

Problem Set #8

due November 26

Problem 1. Rewrite `inferTypes` from PS7 to use a state transformer instead of a “stateful” variable. Signature changes from

```
inferTypes :: TEnv -> Integer -> Exp -> (Subst, Type, Integer)
to
```

```
inferTypes :: TEnv -> Exp -> S.State Integer (Subst, Type)
```

Transfer your previous homework to file `Infer.hs`. The only change should be `inferTypes` and callers of `inferTypes`.

Problem 2. Next, write a parser for the lambda calculus following this grammar:

$$\begin{aligned} \textit{lexp} & ::= \textit{lapp} \mid \textit{labs} \mid \textit{lterm} \\ \textit{lapp} & ::= \textit{lterm} \textit{lterm} \textit{ttail} \\ \textit{ttail} & ::= \textit{lterm} \textit{ttail} \mid \epsilon \\ \textit{labs} & ::= \lambda \textit{identifier} \rightarrow \textit{lexp} \\ \textit{lterm} & ::= (\textit{lexp}) \mid \textit{identifier} \mid \textit{integer} \mid \text{T} \mid \text{F} \end{aligned}$$

The function parses a lambda expression into the `Exp` datatype from PS7 and `Data.hs` is included in the starter code (again).

```
lexp :: P.Parser Exp
lexp = undefined
```

```
> P.runParser (start lexp) "\\f -> \\x -> f (f x)$"
[(ELambda "f" (ELambda "x" (EApp (EVar "f") (EApp (EVar "f") (EVar "x"))))),(")]
```

We can create expressions (by parsing them from text) and infer types a lot easier:

```
> twice = fst $ head $ P.runParser (start lexp) "\\f -> \\x -> f (f x)$"
> I.canonicalize twice -- from PS7, now in Infer.hs
"(t1 -> t1) -> t1 -> t1"
```

Problem 3. Finally, extend the parser with handling of ambiguous grammars in the sense that the parser now finds *all* possible parse trees for a string. I will test with the following grammar that generates all strings of equal number of `a`'s and `b`'s:

$$s ::= a s b s \mid b s a s \mid \epsilon$$

The output of a parse of `s` should be the production sequence the parser finds where productions are numbered from left to right: `s ::= a s b s` is 1, `s ::= b s a s` is 2, and `s ::= ϵ` is 3. Naturally, the autograder expects that your parser respects this order when trying alternatives.

```
s :: P.Parser [Int]
s = undefined
```

```
> P.runParser (start s) "abab$"
([(1,2,3,3,3),(")],([1,3,1,3,3],("))]
```

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E.g., $[1, 2, 3, 3, 3]$ is leftmost derivation $s \xrightarrow{1} asbs \xrightarrow{2} absasbs \xrightarrow{3} abasbs \xrightarrow{3} ababs \xrightarrow{3} abab$.

Note: Download files `Data.hs`, `State.hs`, `Lexer.hs`, `Parser.hs` and `Infer.hs`. Minimal starter code is in `Ps8.hs`.