Talk Outline

Greedy Hot-Potato Routing

One-Bend Algorithm

Multi-Bend Algorithm

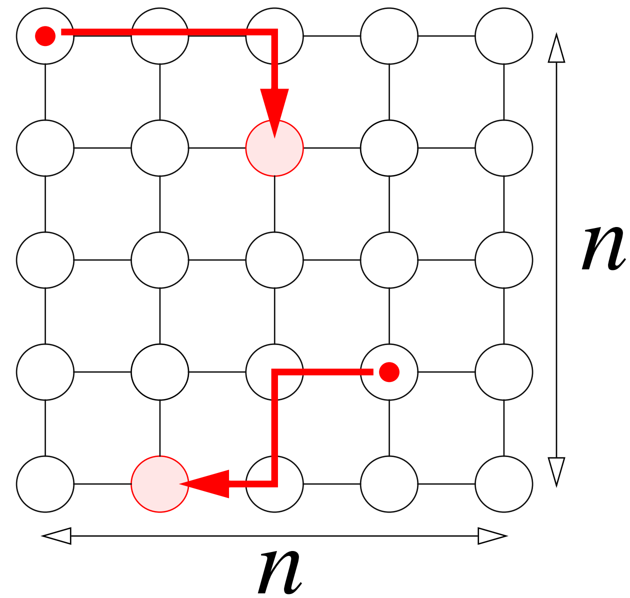
Future Research

Routing

Routing is essential: Internet,
multiprocessors

We consider multiprocessors

Our network: **Mesh**



Synchronous network

From a node to a link:
at most one packet per time step

Local decisions

Hot-Potato Routing

Definition: no buffers

Observation: packets are forwarded immediately

Applications

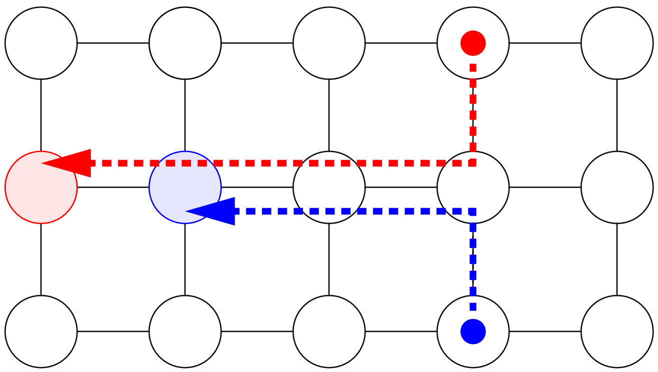
Optical networks

light is difficult to store

Non-optical networks

simple hardware

Conflicts in hot-potato routing



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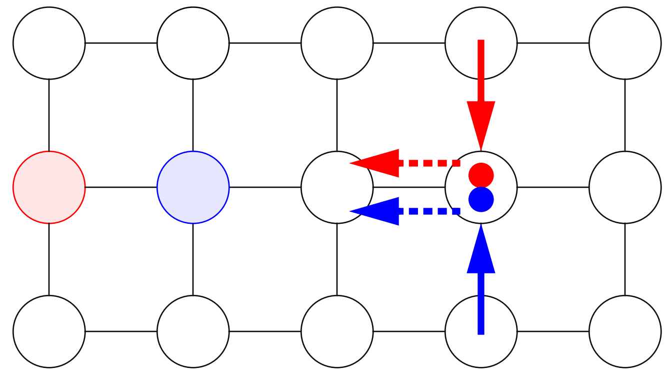

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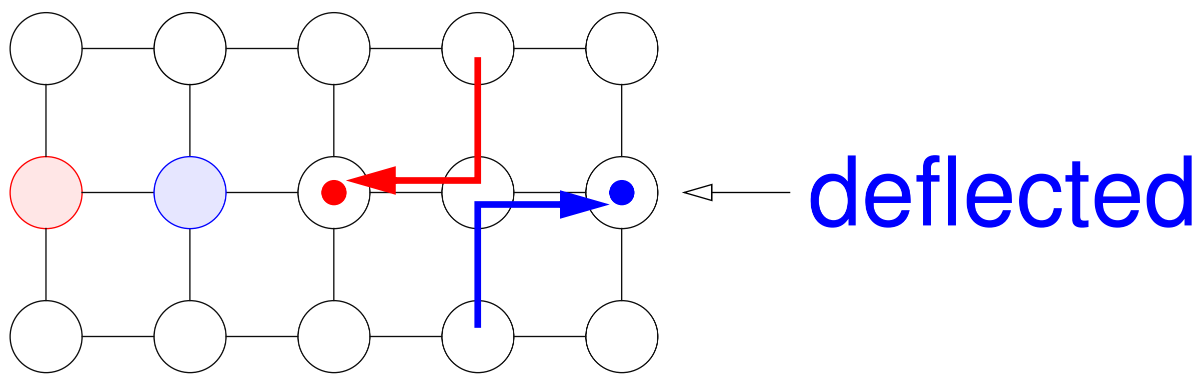


