

Today:

- Projects
- Optimality Conds for smooth cvx optim
- Gradient Descent
- Optimality Conds for non-smooth cvx optim
- Subdifferentials
- Subgradient Descent



Smooth convex optimization

f , our convex is differentiable: $\nabla f(x)$ exists at every x in our domain of optimization

Unconstrained Optim

$$\min_{x \in \mathbb{R}^d} f(x) \text{ or just } \min_x f(x)$$

Constrained Optim

$$\min_{x \in C} f(x)$$

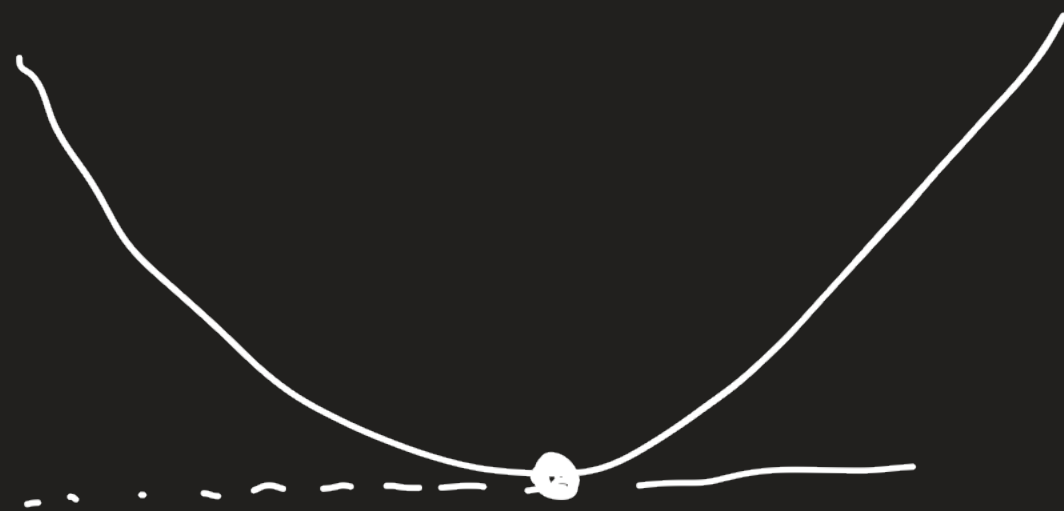
where C is convex

Optimality conditions

Unconstrained: $\nabla f(x^*) = 0$

sufficient & necessary

Sufficiently ($\nabla f(x) = 0 \Rightarrow x$ is a minimizer)



