

Infinity



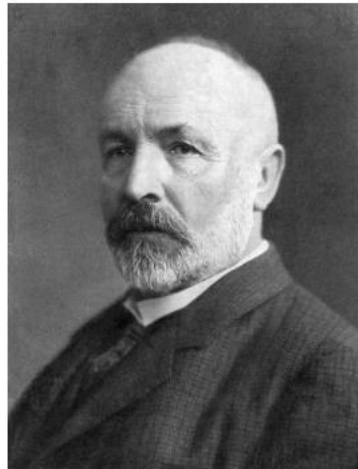
- Malik Magdon-Ismael. Discrete Mathematics and Computing.
 - Chapter 22



- Precise statements, proofs and logic.
- **INDUCTION.**
- Recursively defined structures and Induction. (Data structures; PL)
- Sums and asymptotics. (Algorithm analysis)
- Number theory. (Cryptography; probability; fun)
- Graphs. (Relationships/conflicts; resource allocation; routing; scheduling, . . .)
- Counting. (Enumeration and brute force algorithms)
- Probability. (Real world algorithms involve randomness/uncertainty)
 - Inputs arrive in a random order;
 - Randomized algorithms (primality testing, machine learning, routing, conflict resolution . . .)
 - Expected value is a summary of what happens. Variance tells you how good the summary is.

Today: Infinity

- Comparing “sizes” of sets: countable.
 - Rationals are countable.
- Uncountable
 - Infinite binary strings.
- What does Infinity have to do with computing?

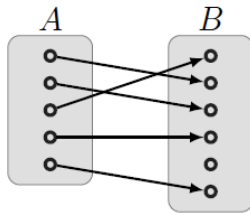


Georg Cantor

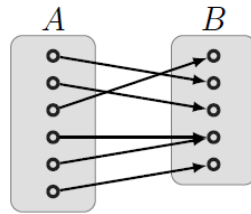


“Size” of a Set: Cardinality

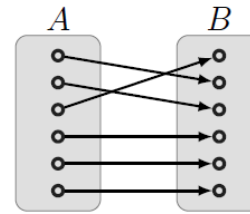
- There’s a reason why small kids use fingers to count
 - They map their intuitive knowledge of 2-fingers to 2 of another object
- You have an *equal* number of fingers on each hand
 - You can map your left hand’s fingers to your right hand’s fingers
- Recall some types of such maps



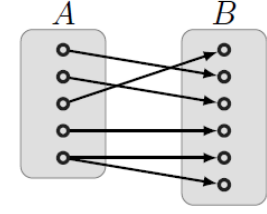
1-1, but **not** onto.
(injection, $A \xrightarrow{INJ} B$)
 $|A| \leq |B|$



onto; **not** 1-1
(surjection, $A \xrightarrow{SUR} B$)
 $|A| \geq |B|$

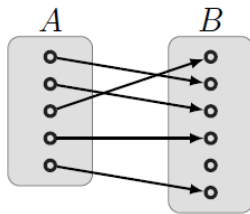


onto **and** 1-1
(bijection, $A \xrightarrow{BIJ} B$)
 $|A| = |B|$

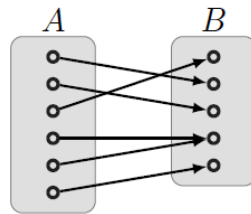


not a function

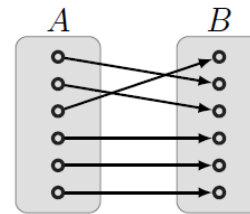
“Size” of a Set: Cardinality



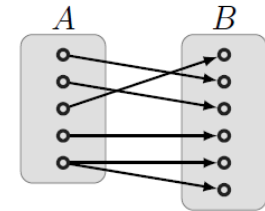
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not a function

- *Cardinality* $|A|$ (“size”), read “cardinality of A ,” is the number of elements for finite sets
- In general, we can define the following relations between sets:

$|A| \leq |B|$ iff there is an injection (1-to-1) from A to B , i.e., $f: A \xrightarrow{INJ} B$

