

Machine Learning & Optimization Lecture 8

- Optimization as minimization
- Convex Sets, Functions ; Convex Optim Probs
- Jensen's Inequality
- First & Second Order Characterizations of Convexity
- Examples
- Operations that preserve convexity
- Examples of Convex Optim Problems

Convex Optimization

- We've seen how to convert ML problems to optim probs using RERM and MLE frameworks
- We also know we only need to ensure

$$R_n(f_T) \leq R_n(\hat{f}_{\hat{T}}) + O(\epsilon)$$

↖ determined by approximate error and generalization error

- In general optimization is NP-Hard
so we restrict ourselves to convex optimization problems (for now)

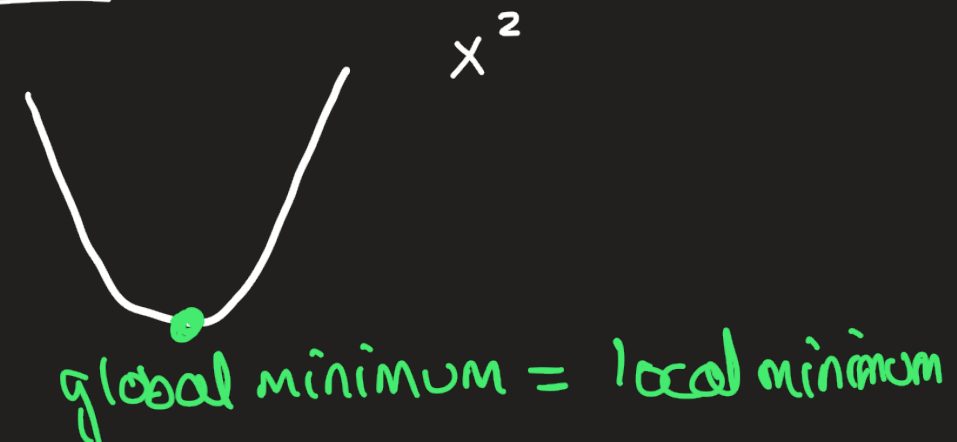
$$x^* = \underset{x \in C}{\operatorname{argmin}} f(x)$$