$$\nabla f(x) = 0$$

Recall f is convex & differentiable $\Rightarrow$

for all y

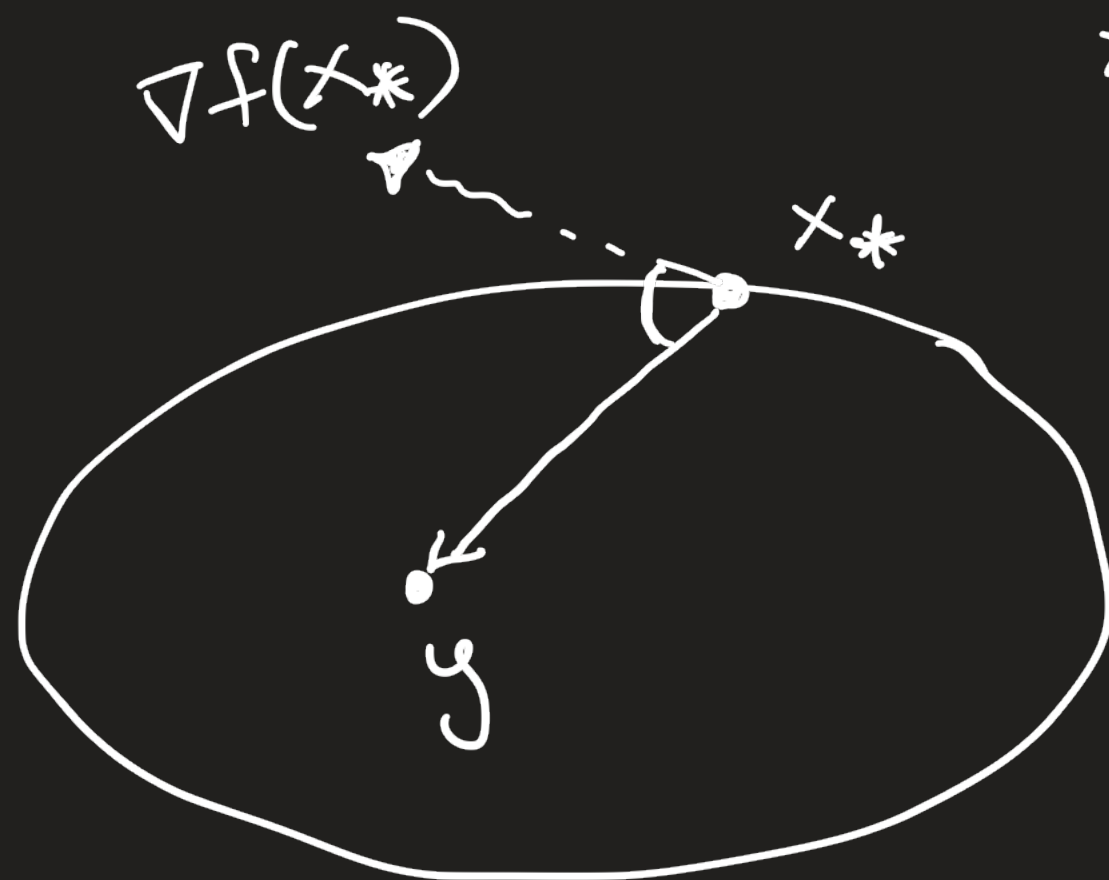
$$f(y) \geq f(x) + \underline{\underline{\nabla f(x)^T (y - x)}}$$

so $\nabla f(x) = 0 \Rightarrow$

$$f(y) \geq f(x)$$

so x is a minimizer.

Constrained Optimization



$$x^* = \operatorname{argmin}_{x \in C} f(x)$$



$$\frac{\langle \nabla f(x^*), y - x^* \rangle \geq 0}{\text{for all } y \in C}$$

Sufficiency:

Assume $\langle \nabla f(x), y - x \rangle \geq 0$ for all $y \in C$

then

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle \geq f(x)$$

so x is a minimizer

Aside inner-products & Cauchy-Schwarz ineq.

$$|\langle x, y \rangle| \leq \|x\|_2 \|y\|_2$$

$$x^T x = \|x\|_2^2$$

Cauchy-Schwarz ineq.
↓ (C-S ineq.)

higher-dim analog of the
fact that

$$|ab| \leq |a||b|$$

$$\Rightarrow \left| \left\langle \frac{x}{\|x\|_2}, \frac{y}{\|y\|_2} \right\rangle \right| \leq 1$$

$\Rightarrow$ can interpret $\left\langle \frac{x}{\|x\|_2}, \frac{y}{\|y\|_2} \right\rangle$ as $\cos \angle x, y$

