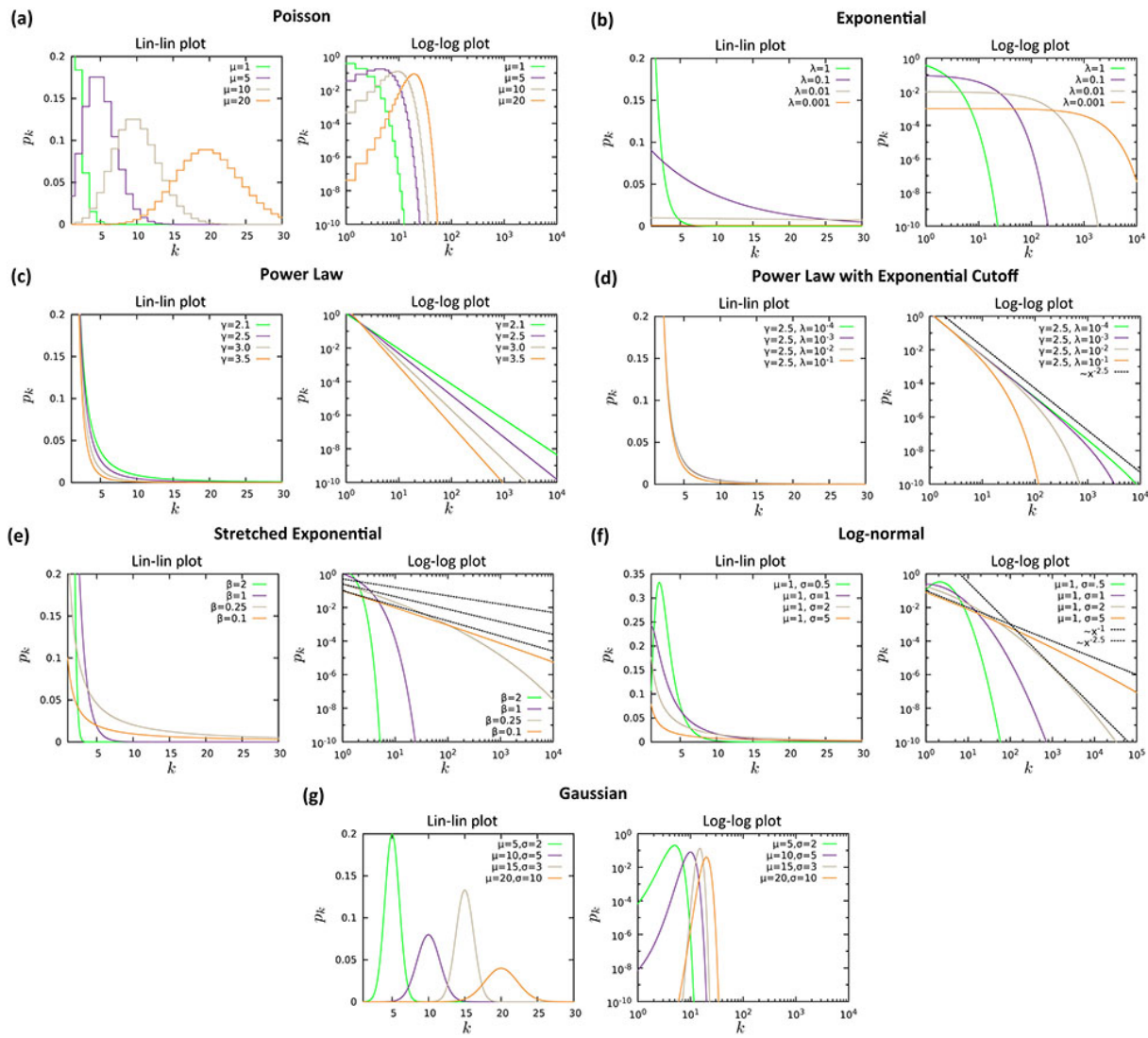
## Bounded Networks

We call a network *bounded* if its degree distribution decrease exponentially or faster for high  $k$ . As a consequence  $\langle k^2 \rangle$  is smaller than  $\langle k \rangle$ , implying that we lack significant degree variations. Examples of  $p_k$  in this class include the Poisson, Gaussian, or the simple exponential distribution (Table 4.2). The Erdős-Rényi and the Watts-Strogatz networks are the best known network models belonging to this class. Bounded networks lack outliers, consequently most nodes have comparable degrees. Real networks in this class include highway networks and the power grid.