f - cvx function
 C - cvx set

noncvx



cvx

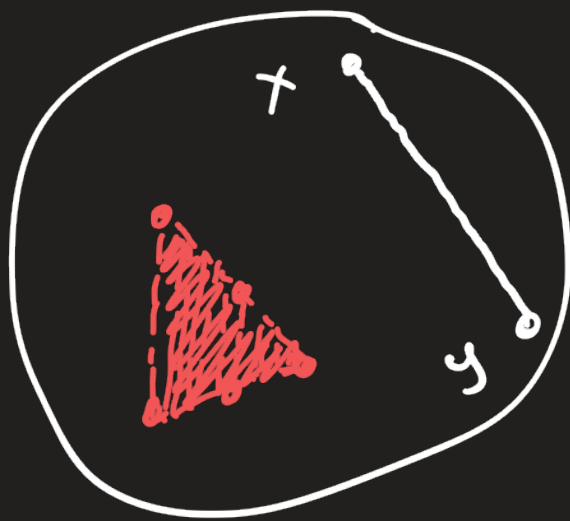


Why restrict to convex optim probs

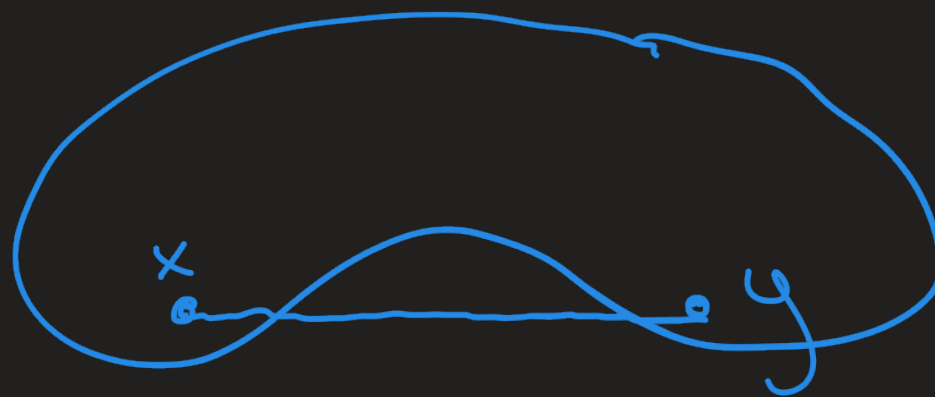
- Very expressive model class
- They are generally tractable
- local minima are global minima
- Many ML problems are convex:
 - GLMs fit using MLE gives convex optim probs
 - many efficient algs exist, tailored for different settings

Convexity

We say a set C is convex if whenever $x, y \in C$, the line segment $[x, y] = \{\alpha x + (1-\alpha)y : \alpha \in [0, 1]\}$ is contained in C .



C - convex



C - nonconvex

We can show that C is convex iff for all $x_1, \dots, x_k \in C$ and any numbers $\theta_1, \dots, \theta_k$ s.t. $\theta_i \geq 0$ for all i and $\sum \theta_i = 1$, the convex combination

$$\theta_1 x_1 + \dots + \theta_k x_k \in C$$

Given $x_1, \dots, x_K$ the set of all convex combinations of these points is called their convex hull

Observe:

If C is a convex set and X is a r.v. taking values in C , then $\mathbb{E}X \in C$

Examples of convex sets

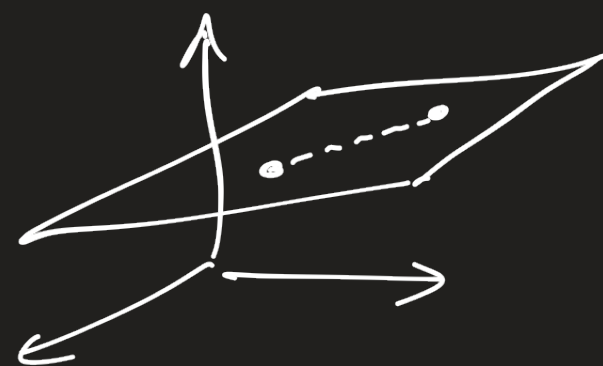
- empty set
- $\mathbb{R}$, $\mathbb{R}_+$, $\mathbb{R}^d$, rays, lines, line segments
- hyperplanes: $\{x \mid a^T x = b\} = H$