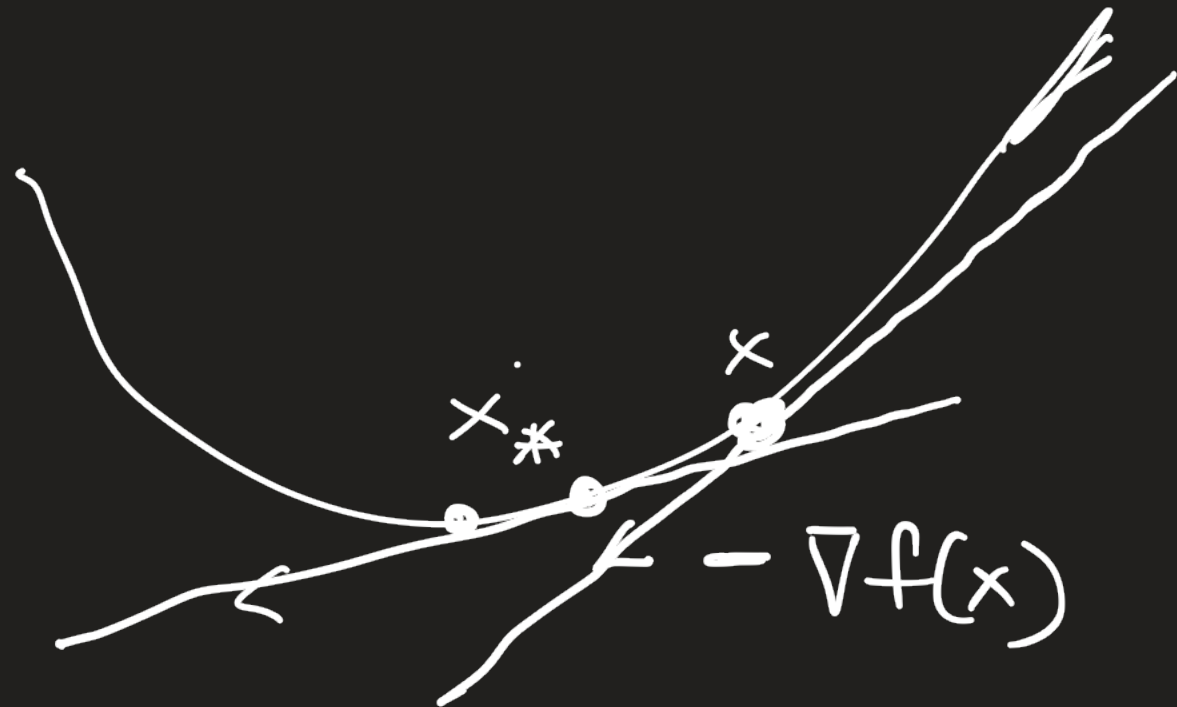
and we have $\langle x, y \rangle = (\cos \angle x, y) \|x\|_2 \|y\|_2$

Note the optimality conditions give us
measures of optimality

Unconst $\|\nabla f(x)\|_2 \rightarrow 0$ as $x \rightarrow x^*$ (*)

Const $\text{dist}(\nabla f(x), S) \rightarrow 0$ (hard-wearing)

Optimization Alg for smooth convex optim: Gradient Descent



The basic idea is to note that

$$f(y) \approx f(x) + \nabla f(x)^T (y-x)$$

so if we want to decrease f by moving to a point

$$y = x + \Delta, \text{ we have}$$

$$f(y) \approx f(x) + \nabla f(x)^T \Delta$$

taking $\Delta = -\alpha \nabla f(x)$, we have

$$f(y) \approx f(x) - \alpha \|\nabla f(x)\|_2^2 \leq f(x)$$

Caveat:

- 1) f is nice
- 2) α is small enough that our linear approx is accurate

Gradient Descent alg

Input: f — convex function on $\mathbb{R}^d$

∇f — gradient

T — number of iterations

x_0 — starting guess

$\alpha_1, \dots, \alpha_T$ — step sizes for each iteration

Output: $x_1, \dots, x_T \in \mathbb{R}^d$ estimates of x_*

for $t=1$ to T

$$x_{t+1} = x_t - \alpha_{t+1} \nabla f(x_t)$$

For now we want to show GD converges under some reasonable assumptions on f ; that tell us

- how to choose our stepsizes
- how many steps we need to take to converge in the sense that

$$f(x_T) \leq f(x_*) + \epsilon$$

A reasonable assumption is that our function's gradient is at least Lipschitz-continuous:

$$\|\nabla f(x) - \nabla f(y)\|_2 \leq \beta \|x - y\|_2$$

We call the gradient β -Lipschitz and we call the function β smooth.