## Unbounded Networks

We call a network *unbounded* if its degree distribution has a fat tail in the high- $k$  region. As a consequence  $\langle k^2 \rangle$  is much larger than  $\langle k \rangle$ , resulting in considerable degree variations. Scale-free networks with a power-law degree distribution (4.1) offer the best known example of networks belonging to this class. Outliers, or exceptionally high-degree nodes, are not only allowed but are expected in these networks. Networks in this class include the WWW, the Internet, the protein interaction networks, and most social and online networks.

# Random Distributions Important for Network Science

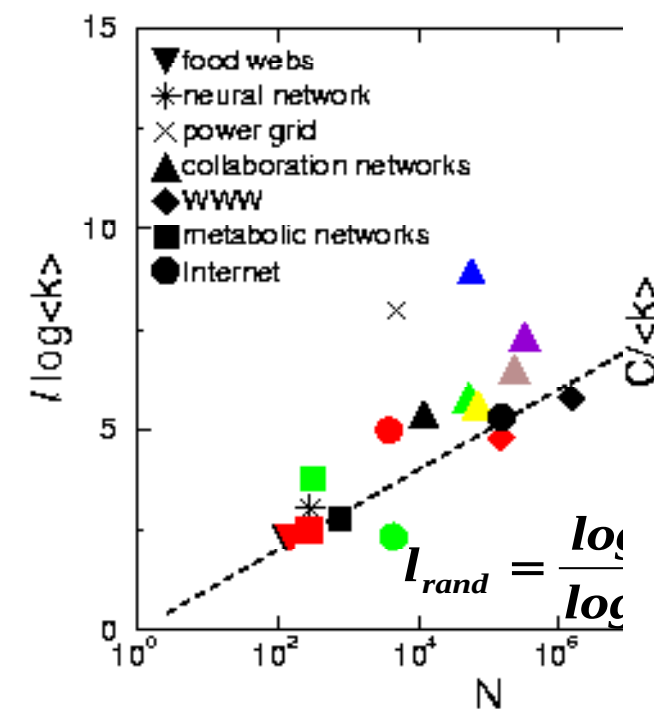


# Introduction to BA Model

Hubs represent the most striking difference between a random and a scale-free network. Their emergence in many real systems raises several fundamental questions:

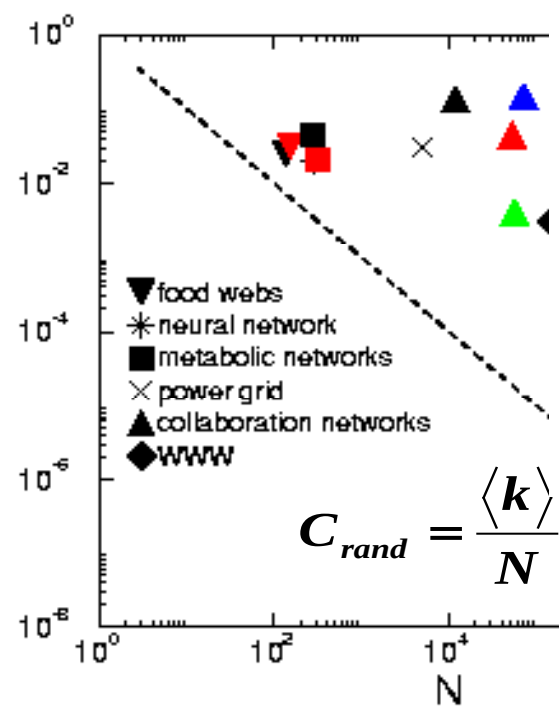
- Why does the random network model of Erdős and Rényi fail to reproduce the hubs and the power laws observed in many real networks?
- Why do so different systems as the WWW or the cell converge to a similar scale-free architecture?