if $x \in y$ in this hyperplane, then

$$a^T [\alpha x + (1-\alpha)y] = \alpha a^T x$$

$$+ (1-\alpha)a^T y = \alpha b + (1-\alpha)b = b$$

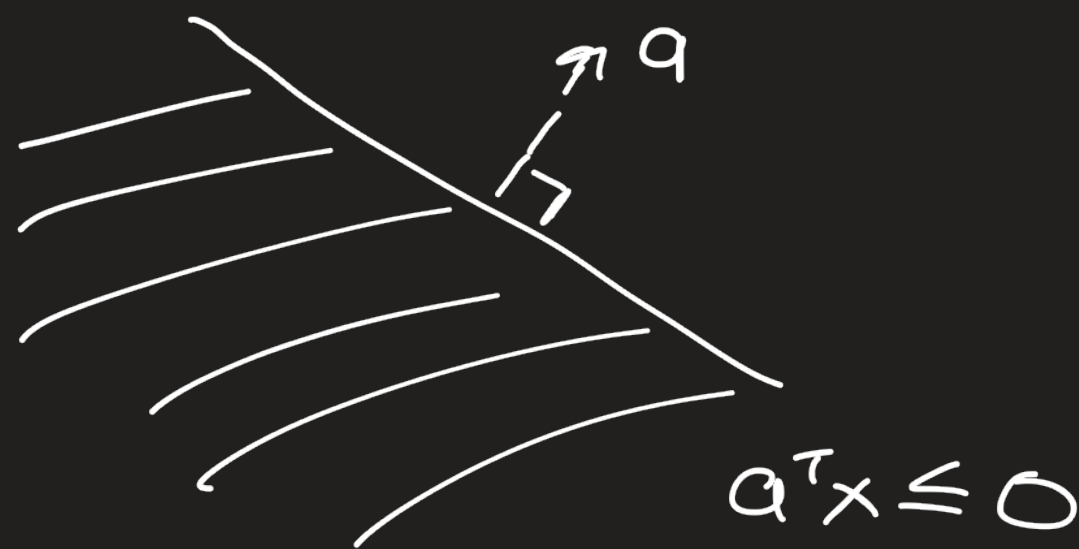
so $[x, y] \subseteq H$



- half-spaces are convex $\{x \mid a^T x \leq b\} = H_{\leq}$

fix $x, y \in H_{\leq}$ and $\alpha \in [0, 1]$,
then

$$\begin{aligned} a^T [\alpha x + (1-\alpha)y] \\ &= \alpha a^T x + (1-\alpha) a^T y \\ &\leq \alpha b + (1-\alpha)b = b \end{aligned}$$



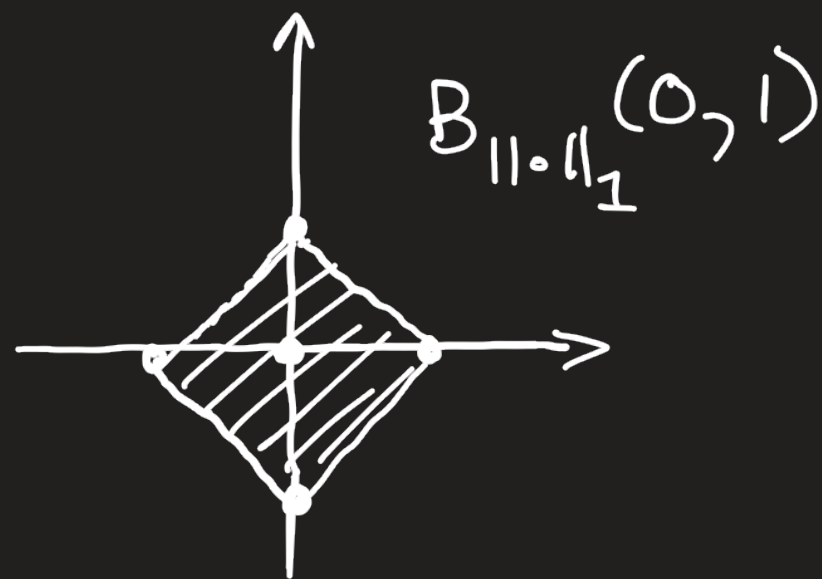
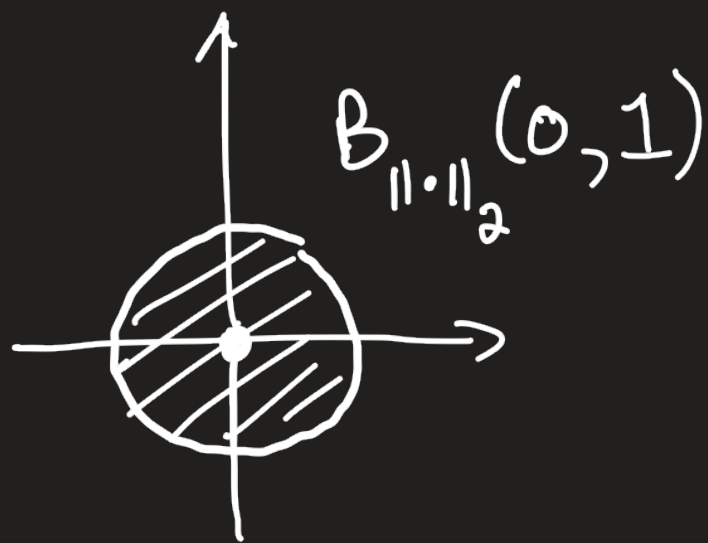
so $[x, y] \in H_{\leq}$