Facts

$$1) \|\nabla f(x)\|_2 \leq \beta \|x - x_*\|_2$$

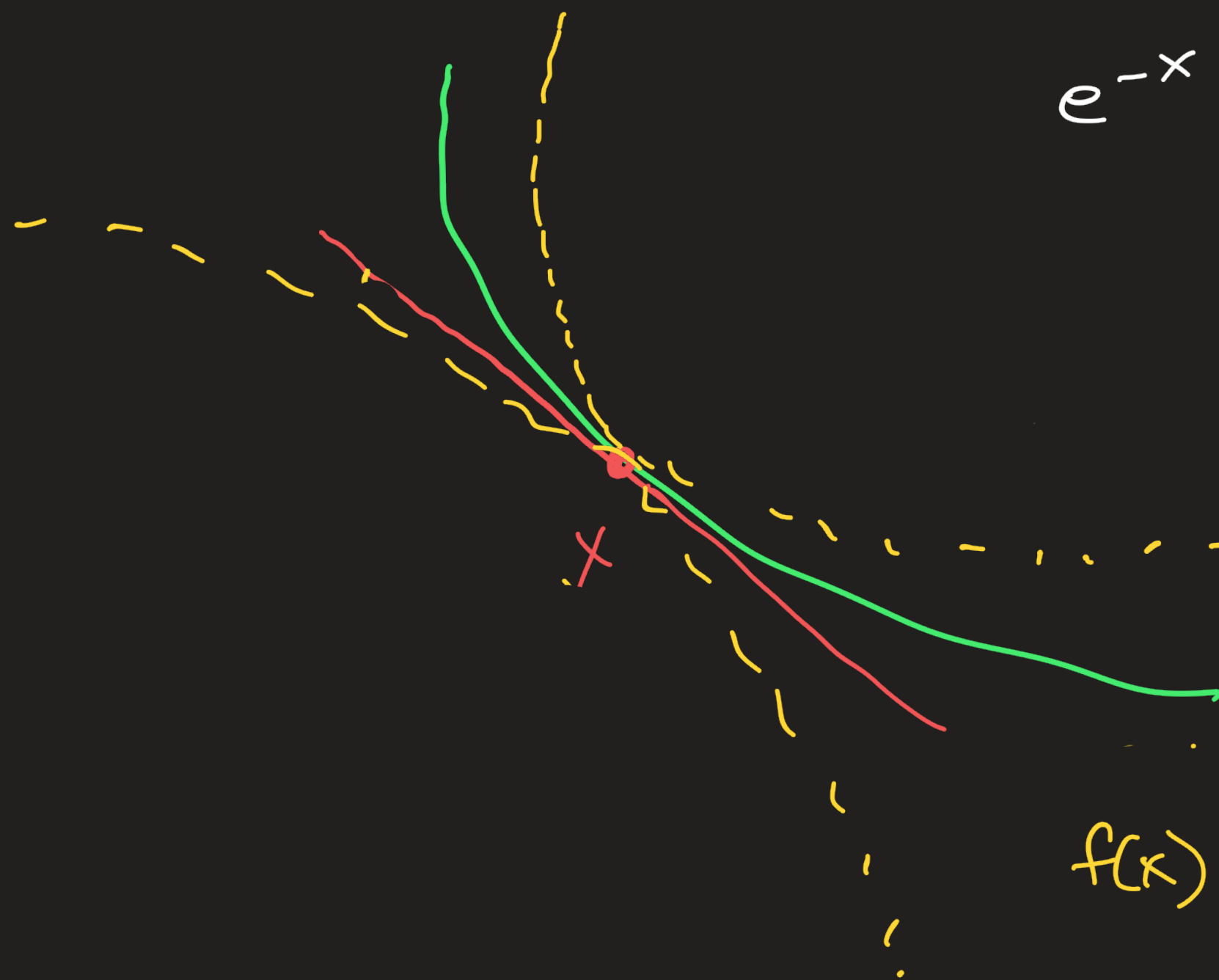
(take $y = x_*$
and note that
 $\nabla f(y) = \nabla f(x_*) = 0$)