(3)

Hot-potato algorithms

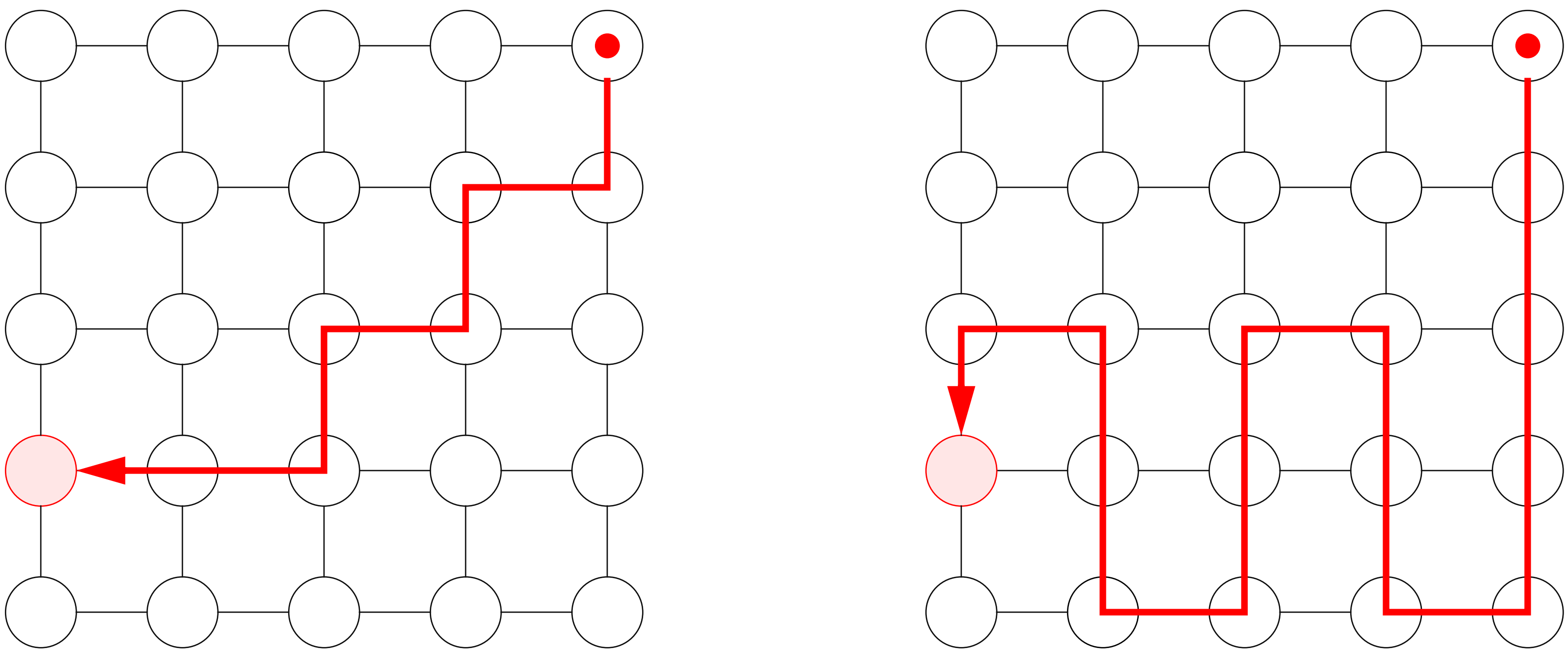
Specify how the packets move

Specify how the conflicts are resolved

Greedy Algorithms

Packets prefer links towards destination

Greedy　　　　Non-Greedy



Advantages of greedy algorithms

Adaptive

Simple implementations

Work very well in practice

Maxemchuk INFOCOM'89

Challenge: hard to analyze

Our Contribution

New greedy hot-potato algorithms

Better bounds for batch routing problems

Batch routing: at time 0 each node
injects a packet

Question: how much time until all
packets reach destinations?