$|A| > |B|$ iff there is no injection (1-to-1) from A to B

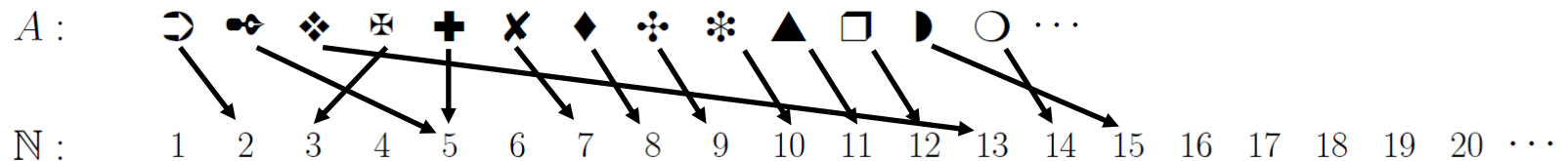
$|A| \geq |B|$ iff there is a surjection (onto) from A to B , i.e., $f: A \xrightarrow{SUR} B$

$|A| = |B|$ iff there is a bijection (1-1 and onto) from A to B , i.e., $f: A \xrightarrow{BIJ} B$

$|A| \leq |B|$ and $|B| \leq |A| \rightarrow |A| = |B|$ [Cantor-Bernstein Theorem]

A Countable Set's Cardinality is at most $|\mathbb{N}|$

- Suppose we have a finite set $A = \{a_1, \dots, a_n\}$
 - Cardinality is $|A| = n$ if and only if there is a bijection from A to $\{1, \dots, n\}$
 - Can you come up with such a function?
$$f(a_i) = i$$
- For infinite sets: the set A is countable if $|A| \leq |\mathbb{N}|$.
 - Intuitively, A is “smaller than” \mathbb{N}
 - Sometimes we say A is at most countable to include both finite and infinite sets that are “smaller than” \mathbb{N}
- To show that A is countable you must find a 1-to-1 mapping from A to \mathbb{N}



- You cannot skip over any elements of A , but you might not use every element of \mathbb{N}

A Countable Set's Cardinality is at most $|\mathbb{N}|$, cont'd



- To prove that a function $f: A \mapsto \mathbb{N}$ is an injection:
 - Assume f is *not* an injection. (Proof by contradiction.)
 - This means there is a pair $x, y \in A$ for which $x \neq y$ and $f(x) = f(y)$
 - Use $f(x) = f(y)$ to prove that $x = y$, a contradiction. Hence, f is an injection

All Finite Sets are Countable

- Suppose $A = \{3,6,8\}$
 - To show $|A| \leq |\mathbb{N}|$, we give an injection from A to \mathbb{N}
 $3 \mapsto 1, 6 \mapsto 23134, 8 \mapsto 8$
- For an arbitrary set $A = \{a_1, \dots, a_n\}$
 $a_1 \mapsto 1, a_2 \mapsto 2, \dots, a_n \mapsto n$

Non-negative integers $\mathbb{N}_0 = \{0, 1, 2, \dots\}$ are countable

- How can this be??
 - I know for a fact that \mathbb{N}_0 contains every element in \mathbb{N} , plus an extra 0
 - It's clearly bigger!!
- Well, if they were finite sets, I would agree. But let's recall the definition.
- To prove that $|\mathbb{N}_0| \leq |\mathbb{N}|$, we need an injection $f: \mathbb{N}_0 \xrightarrow{INJ} \mathbb{N}$
 - Ideas?
 - Let's try $f(x) = x + 1, \forall x \in \mathbb{N}_0$
- *Proof.*
 - Assume f is not an injection. So, there are $x \neq y$ in \mathbb{N}_0 with $f(x) = f(y)$:
$$x + 1 = f(x) = f(y) = y + 1$$
 - But that means $x + 1 = y + 1$, i.e., $x = y$. Contradiction.
- Also, we know that $|\mathbb{N}| \leq |\mathbb{N}_0|$ since $\mathbb{N} \subseteq \mathbb{N}_0$
- By the Cantor-Bernstein Theorem, $\mathbb{N} = \mathbb{N}_0$

$$\begin{array}{cccccccccccc} \mathbb{N}_0 : & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & \dots \\ & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \dots \\ \mathbb{N} : & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & \dots \end{array}$$