2) This implies a concrete bound on how good the approximation

$$f(y) \approx f(x) + \nabla f(x)^T (y - x) \quad \text{is}$$

$$|f(y) - f(x) - \langle \nabla f(x), y-x \rangle| \leq \frac{\beta}{2} \|x-y\|_2^2$$

"f is β -smooth"



e^{-x}

$$f(x) + \nabla f(x)^T(y-x) + \frac{\beta}{2} \|x-y\|_2^2$$

$$f(x) + \nabla f(x)^T(y-x) - \frac{\beta}{2} \|x-y\|_2^2$$

Now assume f is β -smooth. Take $\alpha_t = \frac{1}{\beta}$

Then we see

$$\underbrace{f(x_t - \frac{1}{\beta} \nabla f(x_t))}$$

$$- \left[f(x_t) + \langle \nabla f(x_t), -\frac{1}{\beta} \nabla f(x_t) \rangle \right]$$

$$\leq \frac{\beta}{2} \left\| -\frac{1}{\beta} \nabla f(x_t) \right\|_2^2 = \frac{\beta}{2\beta^2} \left\| \nabla f(x_t) \right\|_2^2$$

$$= \frac{1}{2\beta} \left\| \nabla f(x_t) \right\|_2^2$$

$$f(x_t - \frac{1}{\beta} \nabla f(x_t)) - \left[f(x_t) + \left\langle \nabla f(x_t), -\frac{1}{\beta} \nabla f(x_t) \right\rangle \right] \\ \leq \frac{1}{2\beta} \|\nabla f(x_t)\|_2^2$$

note $\left\langle \nabla f(x_t), -\frac{1}{\beta} \nabla f(x_t) \right\rangle = -\frac{1}{\beta} \|\nabla f(x_t)\|_2^2$

$\Rightarrow$

$$f(x_t - \frac{1}{\beta} \nabla f(x_t)) - f(x_t) \leq -\frac{1}{2\beta} \|\nabla f(x_t)\|_2^2$$

so taking a step of size $\frac{1}{\beta}$ in the negative gradient direction decreases f !

Note: we didn't make use of the fact that the function is convex

Ref: Bubeck. Convex Optimization: Algorithms & Complexity
sec 3.2 for a full proof.

Fact If f is convex & β -smooth then gradient descent with $\alpha_t = \frac{1}{\beta}$ guarantees that

$$f(x_t) - f(x^*) \leq \frac{2\beta \|x_0 - x^*\|^2}{t}$$

rate of convergence is $O\left(\frac{1}{t}\right)$

In general we don't know $\|x_0 - x^*\|_2^2$ so we usually express this result as

$$f(x_T) - f(x^*) \leq \epsilon \quad \leftarrow \text{"}\epsilon\text{-suboptimality"}$$

is achieved when $T = O\left(\frac{P}{\epsilon}\right)$ iterations

What about constrained optimization?

Projected Gradient Descent

$$x_* = \operatorname{argmin}_{x \in C} f(x)$$

Inputs: $f, \nabla f, T, \alpha_1, \dots, \alpha_T, x_0$

P_C — projection operator onto C

Output: $x_1, \dots, x_T$

$$P_C(x) = \operatorname{argmin}_{z \in C} \|x - z\|_2^2$$

Alg:

for t in 1 to T :

$$x_{t+1} = P_C(x_t - \alpha_t \nabla f(x_t))$$

Non-smooth convex optimization

f is convex but not differentiable

Ex:

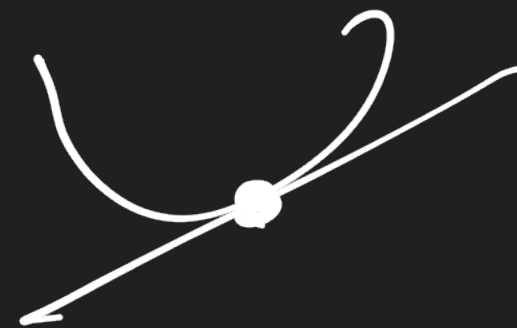
$$F(\omega) = \frac{1}{n} \sum_{i=1}^n (1 - y_i \omega^T x_i)_+ + \lambda \|\omega\|_1$$

In this case, we have no gradient, so can't move in the negative gradient direction to minimize F .

Soln: move in the negative dir of a subgradient

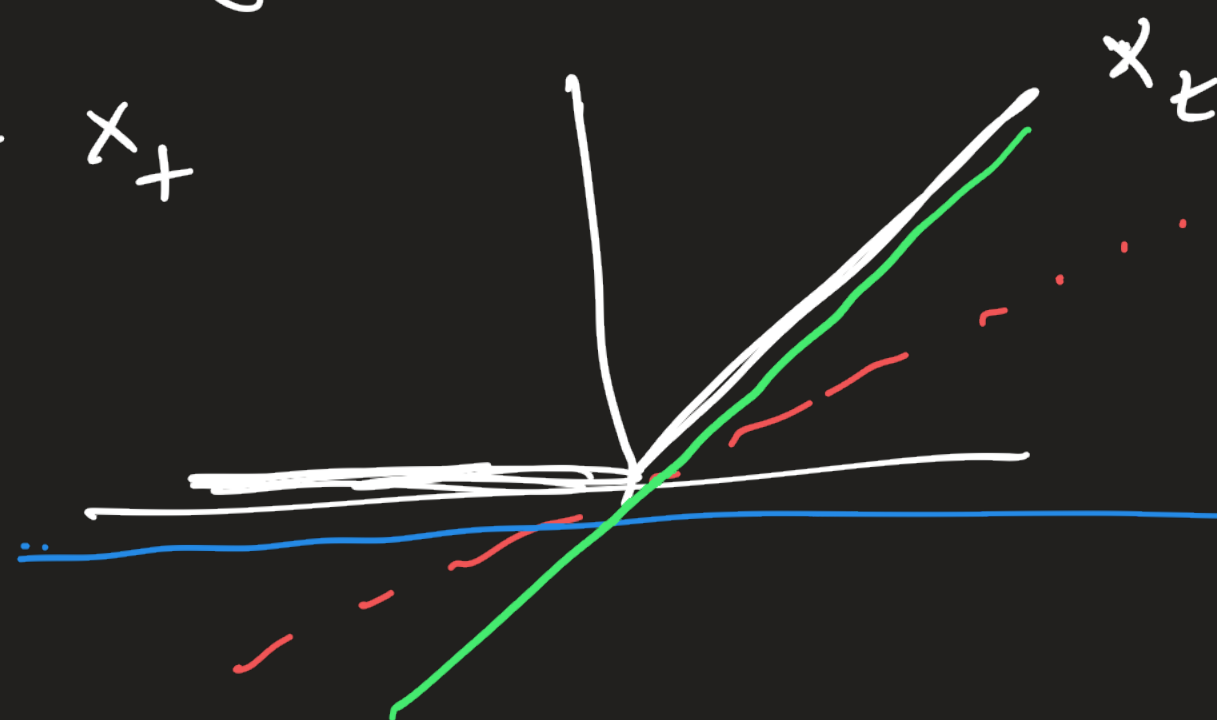
Recall for a smooth convex function

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle$$



If f is convex but not smooth, it may have multiple tangent lines:

consider x_+



multiple tangent lines

In general we define the subdifferential of a function f at x :

$$\partial f(x) = \{ v \mid \text{for all } y, \\ f(y) \geq f(x) + \langle v, y - x \rangle \}$$