Interesting batch problems

Permutation

Random destinations

General batch problem

Previous bounds for greedy algorithms

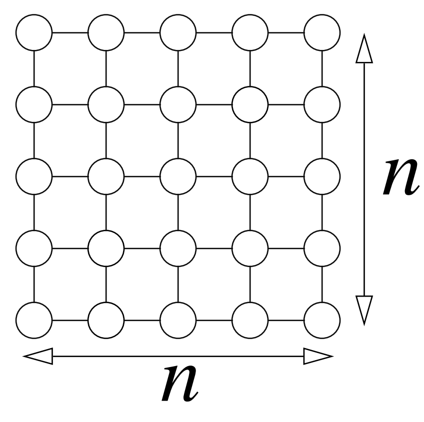
Ben-Dor, Halevi, Schuster PODC '94

Borodin, Rabani, Schieber TPDS '97

Permutation:
Random destinations: $O(n^2)$

Far from lower bound: $\Omega(n)$

Deterministic algorithms



Our bounds: one-bend algorithm

Busch, Herlihy, Wattenhofer SODA '00

Permutation:
Random destinations: $O(n \cdot \log n)$

Close to lower bound: $\Omega(n)$

With high probability: $1 - \frac{1}{n}$

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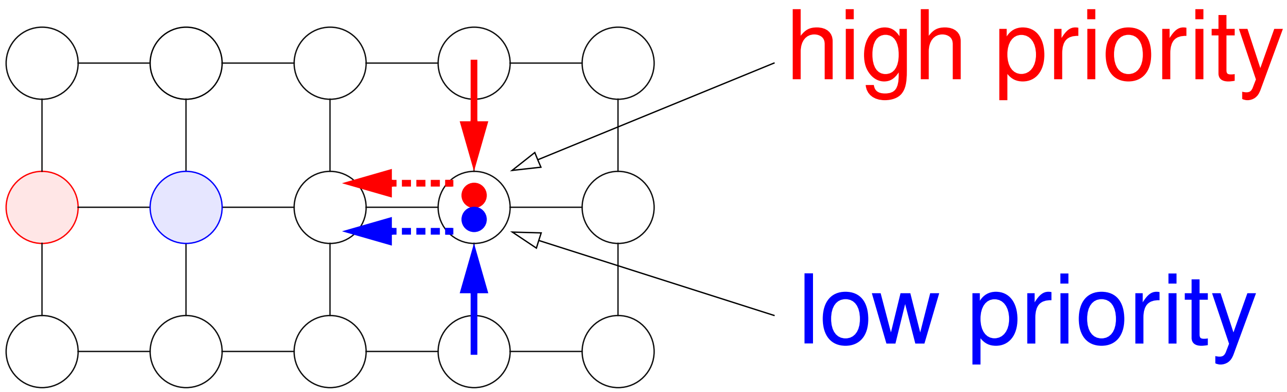
Future Research

One-Bend Algorithm

Packet States	Priority
running	higher
excited	
normal	lower

Conflict

(1)