- convex hull of an arbitrary set S

$$\text{conv}(S) = \left\{ \theta_1 x_1 + \dots + \theta_k x_k : x_i \in S, \theta_i \geq 0 \text{ and } \sum \theta_i = 1 \right\}$$

is a convex set. This is the smallest convex set containing S .

- norm balls $B_{\|\cdot\|}(x, r) = \{y \mid \|x - y\| \leq r\}$



- polyhedrons: set of solutions to a finite number of linear equalities and inequalities

$$\{x : a_1^T x \leq b_1$$

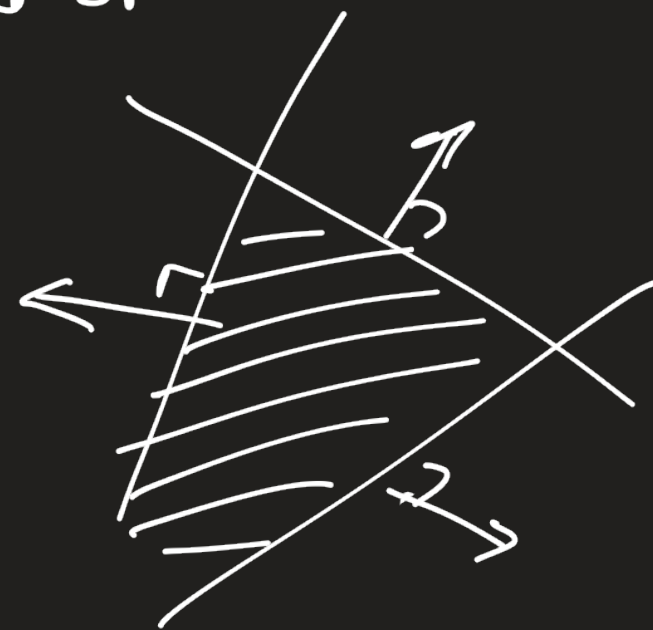
$$a_2^T x \leq b_2$$

$$a_3^T x = b_3$$

$$\vdots$$

$$\}$$

equivalent to $a_3^T x \leq b_3$
and $(-a_3)^T x \leq -b_3$



$$= \{x : Ax \leq b\} \quad \text{where} \quad A = \begin{bmatrix} a_1^T \\ \vdots \end{bmatrix} \quad \text{and} \quad b = \begin{bmatrix} b_1 \\ \vdots \end{bmatrix}$$