Positive Even Numbers and Integers and Countable

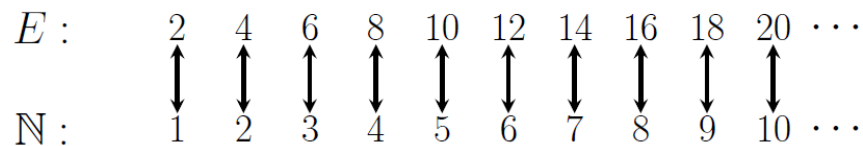


- Huh? The even numbers are exactly half of all natural numbers!!

$$E = \{2, 4, 6, \dots\}, \text{ so surely } |E| = \frac{1}{2}|\mathbb{N}|!!$$

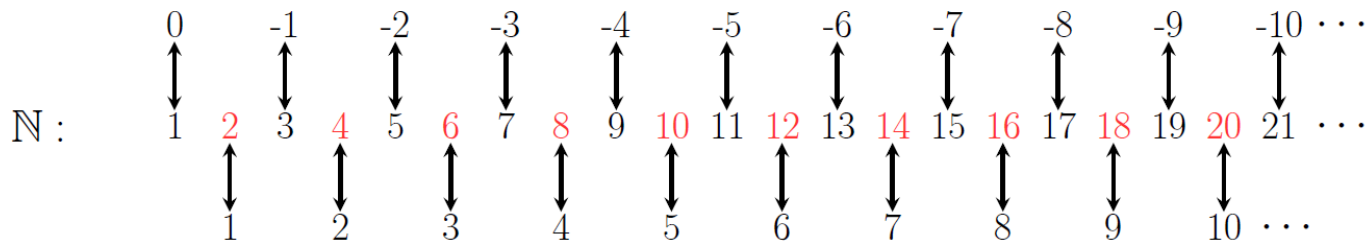
- Turns out, not quite. Can you see a bijection?

- The bijection $f(x) = \frac{1}{2}x$ proves $|E| = |\mathbb{N}|$



- OK, fine, but all integers?? It has to be the case that $|\mathbb{Z}| = 2|\mathbb{N}|!!$

- Is there a bijection? ☺
- Recall the integers are $\mathbb{Z} = \{0, \pm 1, \pm 2, \dots\}$



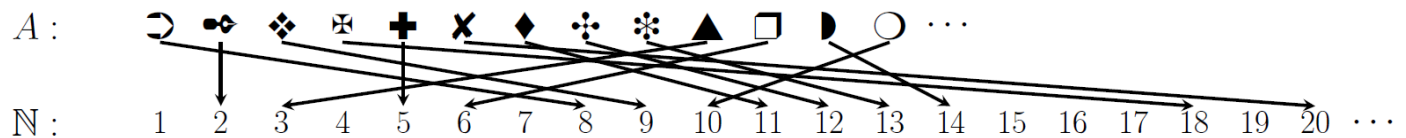
- **Exercise.** What is a mathematical formula for the bijection?

Every Countable Set Can Be “Listed”

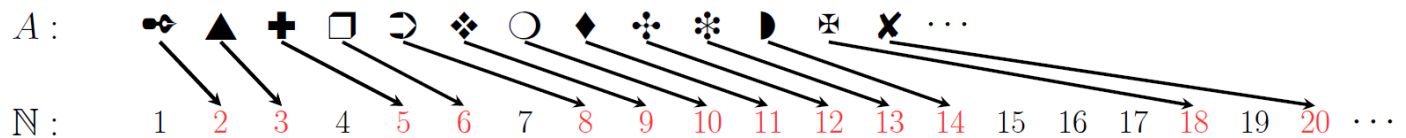
- What does it mean for a set to be “listed”?
 - You know the exact position of each element in the set
 - Regardless of whether the set is finite or infinite
- For example:
 - $\{3,6,8\}$ is a list (why?)
 - (because it’s a finite set)
 - $E = \{2,4,6, \dots\}$ is a list (why?)
 - element i is just $2i$
 - What about \mathbb{Z} ?
 - Suppose I represent it as $\{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$
 - Unclear what the indices are
 - Suppose I represent it as $\{0, -1, 1, -2, 2, \dots\}$
 - Element 1 is 0, o.w. it is $-i/2$ (if i is even), $(i - 1)/2$ (if i is odd)

Every Countable Set Can Be “Listed”, cont’d

- Suppose I give you the following mapping between sets A and \mathbb{N}



- How do I “list” A ?
- Order elements according to their assigned value



- In general, a set can be “listed” if
 - Different elements are assigned to different list-positions.
 - We can determine the list-position of *any* element in the set.