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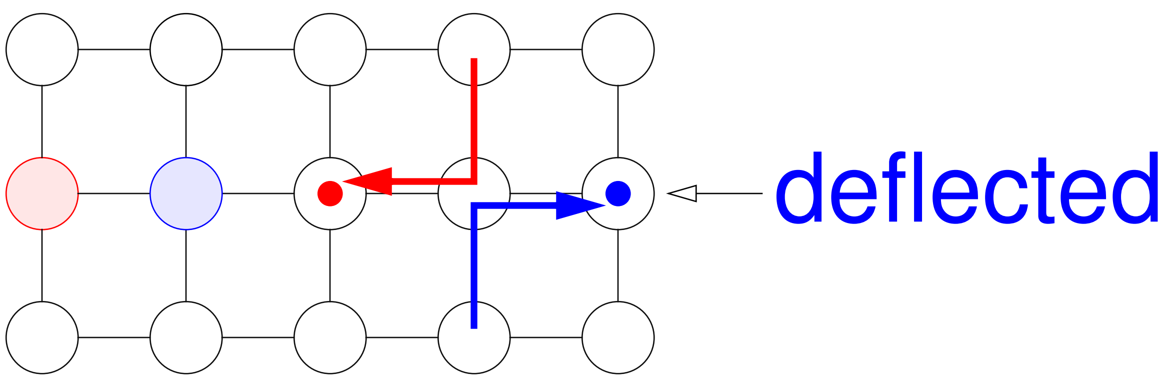