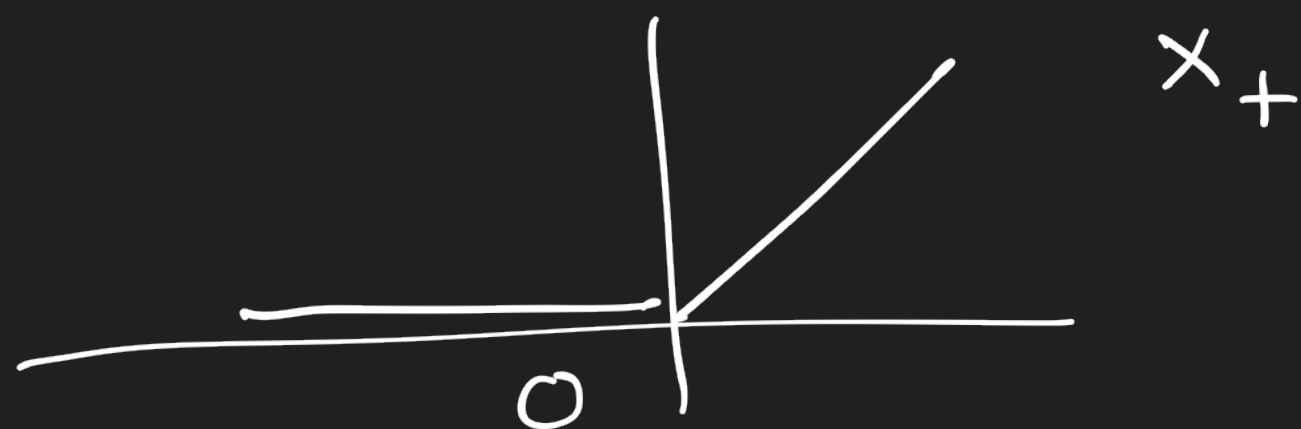
This is a set

Fact For convex functions, the subdifferential is nonempty, and if f is differentiable,