- Intersection of arbitrary # of convex sets is convex

$\bigcap_{i \in I} C_i$ is convex if C_i is convex for all $i \in I$

Convex function

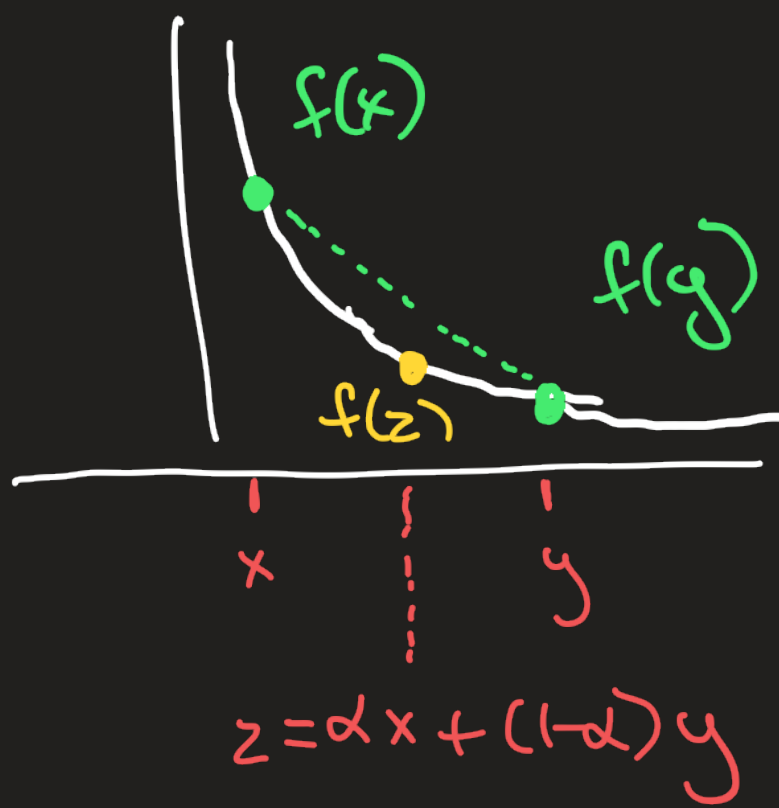
A function f defined on a convex set C is called convex iff it satisfies

$$\underline{f(\alpha x + (1-\alpha)y) \leq \alpha f(x) + (1-\alpha)f(y)}$$

for all $\alpha \in [0,1]$ and $x, y \in C$

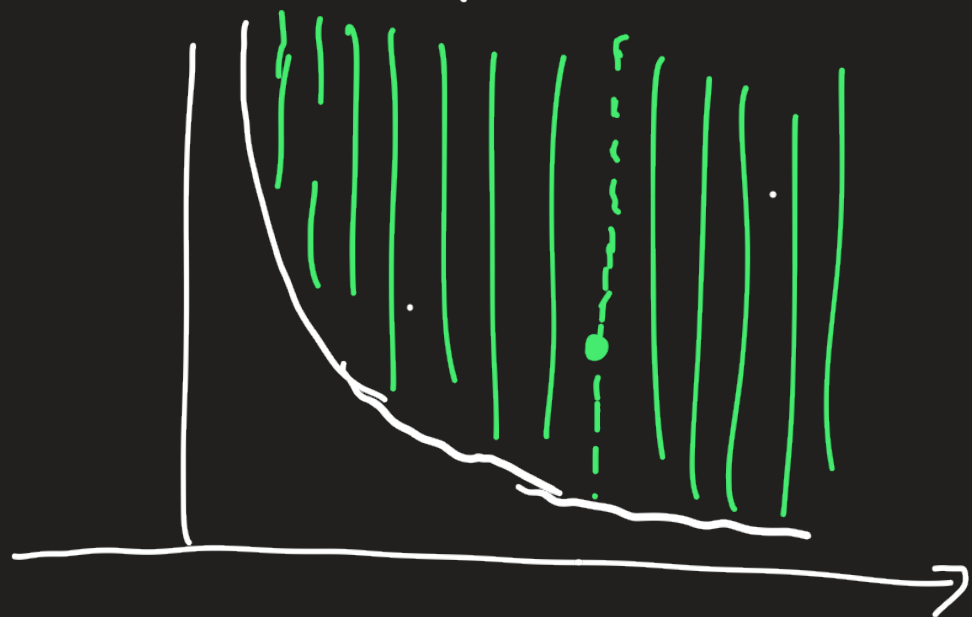
Ex: $\frac{1}{x}$ is convex for $x \in \mathbb{R}_{++}$

Graphically: a function is convex if it lies under its chords



Consider epigraph of f

$$\text{epi}(f) = \{ (x, z) : f(x) \leq z \}$$



Fact: f is convex iff $\text{epi}(f)$ is convex (on the set C)

Prf f convex $\Rightarrow \text{epi}(f)$ convex

Assume f is convex, and (x, z) and $(y, w) \in \text{epi}(f)$ on C .

We need to show that if $\alpha \in [0, 1]$ then $(\alpha x + (1-\alpha)y, \alpha z + (1-\alpha)w) \in \text{epi}(f)$.