Union of Two Countable Sets is Countable

- Consider two countable sets, A and B
 - They are countable, so I can write $A = \{a_1, a_2, a_3, \dots\}$ and $B = \{b_1, b_2, b_3, \dots\}$

- Now, let's look at the union

$$A \cup B = \{a_1, a_2, a_3, \dots, b_1, b_2, b_3, \dots\}$$

- Hm, how do I show this is countable?
 - Can't use "... " twice
 - How do I reorder terms?
 - Need to know the position of each b_i
- Here's a better reordering:
$$A \cup B = \{a_1, b_1, a_2, b_2, a_3, b_3, \dots\}$$
 - Now I know the position of each element
 - List-position of a_i is $2i - 1$
 - List-position of b_i is $2i$
- **Exercise.** Get a list of \mathbb{Z} with $A = \{0, -1, -2, -3, \dots\}$ and $B = \{1, 2, 3, \dots\}$ using union.

Rationals are Countable: $|\mathbb{Q}| = |\mathbb{N}|$



- OK, this one is very surprising!
 - There are **infinitely** many rationals between every two integers!
 - The rationals are dense (there is a rational between any two rationals)!
 - Natural numbers are not!
 - Also, the set of rationals can be expressed as the product of integers and natural numbers:
 - $|\mathbb{Q}| = |\mathbb{N}| \times |\mathbb{Z}|$, so of course $|\mathbb{Q}| \gg |\mathbb{N}|??$
- Well, let's see...

Rationals are Countable: $|\mathbb{Q}| = |\mathbb{N}|$, cont'd

		\mathbb{Z}								
		0	1	-1	2	-2	3	-3	4	-4
\mathbb{N}	1	$\frac{0}{1}$	$\frac{1}{1}$	$\frac{-1}{1}$	$\frac{2}{1}$	$\frac{-2}{1}$	$\frac{3}{1}$	$\frac{-3}{1}$	$\frac{4}{1}$	$\frac{-4}{1}$
	2	$\frac{0}{2}$	$\frac{1}{2}$	$\frac{-1}{2}$	$\frac{2}{2}$	$\frac{-2}{2}$	$\frac{3}{2}$	$\frac{-3}{2}$	$\frac{4}{2}$	$\frac{-4}{2}$
	3	$\frac{0}{3}$	$\frac{1}{3}$	$\frac{-1}{3}$	$\frac{2}{3}$	$\frac{-2}{3}$	$\frac{3}{3}$	$\frac{-3}{3}$	$\frac{4}{3}$	$\frac{-4}{3}$
	4	$\frac{0}{4}$	$\frac{1}{4}$	$\frac{-1}{4}$	$\frac{2}{4}$	$\frac{-2}{4}$	$\frac{3}{4}$	$\frac{-3}{4}$	$\frac{4}{4}$	$\frac{-4}{4}$
	5	$\frac{0}{5}$	$\frac{1}{5}$	$\frac{-1}{5}$	$\frac{2}{5}$	$\frac{-2}{5}$	$\frac{3}{5}$	$\frac{-3}{5}$	$\frac{4}{5}$	$\frac{-4}{5}$

- How do I “list” all rationals?
 - I need a list that visits each rational *exactly* once!
 - I need to know each rational’s position *exactly*

$$\mathbb{Q} = \left\{ \frac{0}{1}, \frac{1}{1}, \frac{1}{2}, \frac{-1}{2}, \frac{2}{3}, \frac{-2}{3}, \dots \right\}$$

$$|\{\text{Rational Values}\}| \leq |\mathbb{Q}| \leq |\mathbb{N}|$$

- **Exercise.** What is a mathematical formula for the list-position of $z/n \in \mathbb{Q}$?