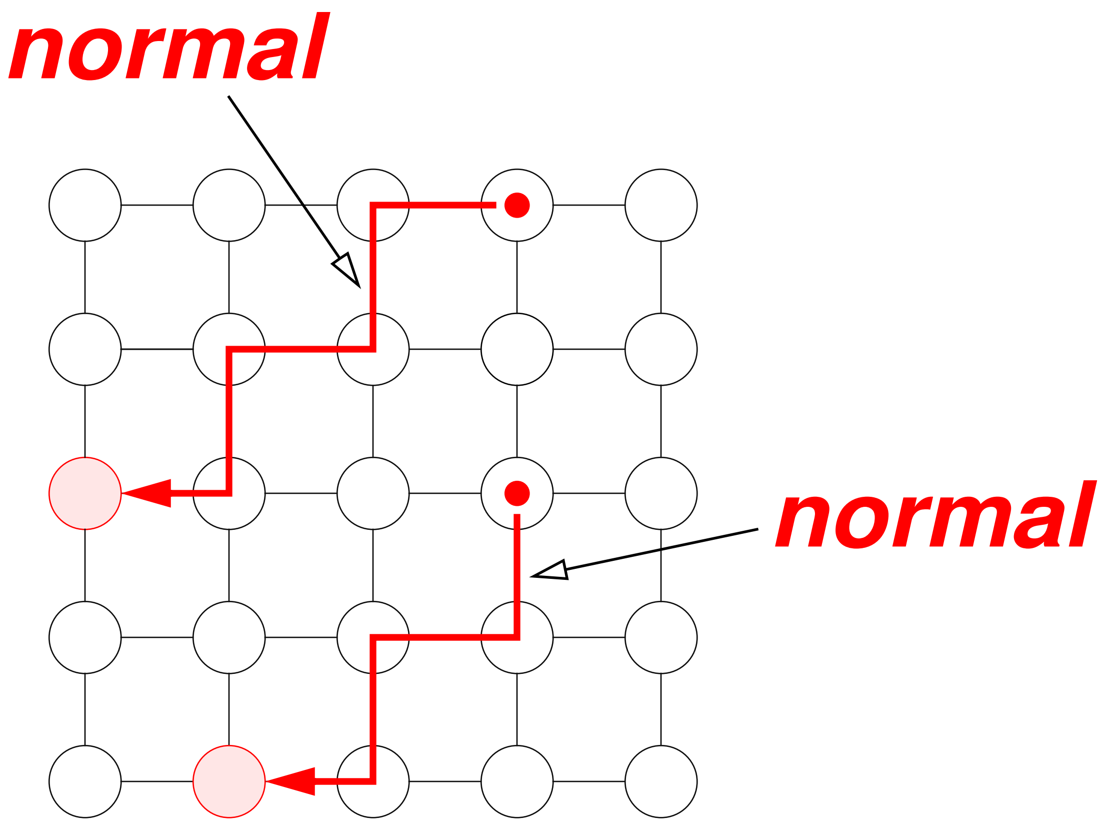
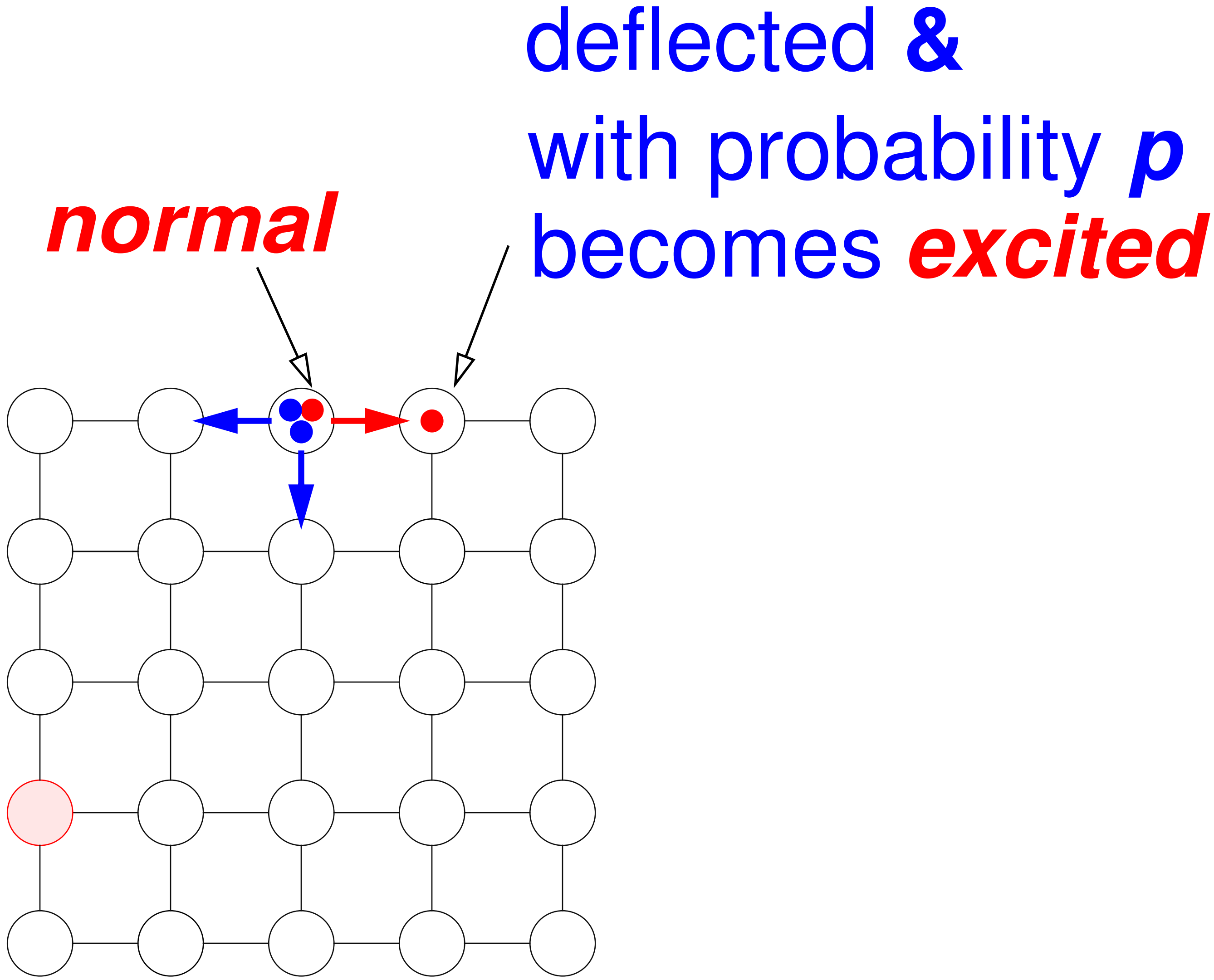