To see this note that

$$f(\alpha x + (1-\alpha)y) \leq \alpha f(x) + (1-\alpha)f(y) \leq \alpha z + (1-\alpha)w.$$

so $(\alpha x + (1-\alpha)y, \alpha z + (1-\alpha)w) \in \text{epi}(f)$ as claimed $\square$

Strict convexity $f(\alpha x + (1-\alpha)y) < \alpha f(x) + (1-\alpha)f(y)$

———— $f=0$
not strictly convex

 $(1-x)_+$
not strictly convex

 x^2
strictly convex

Most important consequence of convexity:

If a function f is convex on a convex set C ,
then any local minimum in C is a global minimum

Prf (contradiction)

Assume that x is a local minimizer (there is a neighborhood of x s.t. $f(x) \leq f(y)$ for any y in that neighborhood) but it is not a global minimizer. This means there exists a $z \in C$ such that $f(z) < f(x)$.

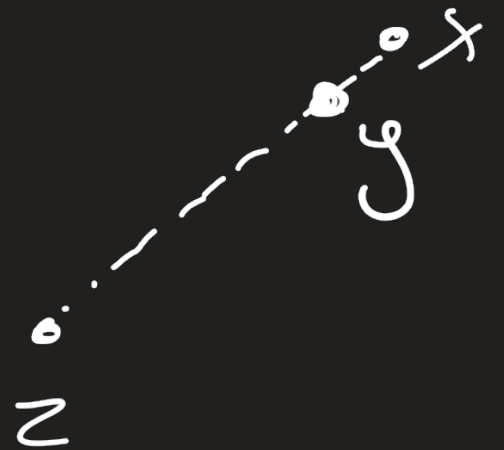
By convexity for an $\varepsilon \in (0, 1]$,

$$y = (1-\varepsilon)x + \varepsilon z$$

then we have by convexity that

$$f(y) \leq (1-\varepsilon)f(x) + \varepsilon f(z) < f(x) \quad \text{because } f(z) < f(x)$$

We've shown that we can a point arbitrarily close to x that satisfies $f(y) < f(x)$. This violates the assumption that x is a local minimizer. We conclude that it's not possible to be a local minimizer without also being a global minimizer in C .