# Empirical findings for real networks



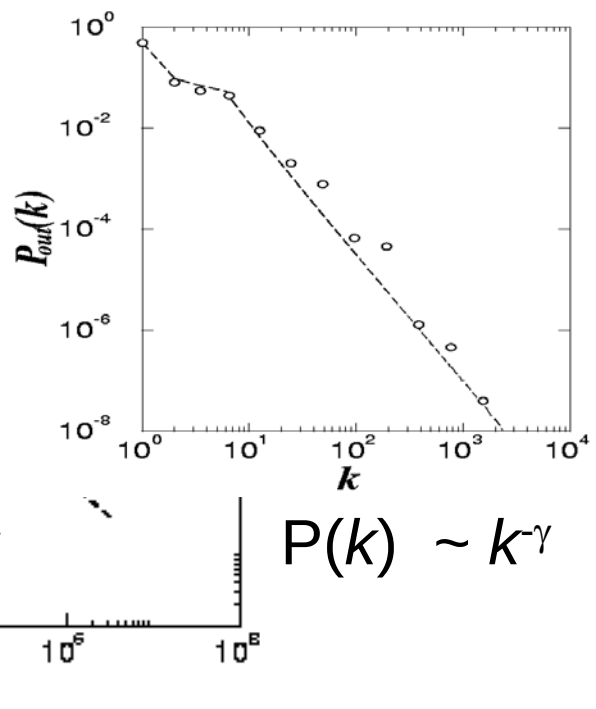
## Small World:

distances scale logarithmically with the network size



## Clustered:

clustering coefficient does not depend on network size.

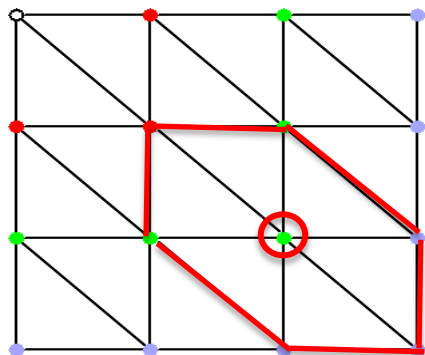


## Scale-free:

The degrees follow a power-laws distribution.

## BENCHMARK 1: Regular Lattices

## Two-dimensional lattice:



Average path-length:

$$l \approx L \approx N^{1/2}$$

### Degree distribution:

$P(k)=\delta(k-6)$  excluding corners  
and boundaries

Clustering coefficient:

$$C = \frac{2}{5}$$

6 neighbors, each with 2 edges = 12/30

### ***D-dimensional lattice:***

The average path-length varies as

$$l \approx N^{1/D}$$



## Constant degree

$$P(k) = \delta(k - k_d)$$



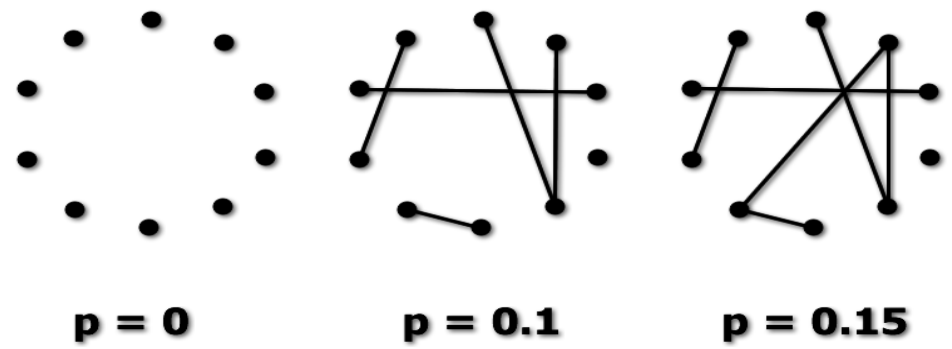
## Constant clustering coefficient

$$C=C_d$$



# BENCHMARK 2: Random Network Model

Erdős-Rényi Model- Publ. Math. Debrecen 6, 290 (1959)