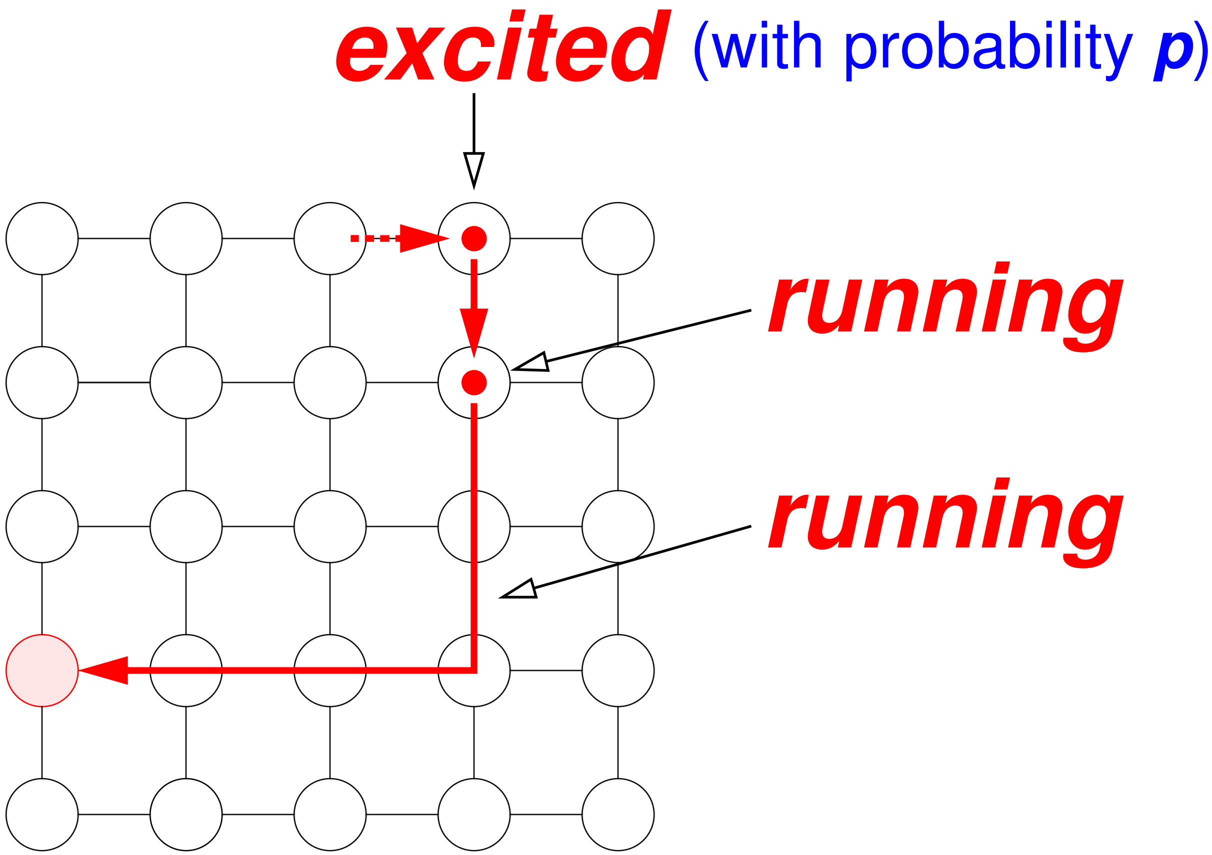
Initially





deflected **&**
with probability ***p***
becomes ***excited***
normal

Home run



If interrupted it becomes ***normal***

Time Analysis for a Packet

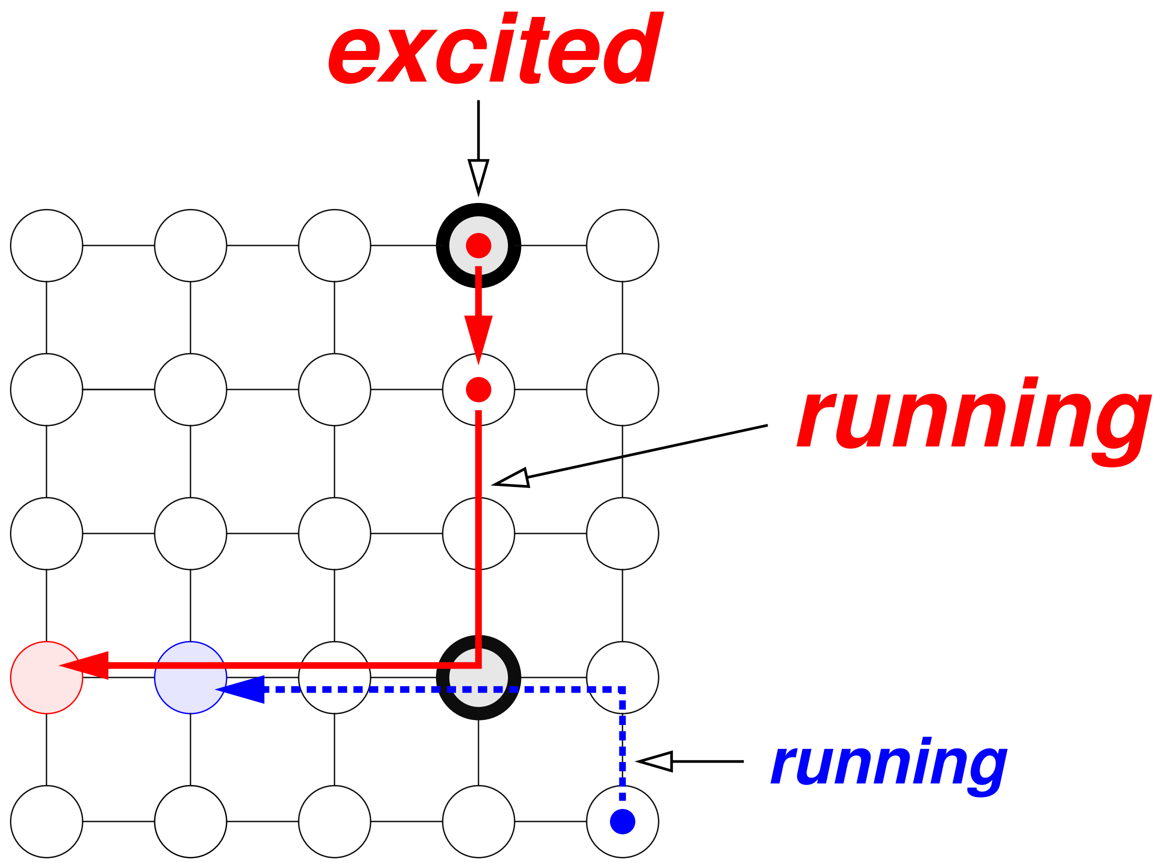
Permutation, Random destinations

Total time until a packet reaches destination:

$O(n)$ expected

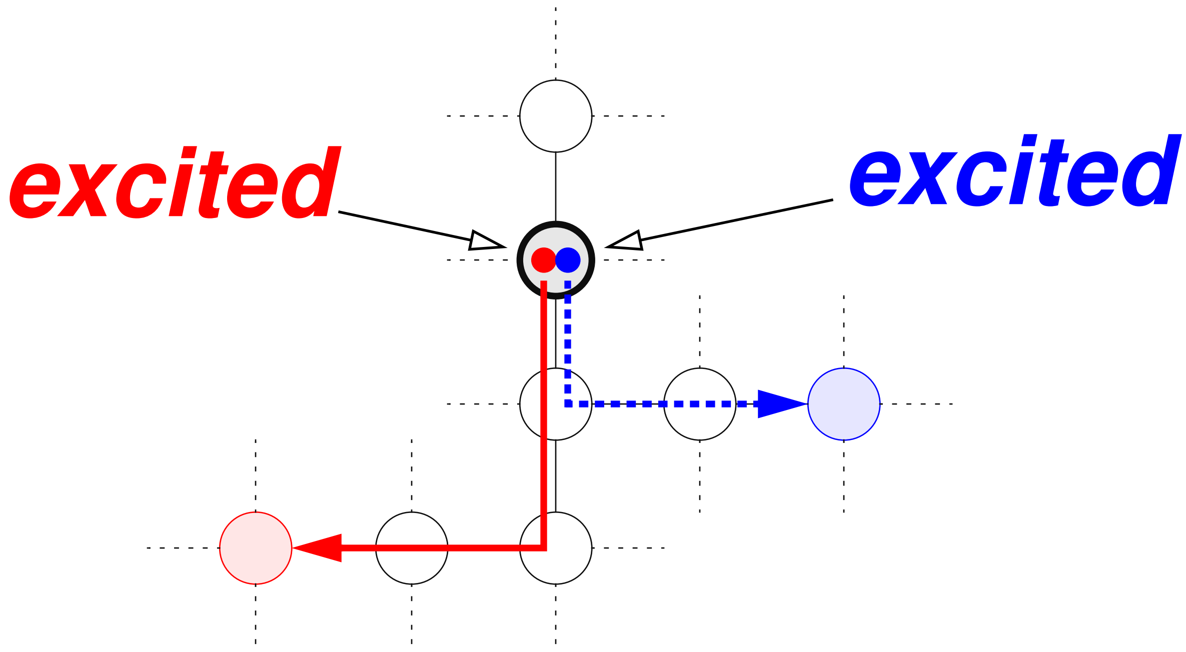
$O(n \cdot \log n)$ with high probability

Interrupting a home run



Interrupting a home run

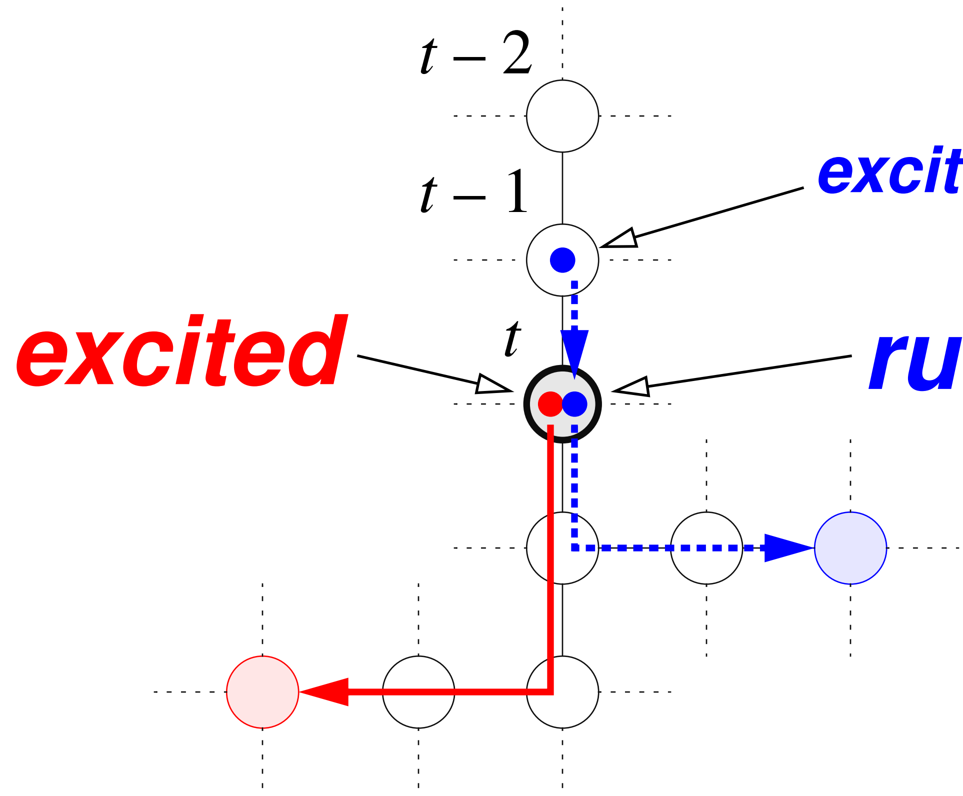
Case1: ***excited*** conflicts ***excited***