- Programs are finite binary strings. We show that all finite binary strings \mathcal{B} are countable

- How do I list them?

- Start with the empty string (duh...)
- Then list all strings of length 1, length 2, etc.

$$\mathcal{B} = \{\varepsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, 011, 100, 101, 110, 111, \dots\}$$

- I now know the exact position of every string!
- **Exercise.** What is the list-position of 0110?
- **Exercise.** For the $(k + 1)$ -bit string $b = b_k b_{k-1} \dots b_1 b_0$, define the string's numerical value:

$$value(b) = b_0 \cdot 2^0 + b_1 \cdot 2^1 + \dots + b_{k-1} 2^{k-1} + b_k 2^k$$

- Show:

$$\text{list-position of } b = 2^{\text{length}(b)} + value(b)$$

- Wait a second... We keep seeing larger and larger sets that are countable!!

$$\mathbb{N}, \mathbb{N}_0, \mathbb{Z}, \mathbb{Q}, \mathcal{B}$$

- SURELY EVERYTHING IS COUNTABLE!?!

Infinite Binary Strings Are Uncountable!

- One of the most cool results in computational theory
- Cantor's Diagonal Argument: Assume there is a list of *all* infinite binary strings

$$\begin{array}{l} b_1: 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 \dots \\ b_2: 0 0 1 1 0 1 0 0 1 0 0 0 0 1 0 \dots \\ b_3: 1 1 0 0 0 0 1 0 0 0 0 1 1 0 0 \dots \\ b_4: 1 0 1 0 0 1 0 0 0 0 1 0 0 0 0 \dots \\ b_5: 0 1 1 0 1 0 1 0 0 0 0 0 0 0 0 \dots \\ b_6: 0 1 0 1 1 0 0 0 1 0 0 0 0 0 0 \dots \\ b_7: 0 0 1 0 0 1 0 0 0 0 0 0 0 0 1 \dots \\ b_8: 0 0 1 0 1 1 0 1 0 0 0 0 1 0 0 \dots \\ b_9: 0 0 0 0 1 1 0 0 0 1 0 0 0 0 0 \dots \\ b_{10}: 1 0 1 1 1 0 1 0 0 1 1 0 0 0 0 \dots \\ \vdots \end{array}$$

- We'll now show that there exists a string that cannot be in that list!
- Look at the (red) diagonal string:

$$b = 0000100101 \dots$$

- What's so special about this string?

– Let's flip the bits

$$\bar{b} = 1111011010 \dots$$

– This string is not in the list!

– Differs from each string b_i in position i

The Real Numbers are Uncountable

- Every real between 0 and 1 has an infinite binary representation and every infinite binary string evaluates to a real number

$$\text{e.g., } 0.001111111111111111 \dots = \frac{1}{2^3} + \frac{1}{2^4} + \frac{1}{2^5} + \frac{1}{2^6} + \dots = \frac{1}{4}$$

- This means

$$|\{\text{reals in } [0,1]\}| = |\{\text{infinite binary strings}\}| > |\mathbb{N}|$$

- Aha, found one!
- **Brain-breaking exercise (Continuum Hypothesis).** Prove that there is no set \mathcal{R} s.t.

$$|\{\text{reals in } [0,1]\}| > |\mathcal{R}| > |\mathbb{N}|$$

- Cantor took on the abstract beast Infinity. (1874)
- ~60 years later, Alan Turing asked the abstract question: What can we compute? (1936)

- For example, consider the set of binary functions f defined on \mathbb{N}

$n:$	1	2	3	4	5	6	7	8	9	10	...
$f(n):$	0	1	1	0	1	0	0	0	1	1	...

- Turns out the set of all such functions is uncountable
 - Corresponds to the set of all infinite binary strings
- Every program is a finite binary string. For example:

```
int main(); //a program that does nothing
```

- This program corresponds to the finite binary string (ASCII code)

```
0110100101101110011101000010000001101101011000010110100101101110001010000010100100111011
```

- So, the number of programs is countable
 - But the number of functions is uncountable!
- There are many more functions than we can write/compute!

- There are MANY MANY functions that cannot be computed by programs!
- Are there interesting, useful functions that cannot be computed by programs?