- fixed node number *N*
- connecting pairs of nodes with probability *p*

Degree distribution:

$$P_{rand}(k) \cong C_{N-1}^k p^k (1-p)^{N-1-k}$$



Path length:

$$l_{rand} \approx \frac{\log N}{\log \langle k \rangle}$$



Clustering coefficient:

$$C_{rand} = p = \frac{\langle k \rangle}{N}$$



# BENCHMARK 3: Small World Model

## Watts-Strogatz algorithm – Nature 2008



- For fixed node number  $N$ , first connect them into even number,  $k$ , degree ring in which  $k/2$  nearest neighbors on each side of each node are connected to it
- Then, with probability  $p$  re-wire ring edges of each node to nodes not currently connected to and different from it

Degree distribution:

Exponential



Path length:

$$l_{rand} \approx \frac{\log N}{\log \langle k \rangle}$$



Clustering coefficient:

$$C(\beta) = C(0)(1 - \beta)$$



# Growth and preferential attachment

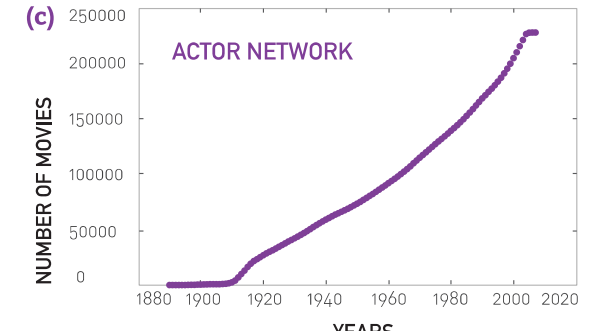
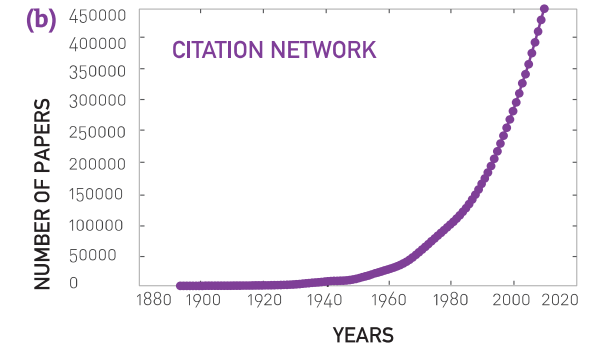
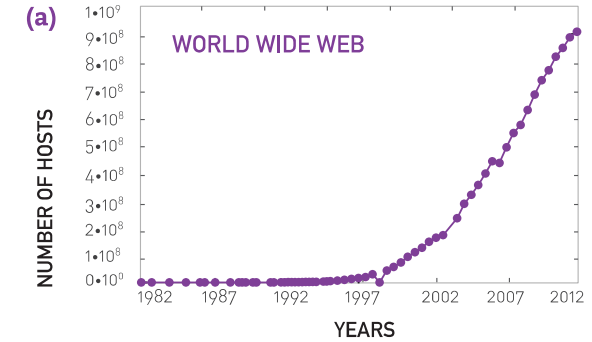
ER, WS models: the number of nodes,  $N$ , is fixed (static models)

**Real networks continuously expand by the addition of new nodes**

## ER model:

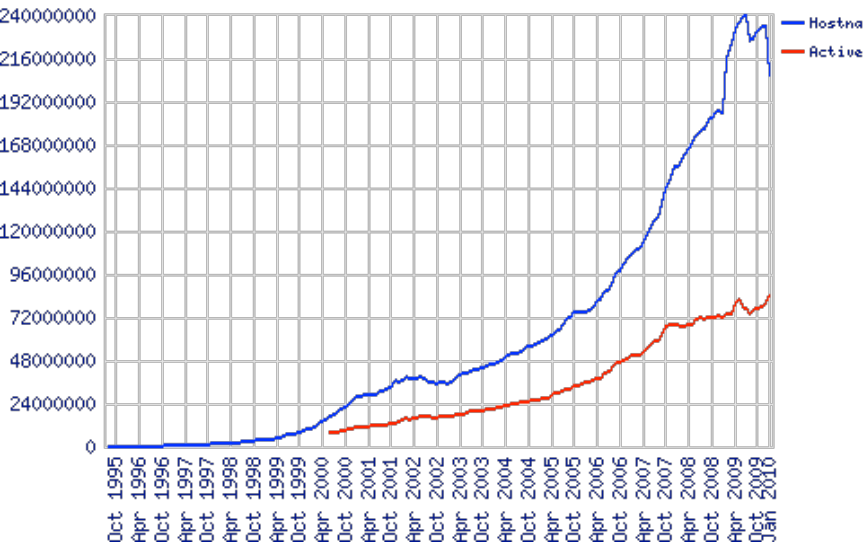
the number of nodes,  $N$ , is fixed (static models)