Probability of no conflict $\geq (1-p)^3$

Interrupting a home run

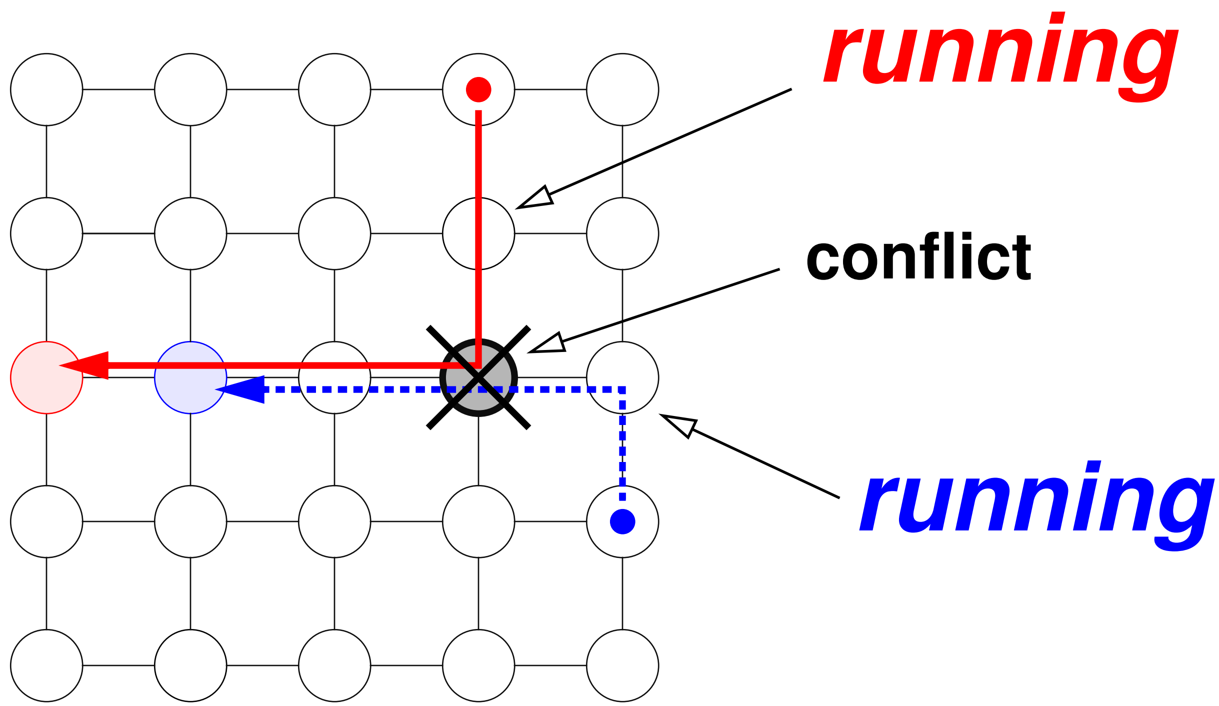
Case 2: ***excited*** conflicts ***running***



Probability of no conflict $\geq (1-p)^{4 \cdot n}$

Interrupting a home run

Case 3: ***running*** conflicts ***running***



Probability of no conflict $\geq (1-p)^n$

Succeeding in a home run

Take $p = \frac{1}{n}$

Probability of success:

$$\underbrace{(1-p)^3}_{\text{case 1}} \cdot \underbrace{(1-p)^{4 \cdot n}}_{\text{case 2}} \cdot \underbrace{(1-p)^n}_{\text{case 3}} \geq \underbrace{\frac{1}{c}}_{\text{constant}}$$

Success after a deflection

deflected $\longrightarrow$ ***excited*** $\longrightarrow$ successful home run

Probability of success: $$p \cdot \frac{1}{c} = \frac{1}{n} \cdot \frac{1}{c}$$

Expected deflections: $$n \cdot c = O(n)$$

Total time for a packet

Expected deflections $O(n)$ $\Rightarrow$ Expected total time $O(n)$

With similar analysis:

Total time: $O(n \cdot \log n)$

With probability $\geq 1 - \frac{1}{n^3}$

Total time for all packets

Permutation:
Random destinations:

$$O(n \cdot \log n)$$

With high probability $\geq \left(1 - \frac{1}{n^3}\right)^{n^2} \geq 1 - \frac{1}{n}$

Lower bound: $\Omega(n)$

Previous algorithms: $O(n^2)$