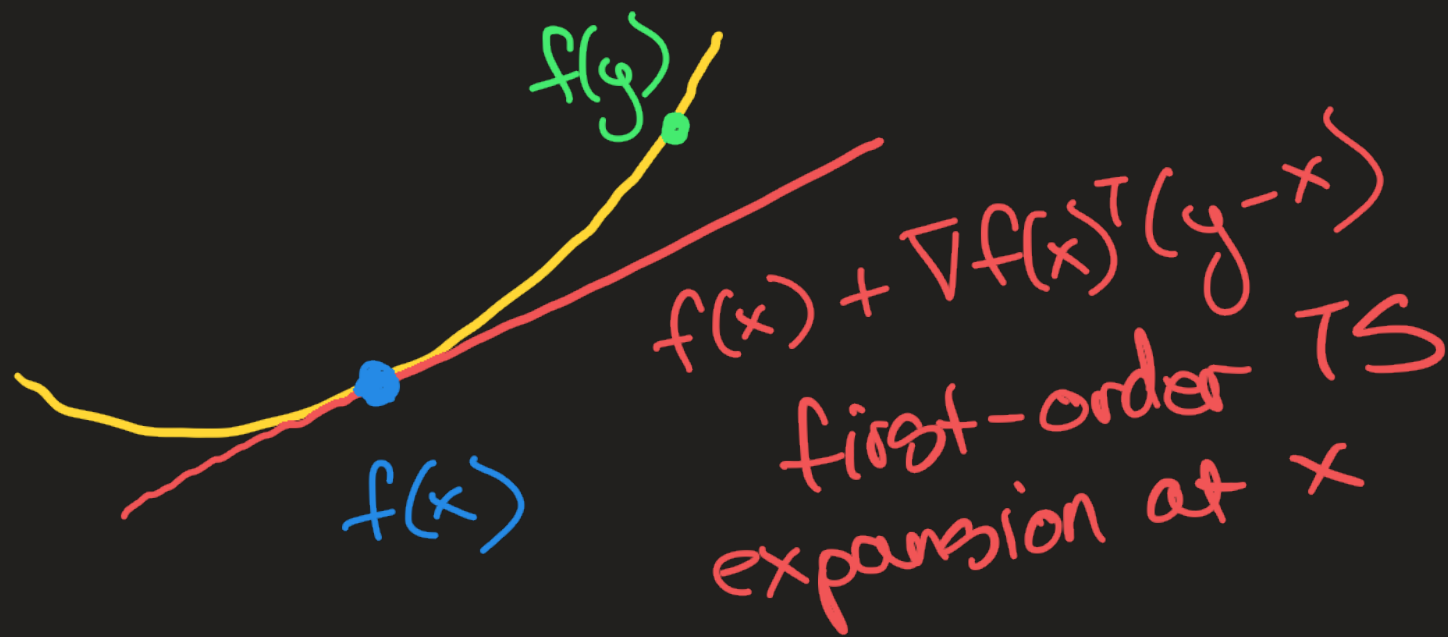


Fact: if f is strictly convex on C then the minimizer of f on C (if any exists) is unique.

First and Second-Order Characterizations of Convexity

1) First-order: if f is differentiable, then f is convex
iff $\forall x, y \in C$

$$f(y) \geq f(x) + \nabla f(x)^T (y - x)$$



That is, f is convex iff it lies above all of its tangent lines

Ex $f(x) = x^2$ is convex on $\mathbb{R}$

note $\frac{df}{dx} = 2x$

Check $f(y) \geq f(x) + \frac{df}{dx}(x)(y-x)$ for all $x, y \in \mathbb{R}$:

$$\Leftrightarrow y^2 \geq x^2 + 2x(y-x)$$

$$\Leftrightarrow y^2 \geq 2xy - x^2$$

$$\Leftrightarrow x^2 + y^2 - 2xy \geq 0$$

$$\Leftrightarrow (x-y)^2 \geq 0 \quad \checkmark$$

conclude that $f(x) = x^2$ is convex on $\mathbb{R}$

2) Second-order condition: if f is twice-differentiable then f is convex iff

$$\nabla^2 f(x) \succeq 0 \quad \left(\text{Hessian is positive semidefinite (PSD)} \right)$$

Recall the Hessian of a function $f: \mathbb{R}^d \rightarrow \mathbb{R}$ is the matrix of second derivatives:

$$\left[\nabla^2 f(x) \right]_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j}(x) = \frac{\partial^2 f}{\partial x_j \partial x_i}(x) = \left[\nabla^2 f(x) \right]_{ji}$$

so $\nabla^2 f(x)$ is always symmetric

Characterizations of PSD matrices M ($M \succeq 0$):