networks expand through the addition  
of new nodes

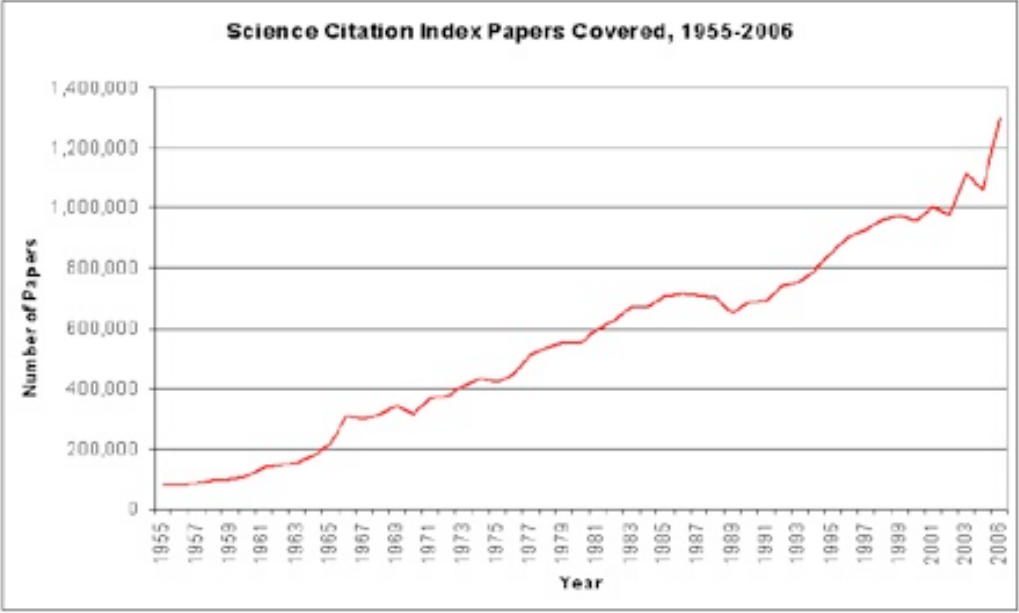


# Growth (www/Pubs)

## WWW



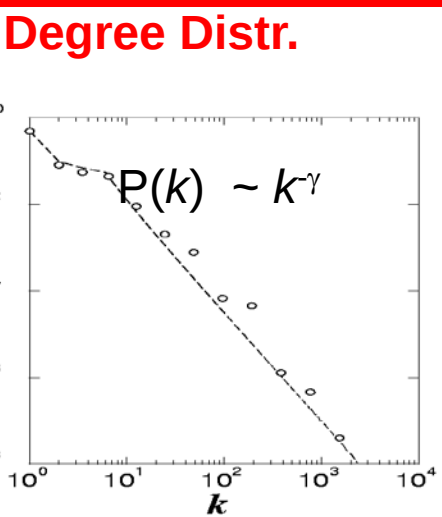
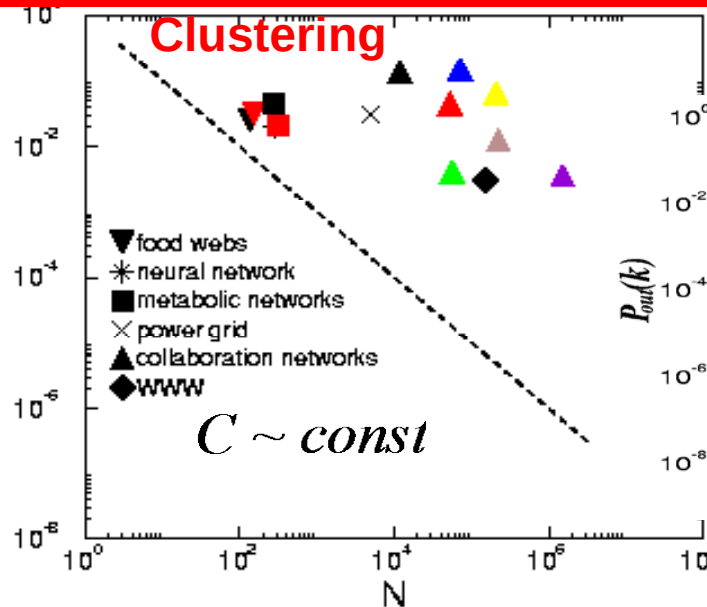
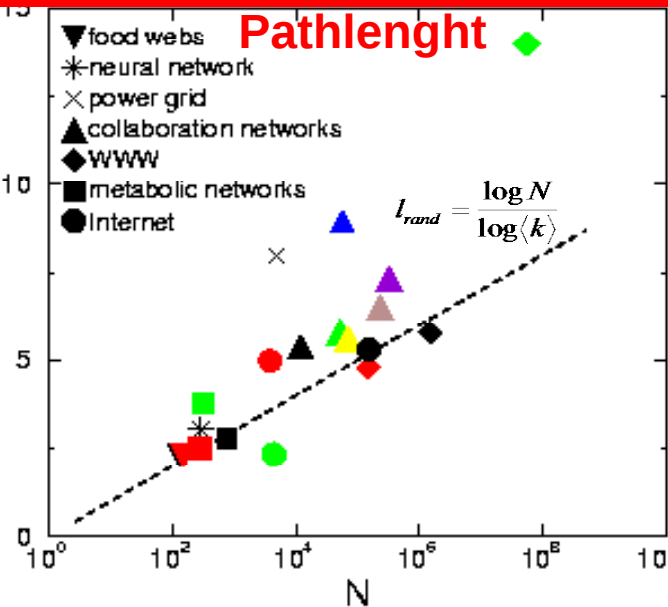
## Scientific Publications



<http://website101.com/define-ecommerce-web-terms-definitions/>

[http://www.kk.org/thetechnium/archives/2008/10/the\\_expansion\\_o.php](http://www.kk.org/thetechnium/archives/2008/10/the_expansion_o.php)

# EMPIRICAL DATA FOR REAL NETWORKS



|                 |                                                          |   |                                              |   |                                                                |   |
|-----------------|----------------------------------------------------------|---|----------------------------------------------|---|----------------------------------------------------------------|---|
| Regular network | $l \approx N^{1/D}$                                      | 👎 | $C \sim const$                               | 👍 | $P(k) = \delta(k - k_d)$                                       | 👎 |
| Erdos-Renyi     | $l_{rand} \approx \frac{\log N}{\log \langle k \rangle}$ | 👍 | $C_{rand} = p = \frac{\langle k \rangle}{N}$ | 👎 | $P(k) = e^{-\langle k \rangle} \frac{\langle k \rangle^k}{k!}$ | 👎 |
| Watts-Strogatz  | $l_{rand} \approx \frac{\log N}{\log \langle k \rangle}$ | 👍 | $C \sim const$                               | 👍 | Exponential                                                    | 👎 |

ER model: links are added randomly to the network

**New nodes prefer to connect to the more connected nodes**

# The Barabási-Albert model

## Section 2: Growth and Preferential Attachment

The random network model differs from real networks in two important characteristics:

**Growth:** While the random network model assumes that the number of nodes is fixed (time invariant), real networks are the result of a growth process that continuously increases.

**Preferential Attachment:** While nodes in random networks randomly choose their interaction partner, in real networks new nodes prefer to link to the more connected nodes.