$$\partial f(x) = \{ \nabla f(x) \}$$

We call the vectors in $\partial f(x)$ subgradients

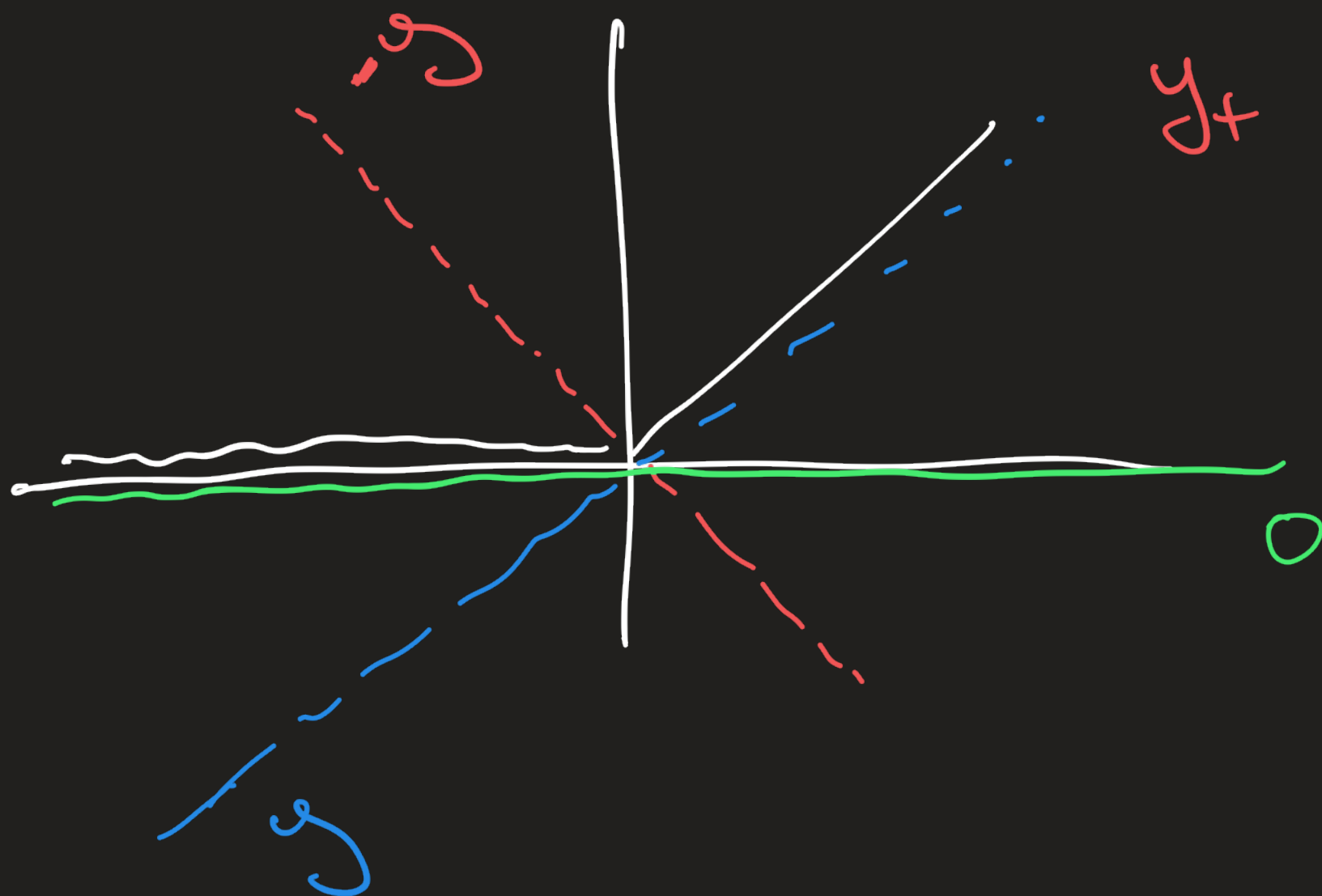
Ex subdifferential of positive part $f(x) = x_+$



Clearly when $x < 0$, $\partial f(x) = \{0\}$
 $x > 0$, $\partial f(x) = \{1\}$

when $x = 0$,

$$\begin{aligned} \partial f(0) &= \left\{ v : f(y) \geq f(0) + v y \right\}^{\text{for all } y} \\ &= \left\{ v : y_+ \geq v y \right\}^{\text{for all } y} \end{aligned}$$



$$\text{try } v = -1 \\ \Rightarrow v \notin \partial f(0)$$

$$\text{try } v = 1 \\ \Rightarrow v \in \partial f(0)$$

$$\text{try } v = 0 \\ \Rightarrow v \in \partial f(0)$$

convince yourself that $\partial f(0) = [0, 1]$

$$\partial f(x) = \begin{cases} 1 & \text{if } x > 0 \\ [0, 1] & \text{if } x = 0 \\ 0 & \text{if } x < 0 \end{cases}$$

Optimality Conditions (Necessary & Sufficient)

Unconstrained

$$x_* = \underset{x}{\operatorname{argmin}} f(x)$$

$$\Leftrightarrow 0 \in \partial f(x_*)$$

Constrained

$$x_* = \underset{x \in C}{\operatorname{argmin}} f(x)$$

$\Leftrightarrow$ there is a subgradient g
in the subdifferential at x_* ,
 $g \in \partial f(x_*)$, for which

$$\langle g, y - x_* \rangle \geq 0$$

for all $y \in C$

Ex $f(x) = |x|$

$$\partial f(x) = \begin{cases} 1 & x > 0 \\ [-1, 1] & x = 0 \\ -1 & x < 0 \end{cases}$$