1) M is symmetric and all of its eigenvalues are nonnegative

2) $M = CC^T$ for some matrix C

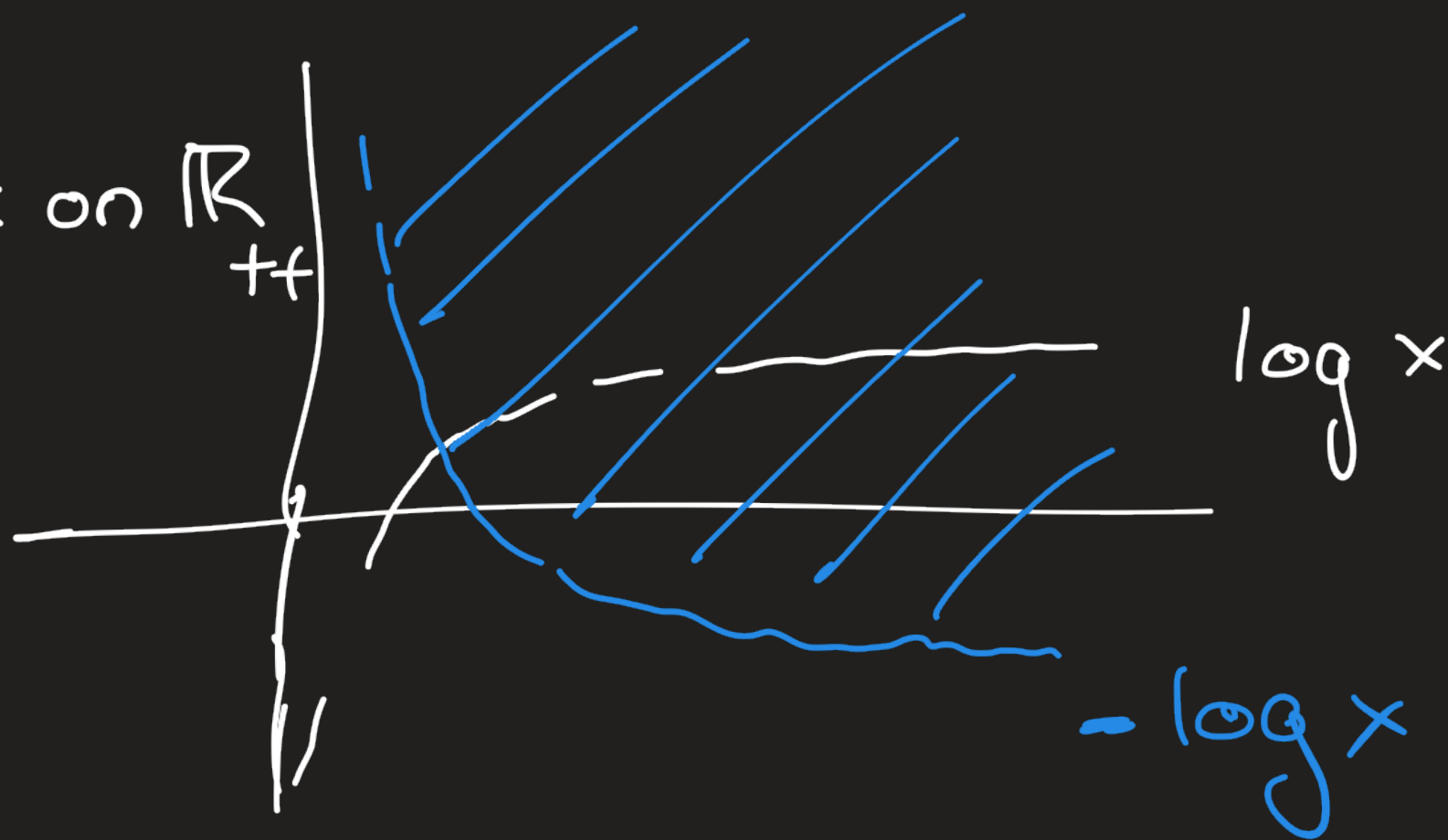
If $f: \mathbb{R} \rightarrow \mathbb{R}$ then $\nabla^2 f = [\nabla^2 f]_{11} = \frac{d^2 f}{dx^2}$

Ex

$f(x) = -\log x$ is convex on $\mathbb{R}_{++}$

$$\frac{df}{dx} = -\frac{1}{x} = -x^{-1}$$

$$\frac{d^2 f}{dx^2} = \frac{1}{x^2} \geq 0 \text{ for } x \in \mathbb{R}_{++}$$



so by the 2nd-order condition, $-\log x$ is convex on $\mathbb{R}_{++}$

Ex (next time) $\log \text{sumexp}(x)$ is convex on $\mathbb{R}^d$

- see Jensen's inequality first

- implies multiclass logistic regression is a convex optim problem