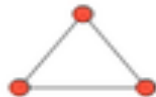
# Origin of SF networks: Growth and preferential attachment

(1) Networks continuously expand by the addition of new nodes

WWW : addition of new documents

(2) New nodes prefer to link to highly connected nodes.

WWW : linking to well known sites

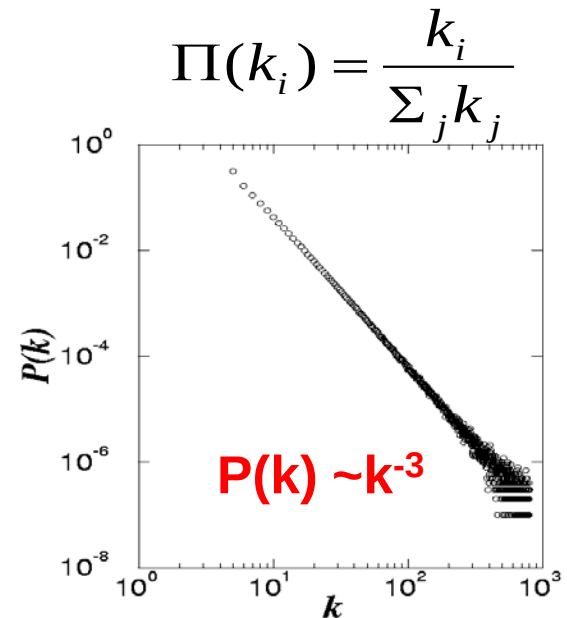


**GROWTH:**

add a new node with  $m$  links

**PREFERENTIAL ATTACHMENT:**

the probability that a node connects to a node with  $k$  links is proportional to  $k$ .





György Pólya  
**PÓLYA PROCESS**  
MATHEMATICIAN



George Kinsley Zipf  
**WEALTH DISTRIBUTION**  
ECONOMIST



Herbert Alexander Simon  
**MASTER EQUATION**  
POLITICAL SCIENTIST



Robert Merton  
**MATTHEW EFFECT**  
SOCIOLOGIST



Albert-László Barabási & Réka Albert  
**PREFERENTIAL ATTACHMENT**  
NETWORK SCIENTISTS



George Udny Yule  
**YULE PROCESS**  
STATISTICIAN



Robert Gibrat  
**PROPORTIONAL GROWTH**  
ECONOMIST



Derek de Solla Price  
**CUMULATIVE ADVANTAGE**  
PHYSICIST

## MILESTONES

PUBLICATION  
DATE

1923 1925 1931 1935 1941 1945 1950 1955 1960 1968 1970 1976 1980 1985 1990 1995 1999 2000 2005 2010

**György Pólya (1887-1985)**  
Preferential attachment made its first appearance in 1923 in the celebrated urn model of the Hungarian mathematician György Pólya [2]. Hence, in mathematics preferential attachment is often called a **Pólya process**.

**George Udny Yule (1871-1951)**  
used preferential attachment to explain the power-law distribution of the number of species per genus of flowering plants [3]. Hence, in statistics preferential attachment is often called a **Yule process**.

**Robert Gibrat (1904-1980)**  
proposed that the size and the growth rate of a firm are independent. Hence, larger firms grow faster [4]. Called **proportional growth**, this is a form of preferential attachment.

**George Kinsley Zipf (1902-1950)**  
used preferential attachment to explain the fat tailed distribution of wealth in the society [5].

**Herbert Alexander Simon (1916-2001)**  
used preferential attachment to explain the fat-tailed nature of the distributions describing city sizes, word frequencies, or the number of papers published by scientists [6].

**Derek de Solla Price (1922-1983)**  
used preferential attachment to explain the citation statistics of scientific publications, referring to it as **cumulative advantage** [7].

**Robert Merton (1910-2003)**  
In sociology preferential attachment is often called the **Matthew effect**, named by Merton [8] after a passage in the Gospel of Matthew.

**Barabási (1967) & Albert (1972)**  
introduce the term **preferential attachment** in the context of networks [1] to explain the origin of their power-law degree distribution.

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